

# COMP2421 – DATA STRUCTURES

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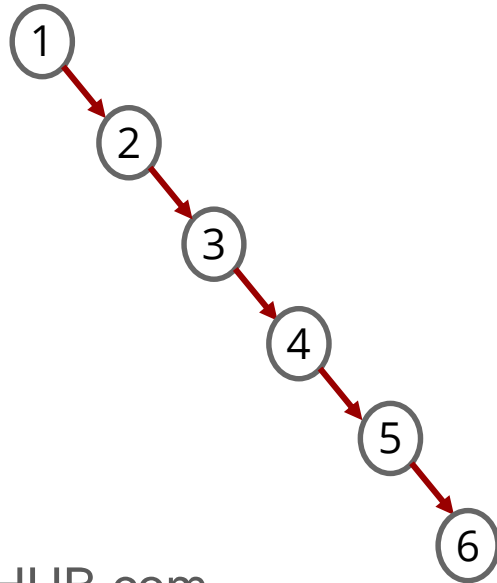
## AVL Trees

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# AVL

- AVL (Adelson-Velskii and Landis) tree is a binary search tree with additional balance property.
- It actually ensures that the depth of the tree is  $O(\log n)$ .
- Why do we need balance?



## AVL (2)

- Height of the left subtree and height of the right subtree differ by at most 1.
- For a tree to be balanced, the value has to only be -1, 0, 1 at each node.
- $\text{Balance}(\text{tree}) = \text{height of } (\text{tree.left}) - \text{height of } (\text{tree.right})$   
(height of the root's left branch and the height of the root's right branch).

## AVL (3)

- This means that all elements in the tree can be ordered in some consistent manner.

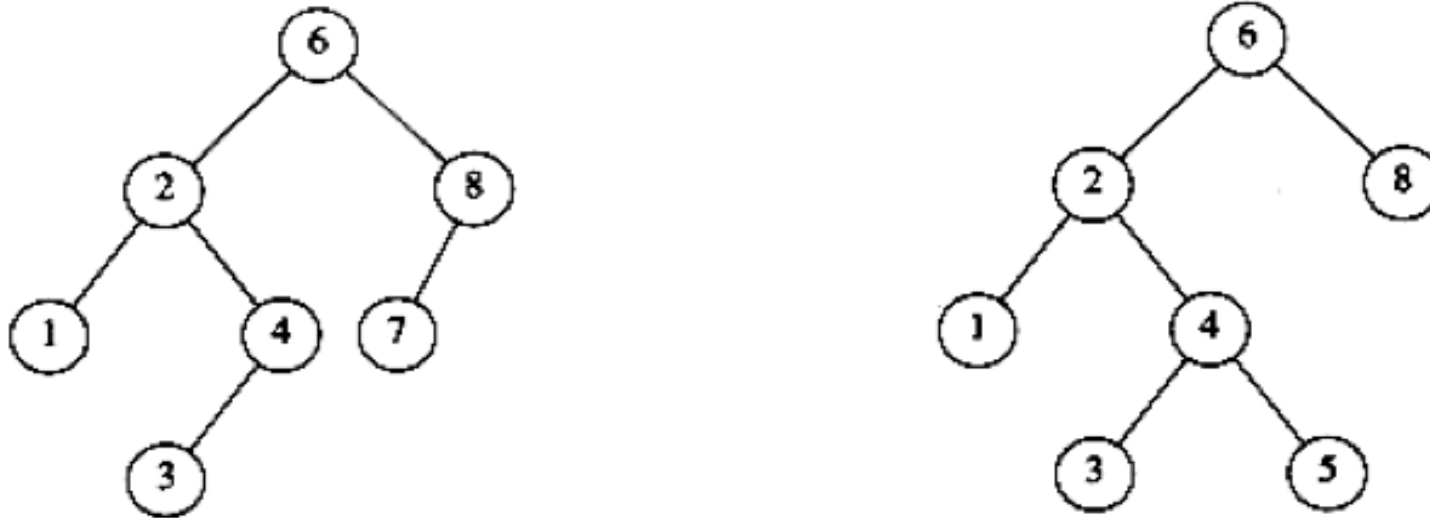


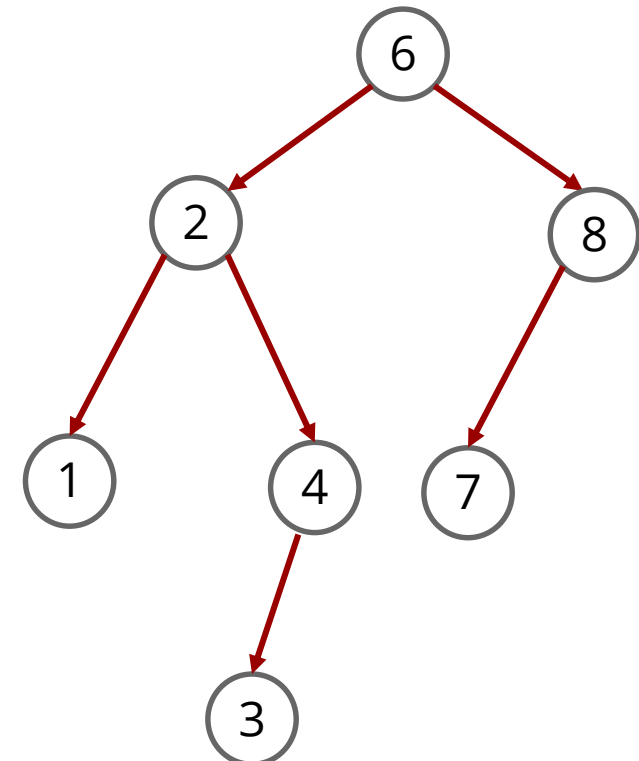
Figure 4.29 Two binary search trees. Only the left tree is **AVL**.

# Rotations

- The balance of an AVL-tree is made via rotations.
- **The algorithm:** start at the node inserted and travel up the tree, updating the balance information at every node on the path.
- If we get to the root without having found any badly balanced nodes, then nothing to update.
- Otherwise, we do rotation at the first badly balanced node found, adjust the balance, and we are done.

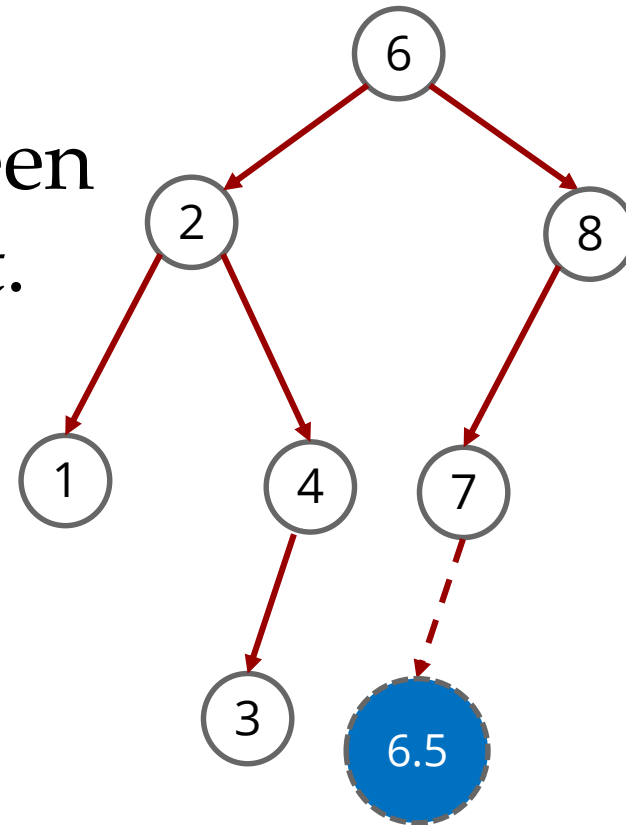
## Rotations (2)

- If the insertion causes some node in the AVL tree to lose balance, perform rotation at that node. In most cases, this is sufficient to rebalance the tree.
- There are two types of rotations
  - Single Rotation
  - Double Rotation
- Example: Insert the node 6.5 to the following tree



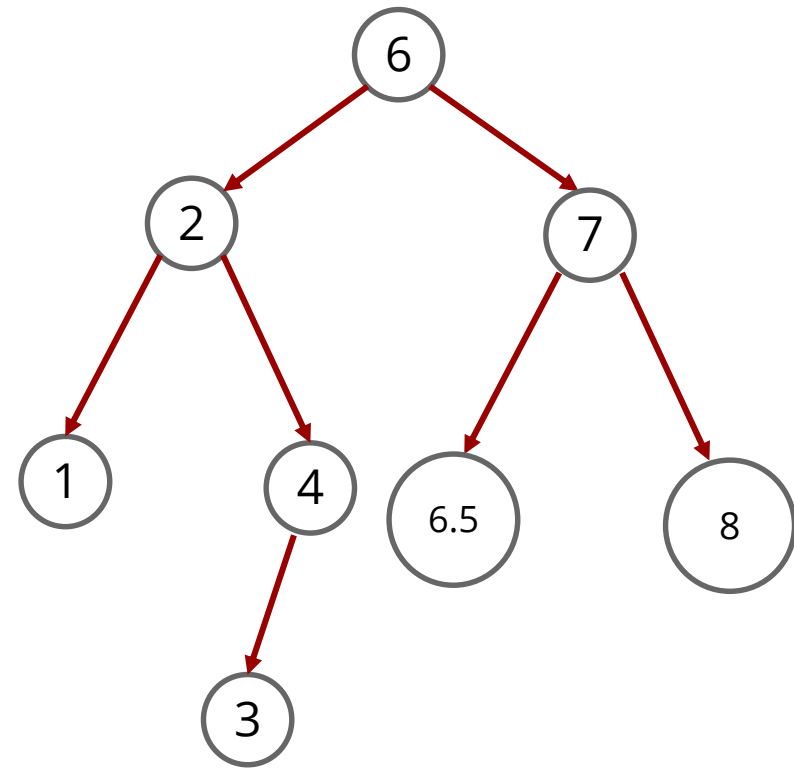
## Rotations (3)

- After inserting 6.5, the original AVL tree on the left, node 8 becomes unbalanced.
- Thus, we do a single rotation between 7 & 8 obtaining the tree on the right.



## Rotations (4)

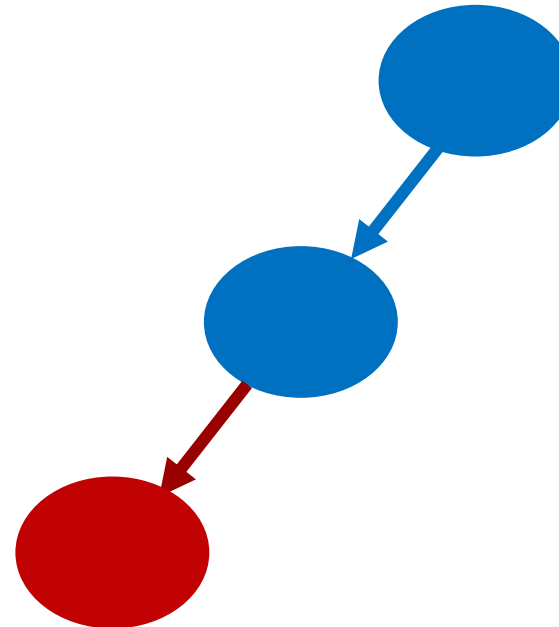
- After rotation the tree becomes





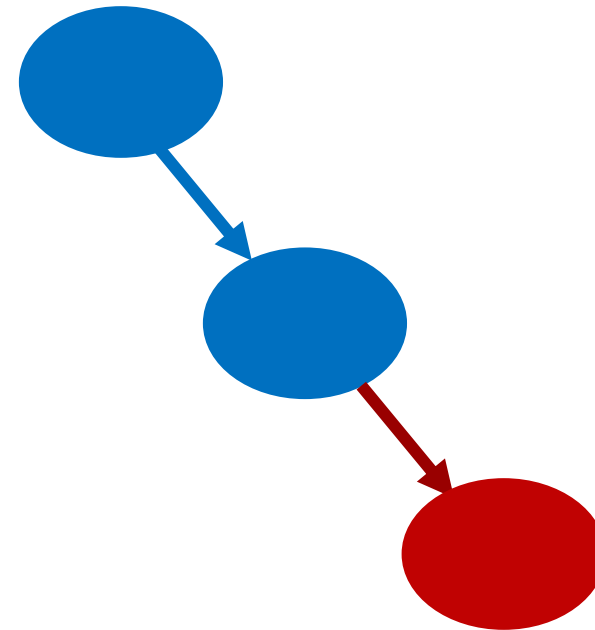
## Rotations (5)

- To re-balance a tree there are 4 different situations:
  1. insert into the left sub-tree of the left child (Right single rotation)



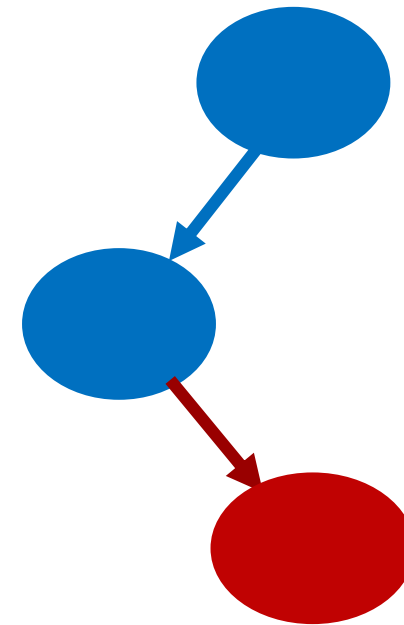
## Rotations (5)

- To re-balance a tree there are 4 different situations:
  1. insert into the left sub-tree of the left child (Right single rotation)
  2. insert into the right sub-tree of the right child (Left single rotation)



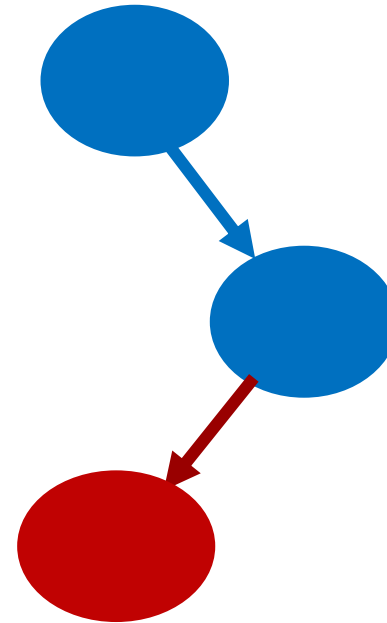
## Rotations (5)

- To re-balance a tree there are 4 different situations:  
3. insert into the right of the left sub-tree (LR - double rotation)



## Rotations (5)

- To re-balance a tree there are 4 different situations:  
4. insert into the left of the right sub-tree (RL – double rotation)



# Example

- Create an AVL tree and insert the values from 1 to 7.
- Insert( **1** )

# Example

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- Insert( **1** )



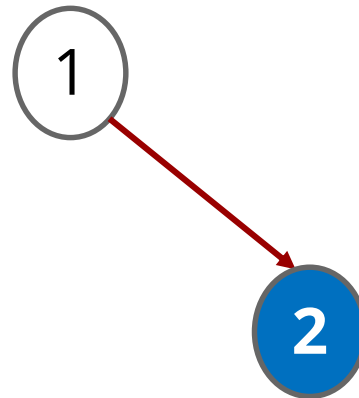
# Example

- Create an AVL tree and insert the values from 1 to 7.
- Insert( **2** )



# Example

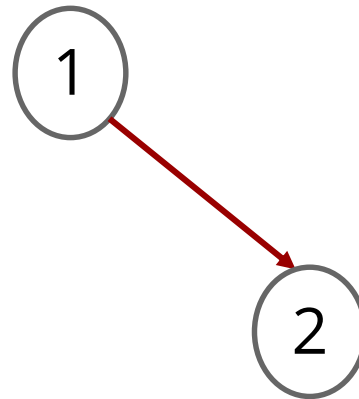
- Create an AVL tree and insert the values from 1 to 7.
- Insert( **2** )





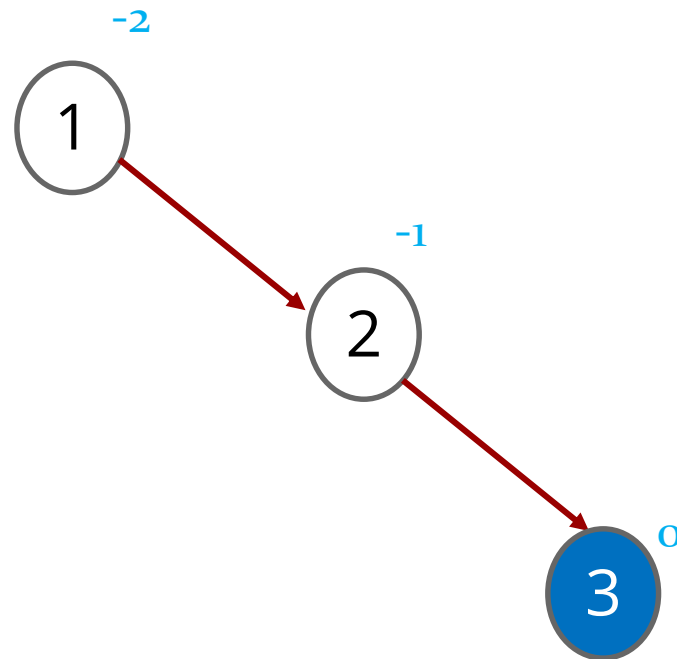
# Example

- Create an AVL tree and insert the values from 1 to 7.
- Insert( **3** )



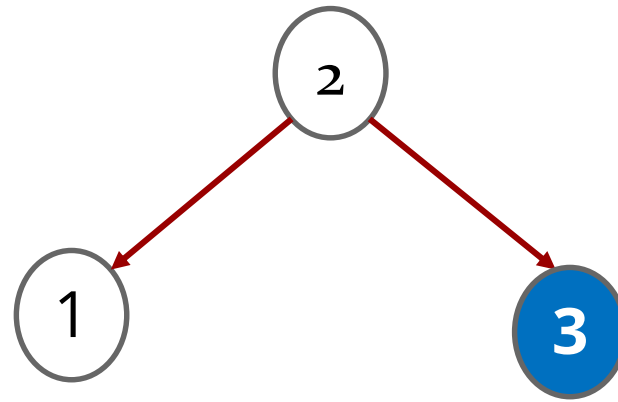
# Example

- Create an AVL tree and insert the values from 1 to 7.
- Insert( **3** )



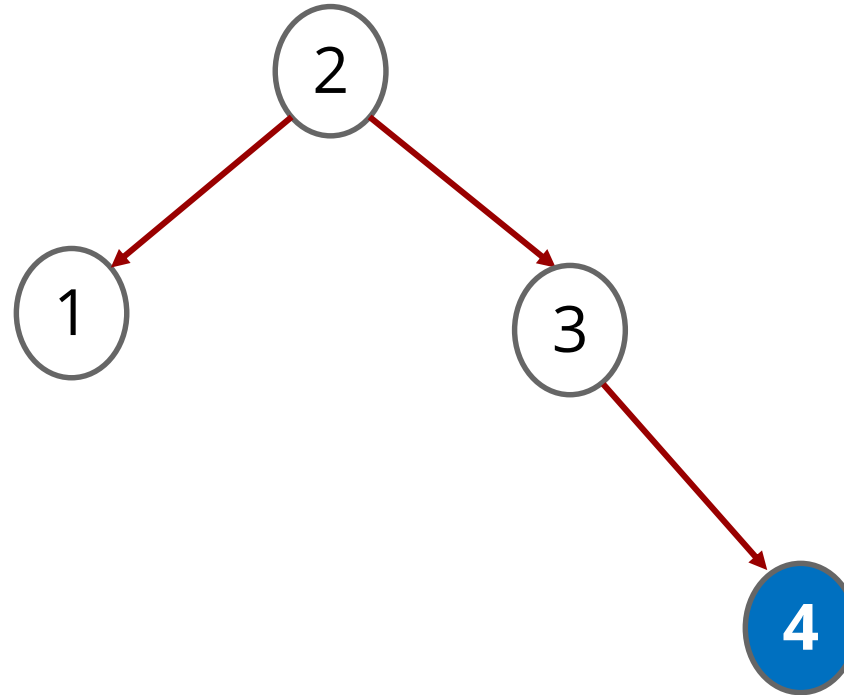
# Example

- Insert( **4** )



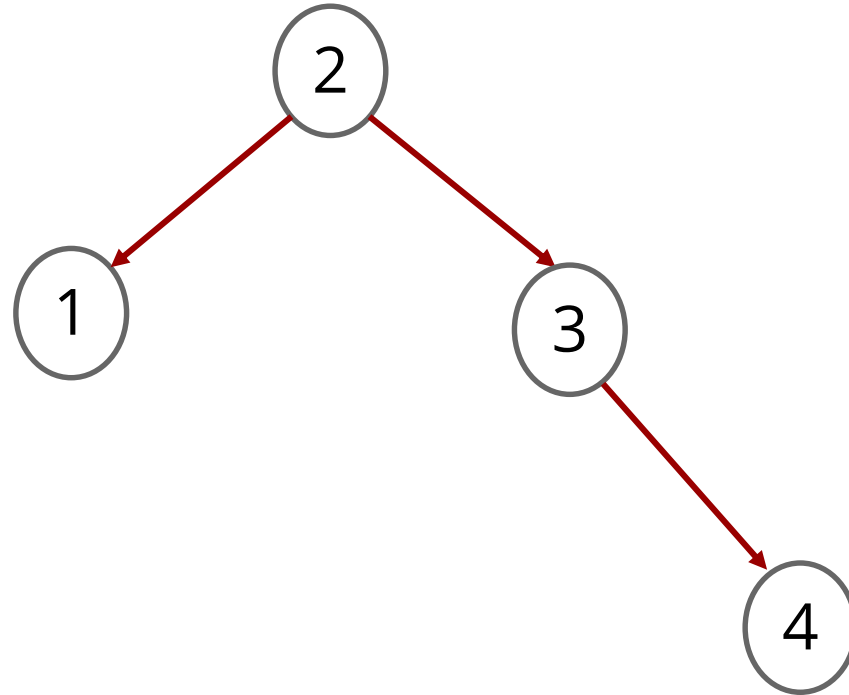
# Example

- Insert( **4** )



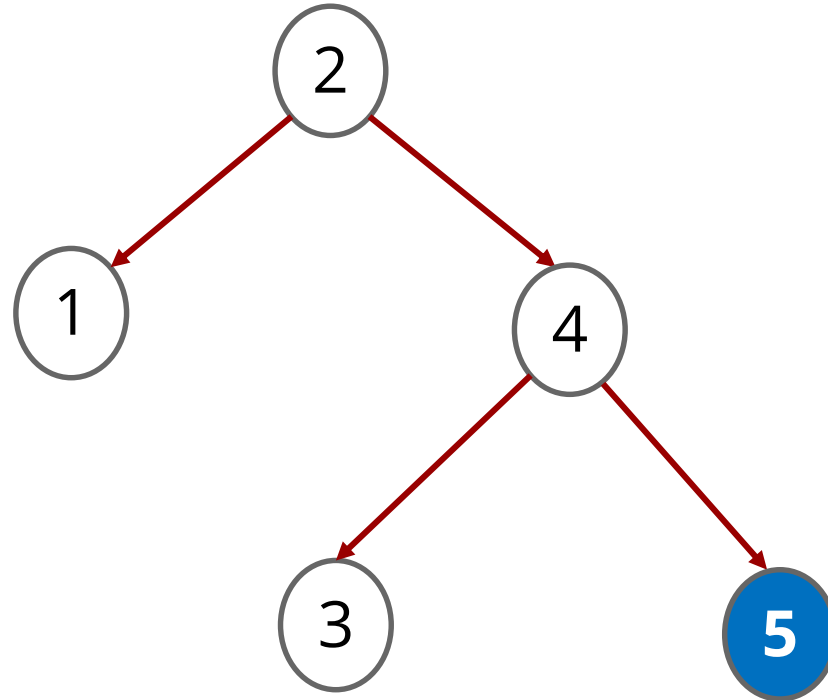
# Example

- Insert( **5** )



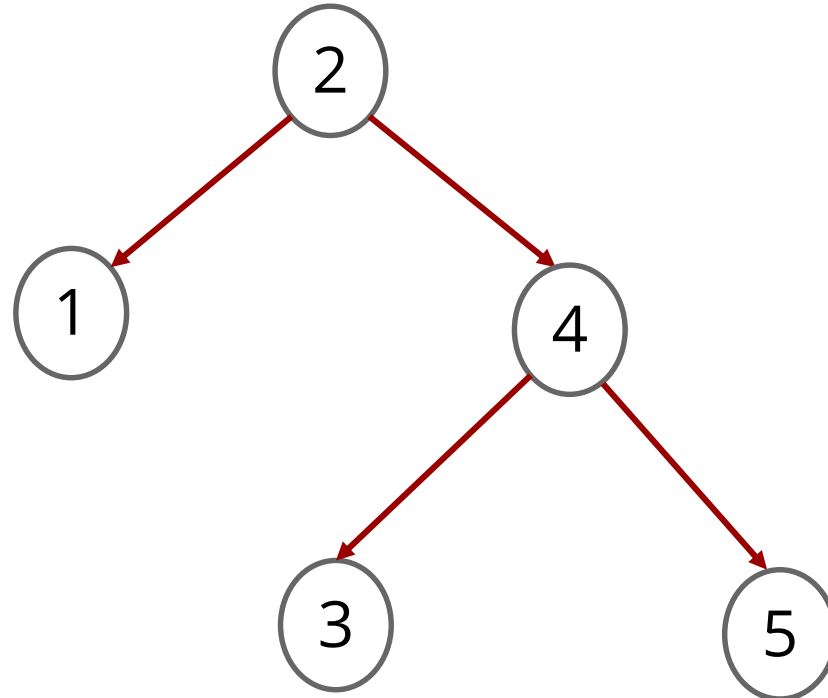
# Example

- Insert( **5** )



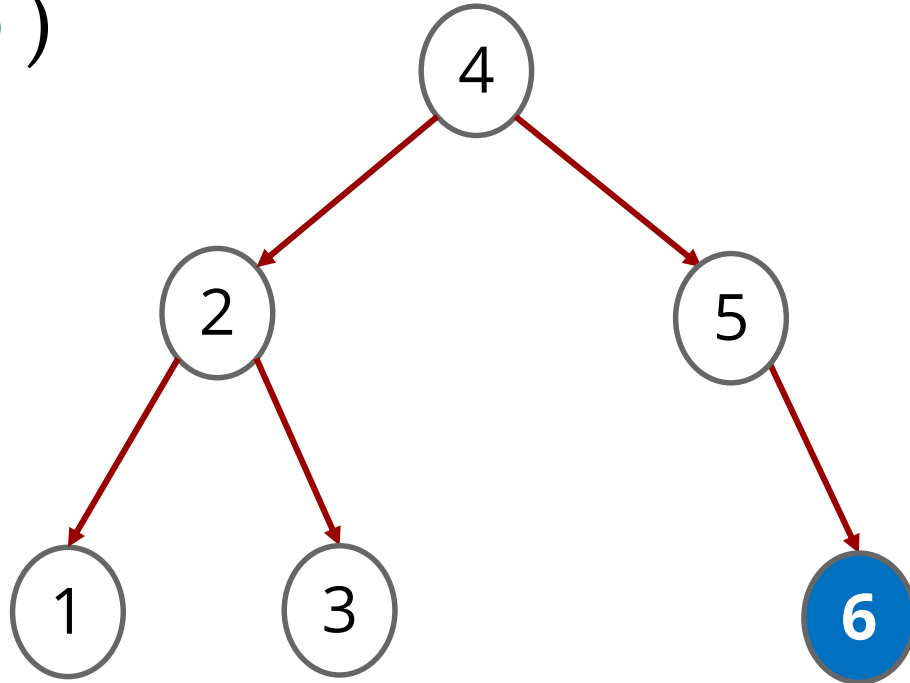
# Example

- Insert( **6** )



# Example

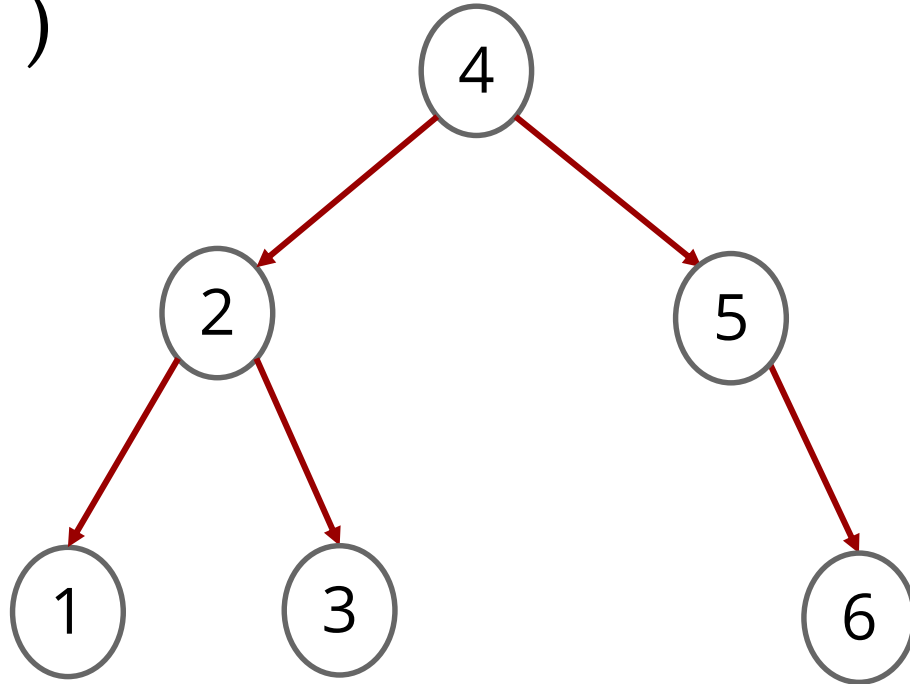
- Insert( **6** )





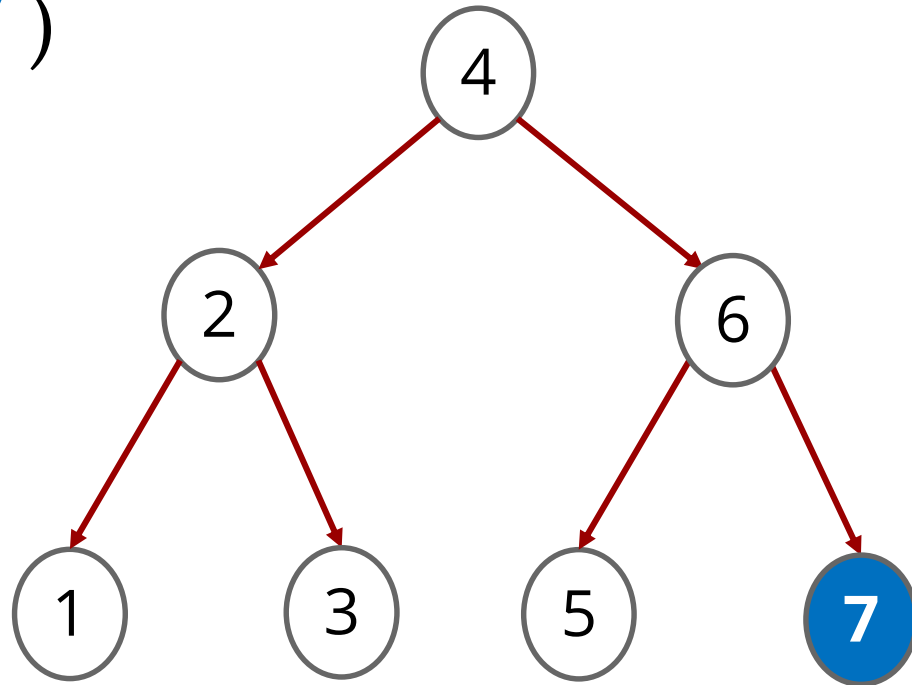
# Example

- Insert( **7** )



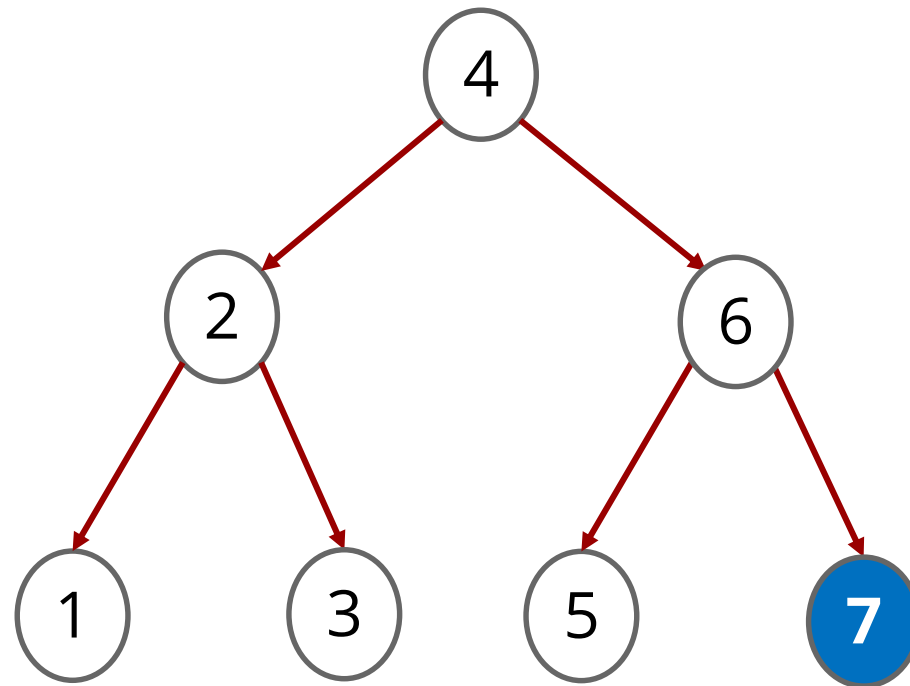
# Example

- Insert( **7** )



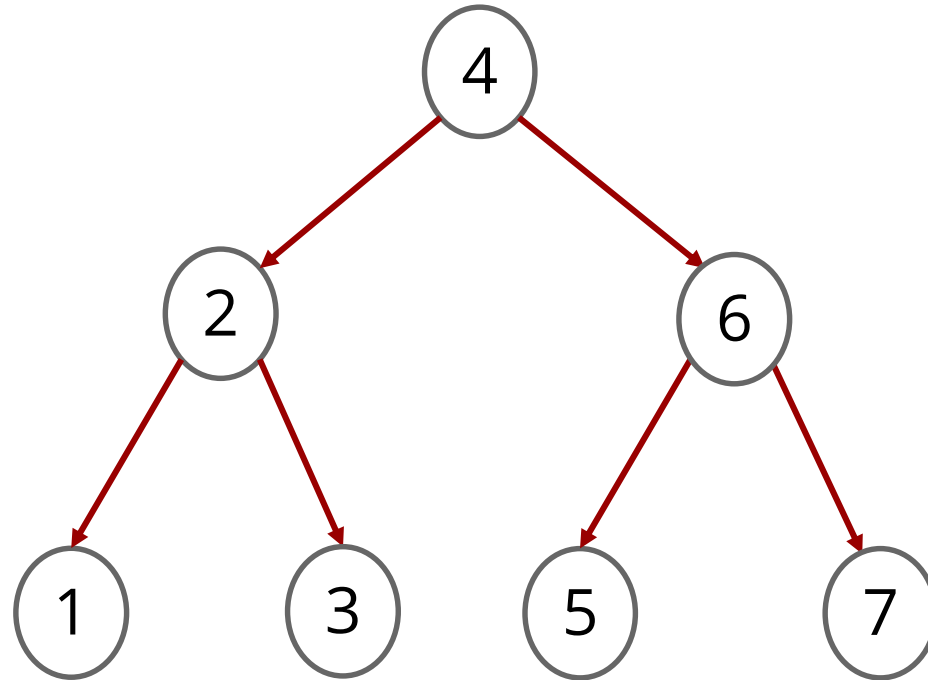
# Example

- Now insert the numbers 10 – 16 backward to the tree



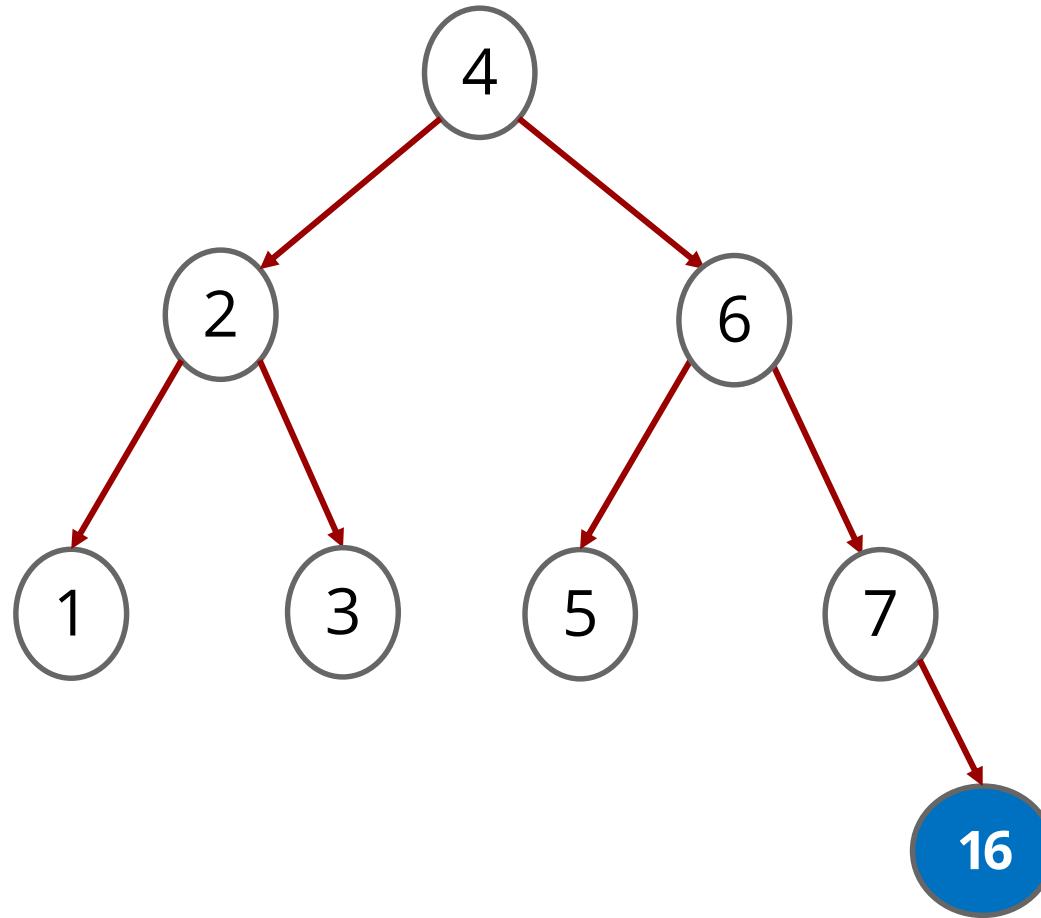
# Example

- Insert( **16** )



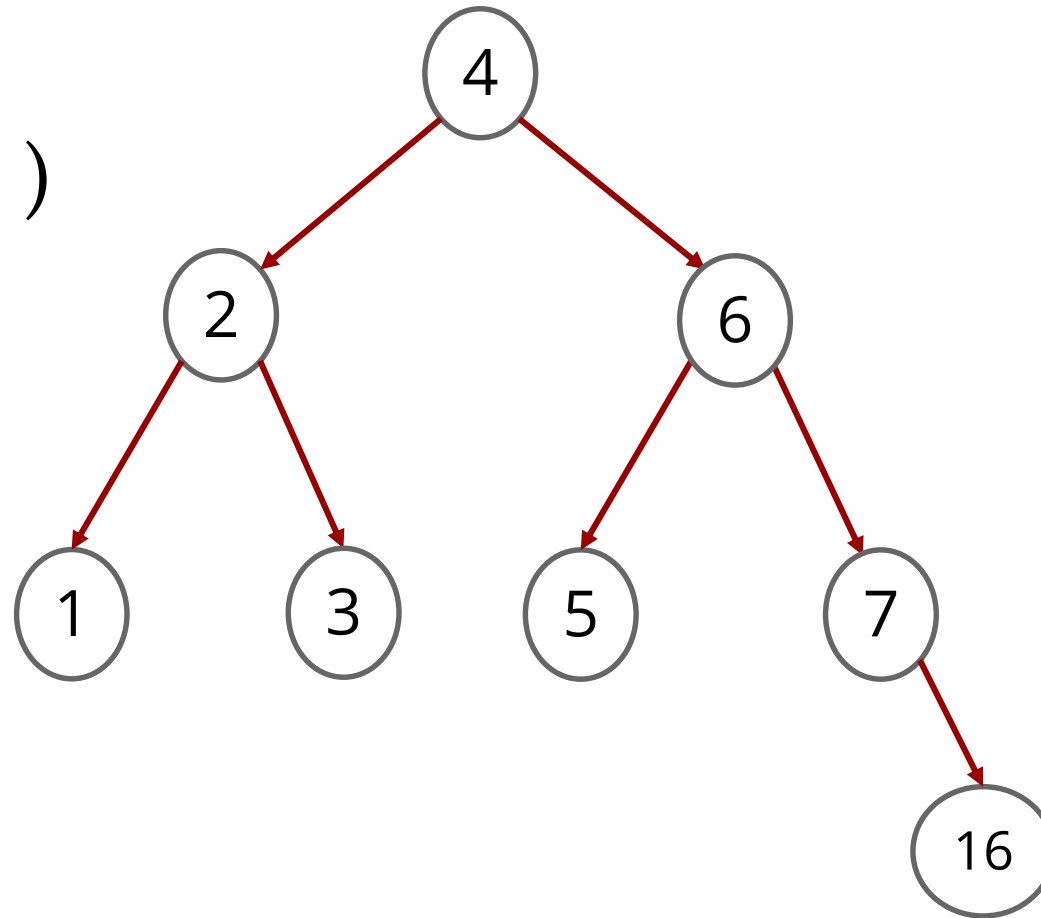
# Example

- Insert( **16** )



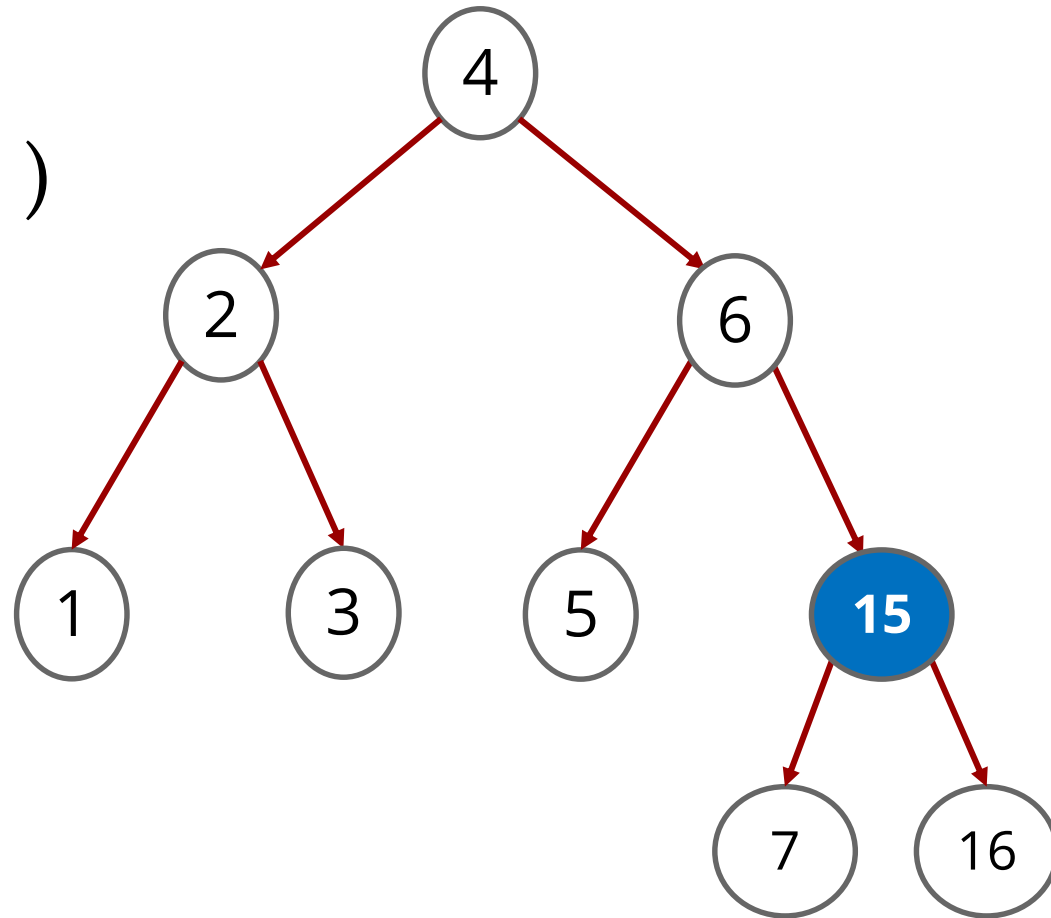
# Example

- Insert( **15** )



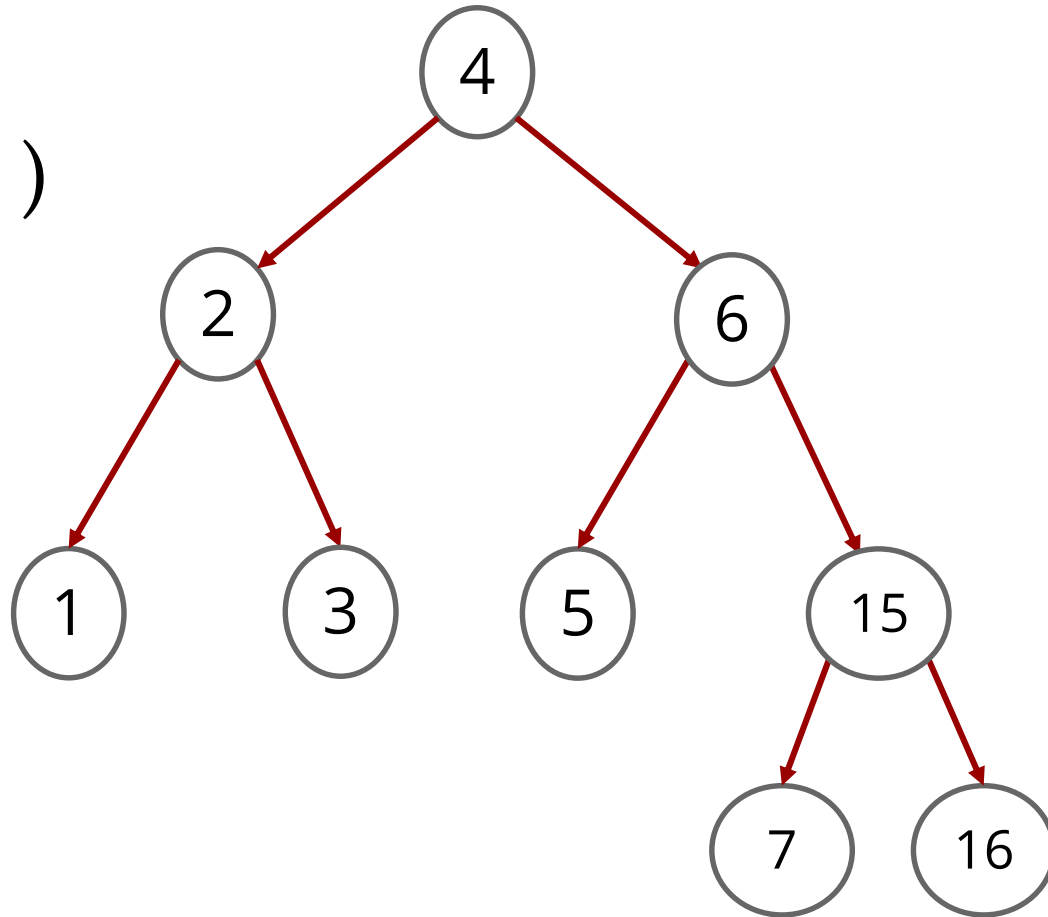
# Example

- Insert( **15** )



# Example

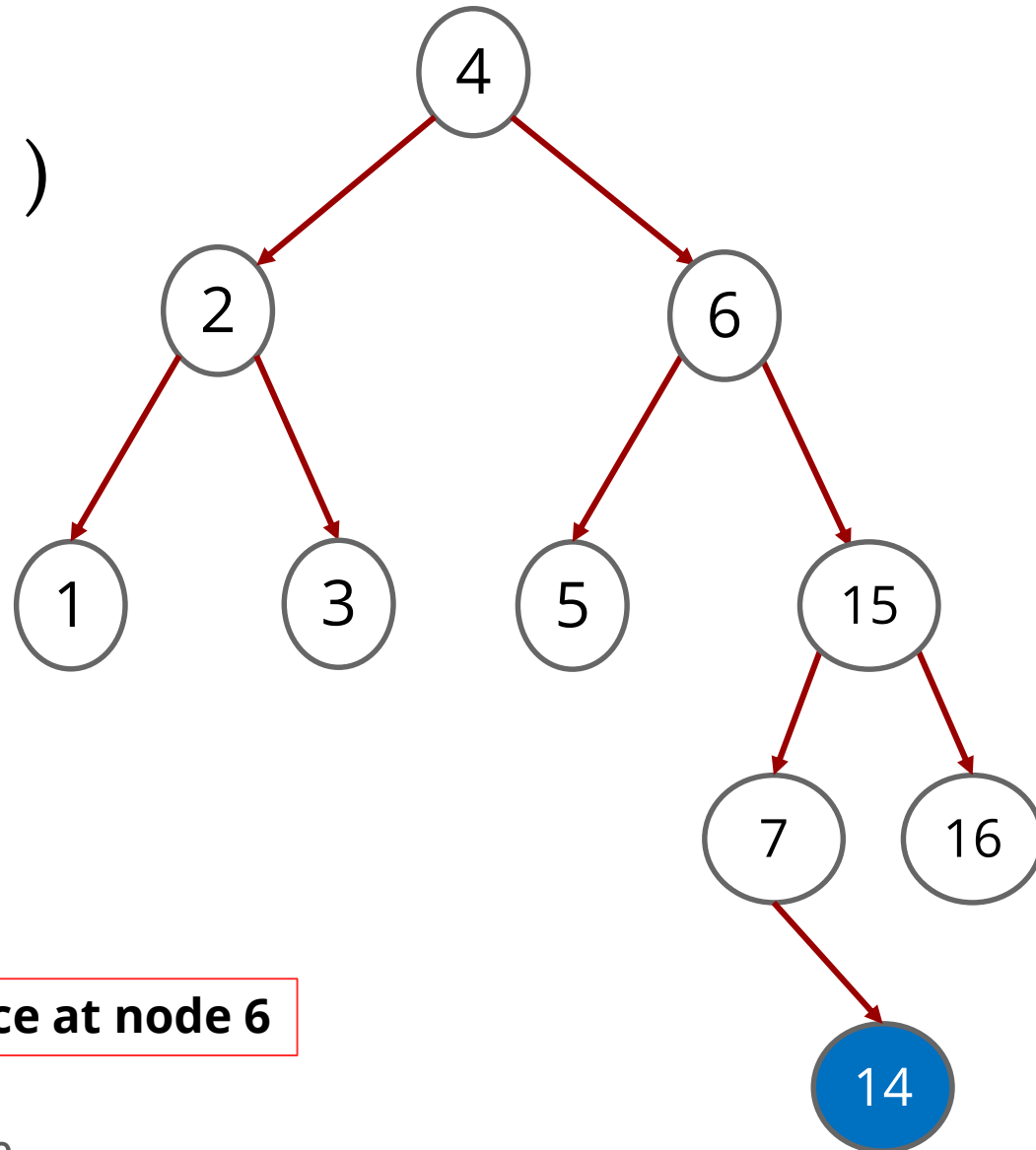
- Insert( **14** )





# Example

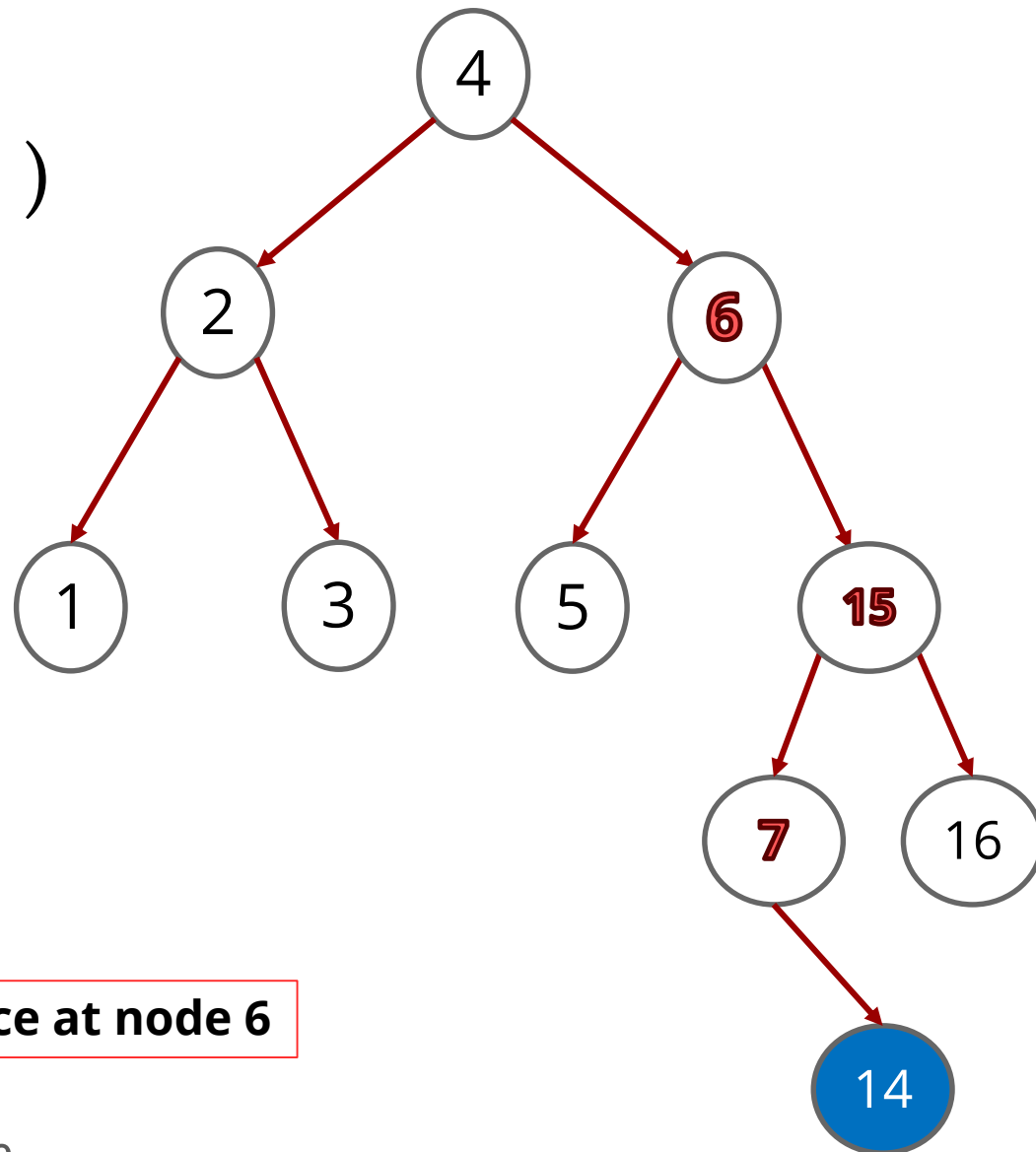
- Insert( **14** )



**14 causes imbalance at node 6**

# Example

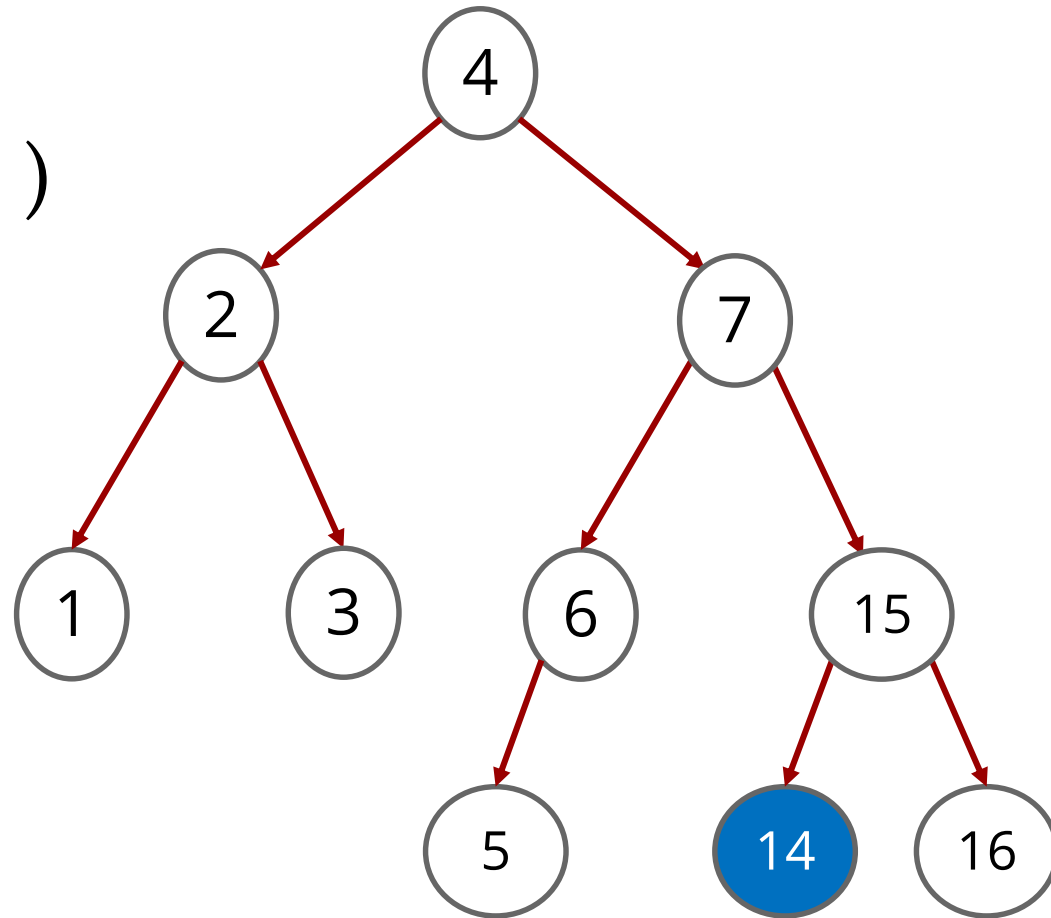
- Insert( **14** )



**14 causes imbalance at node 6**

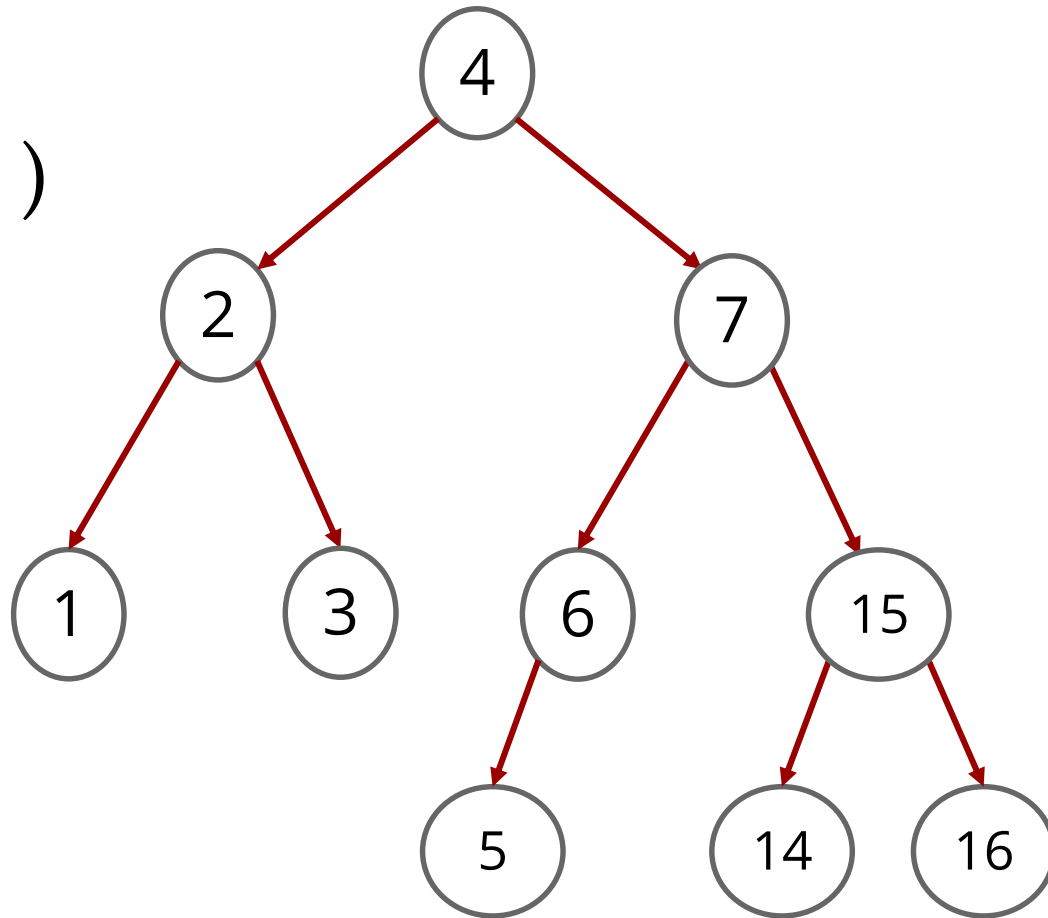
# Example

- Insert( **14** )



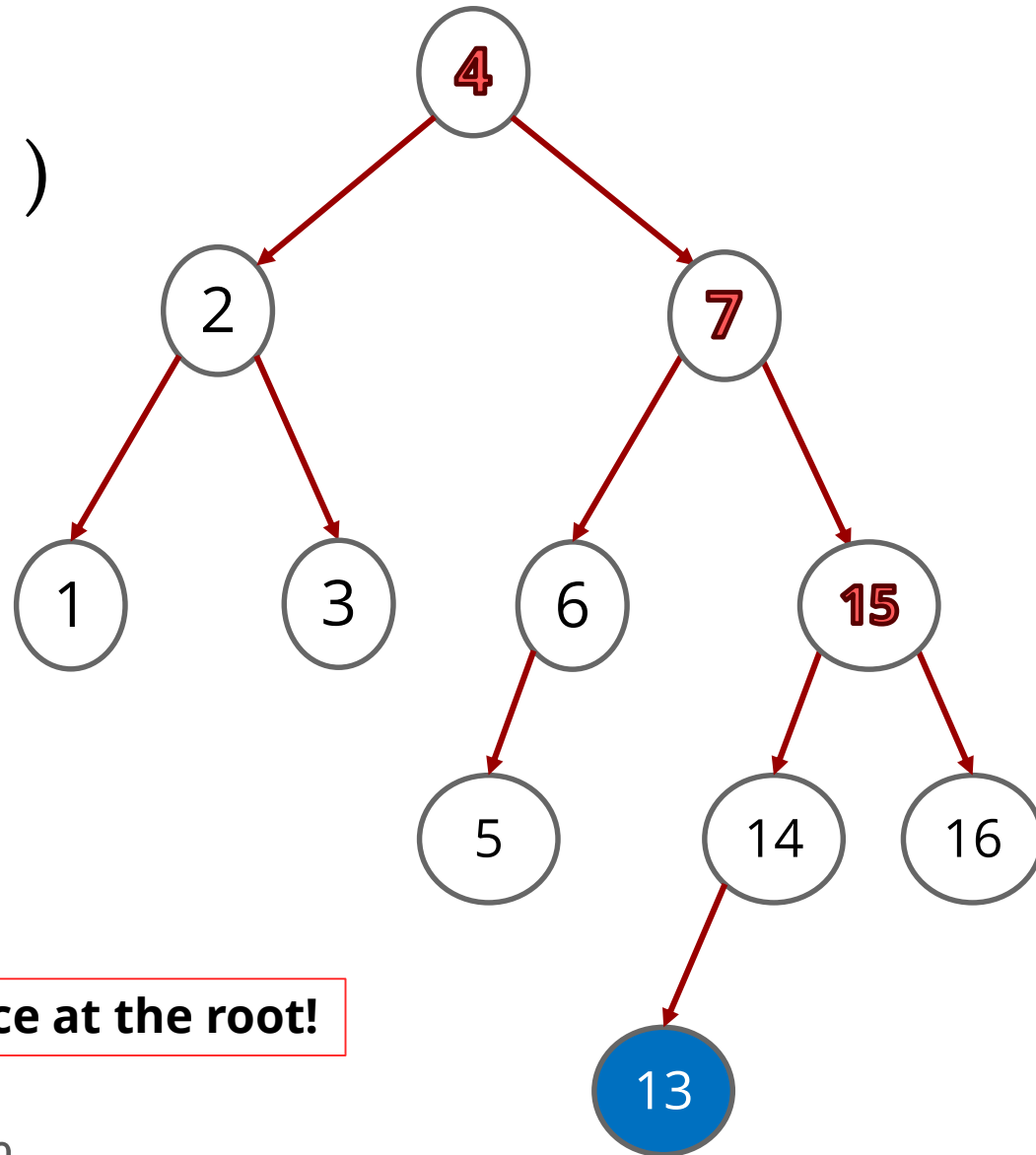
# Example

- Insert( **13** )



# Example

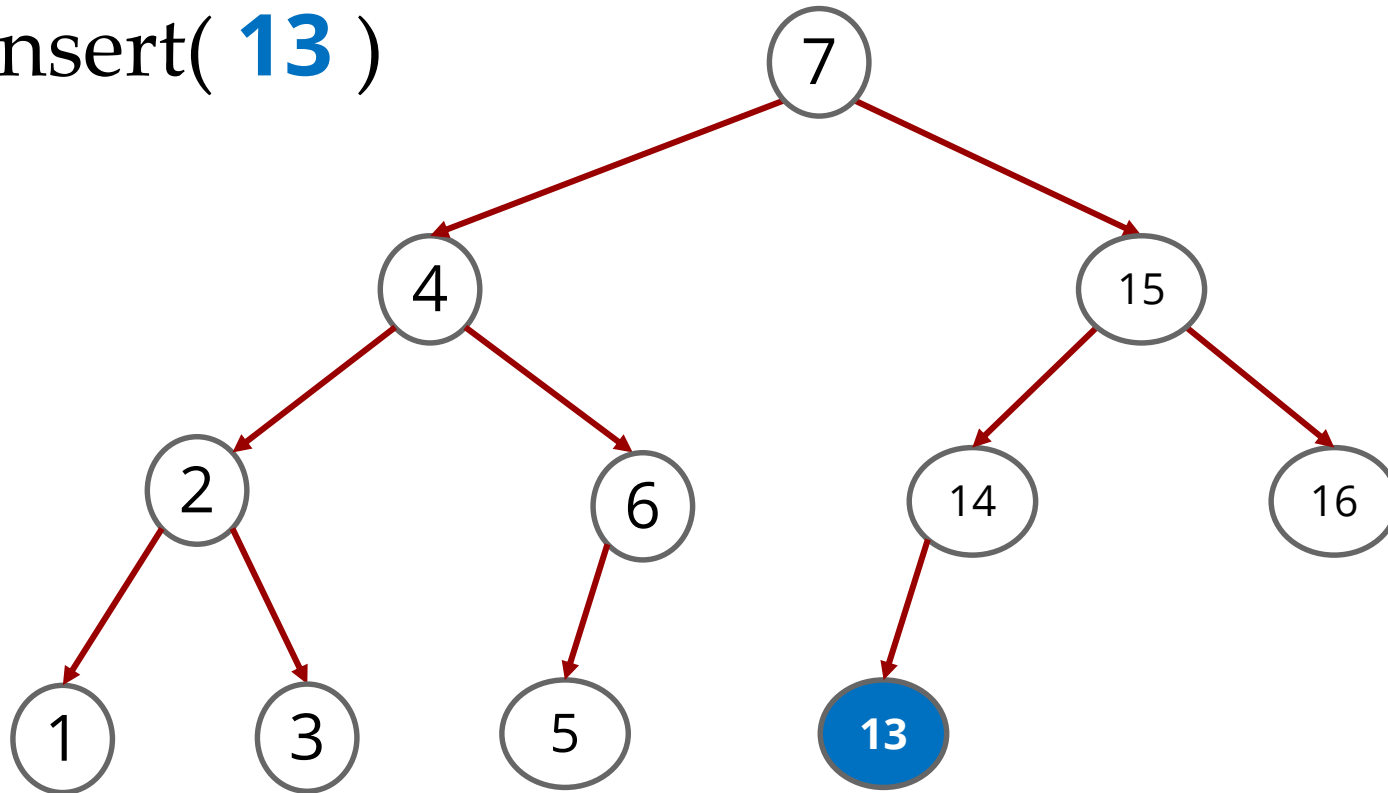
- Insert( **13** )



**13 causes imbalance at the root!**

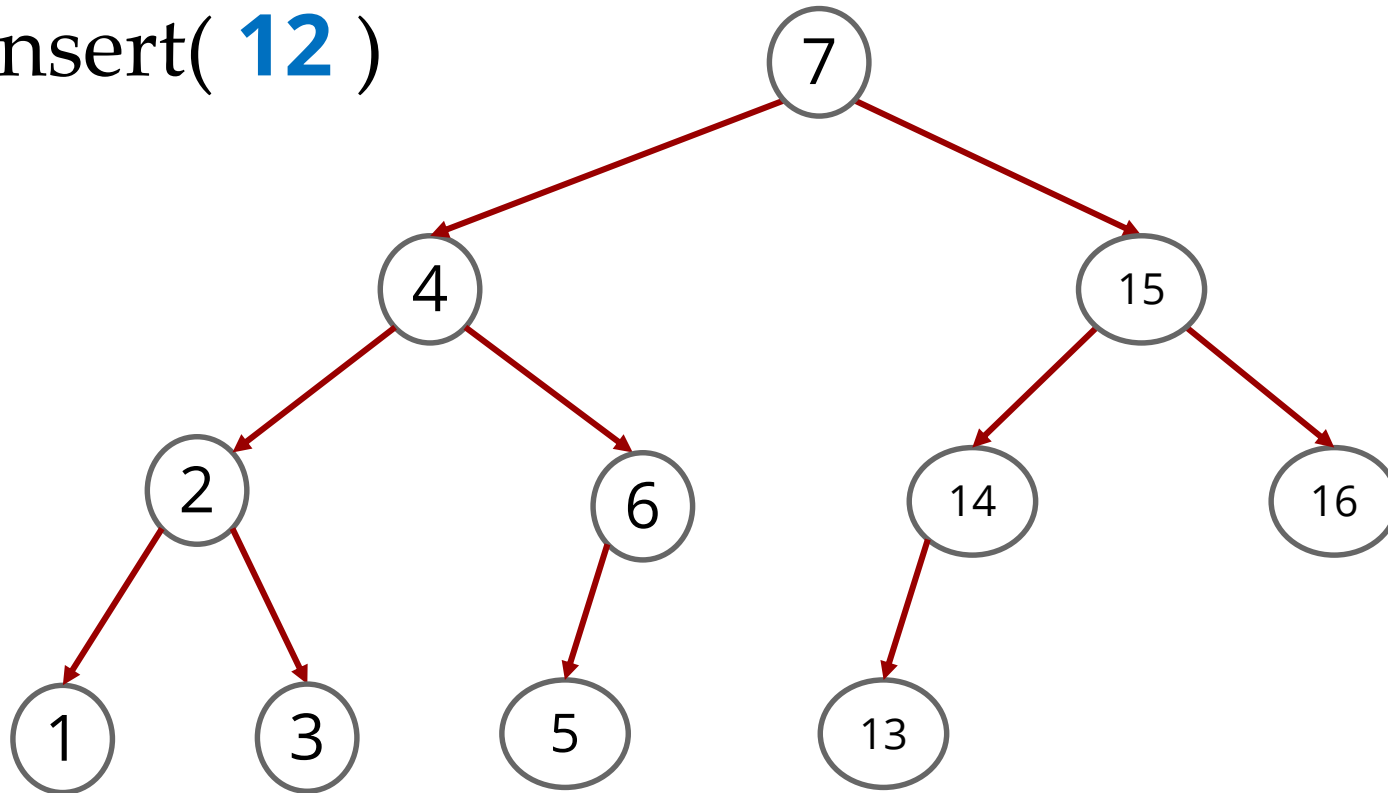
# Example

- Insert( **13** )



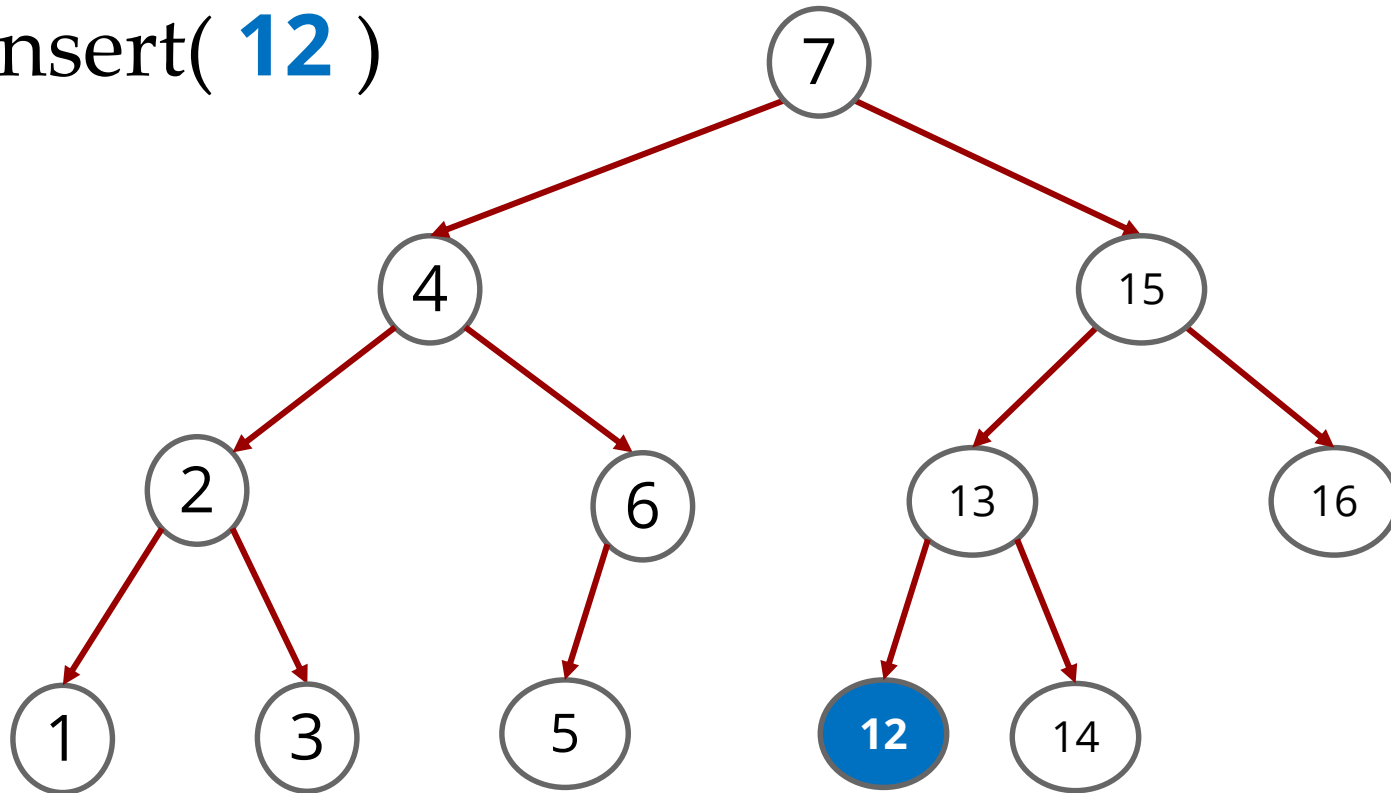
# Example

- Insert( **12** )



# Example

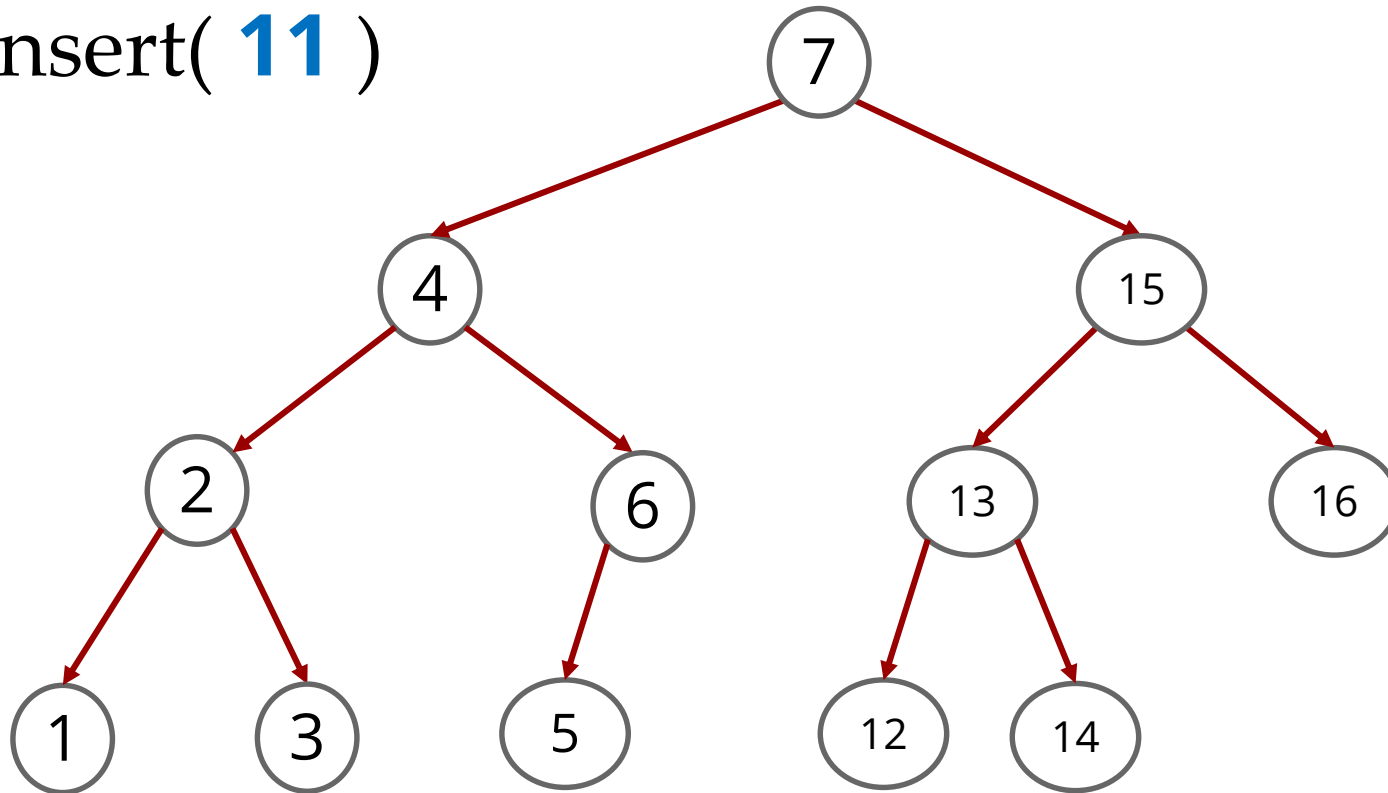
- Insert( **12** )





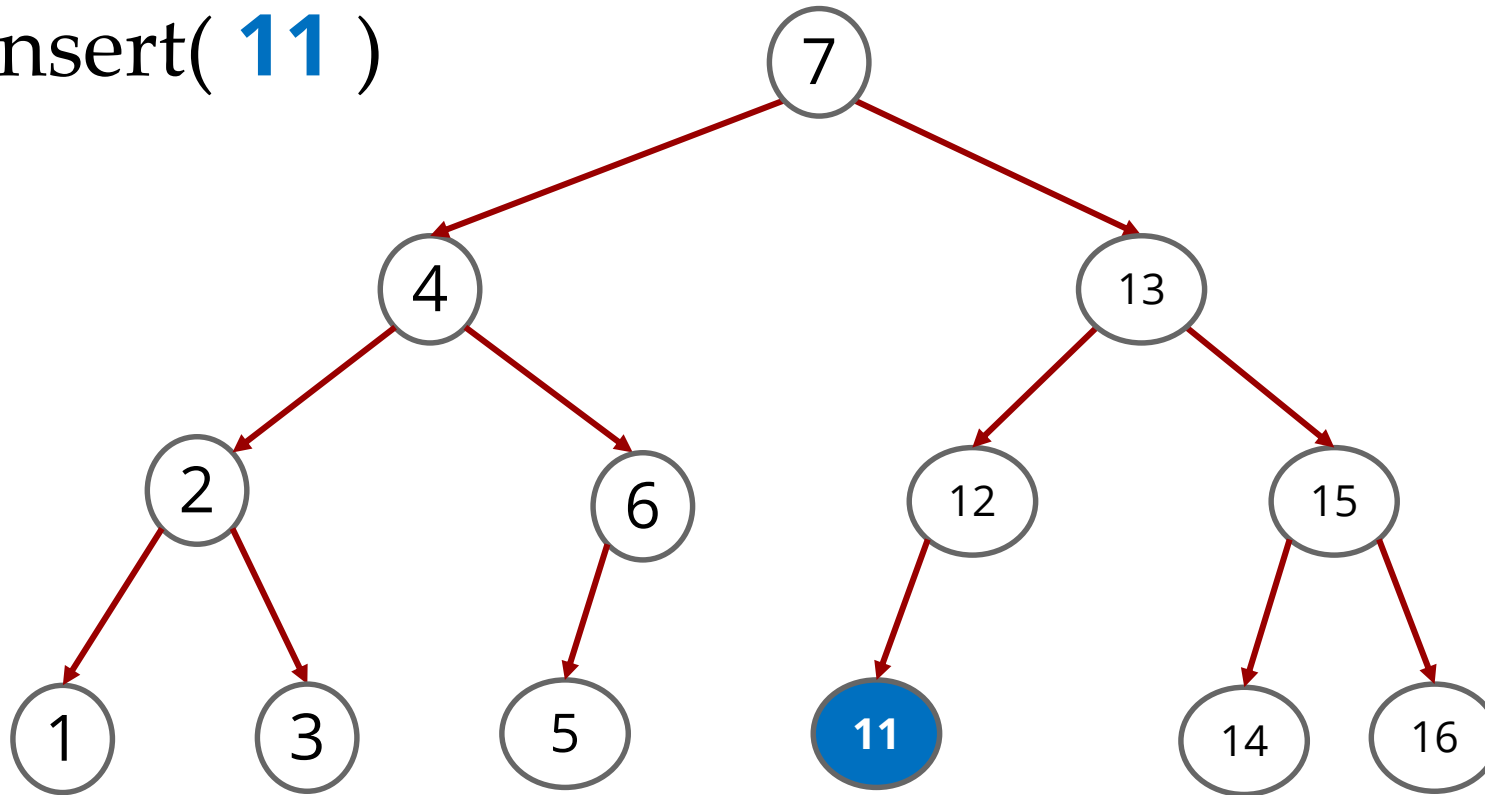
# Example

- Insert( **11** )



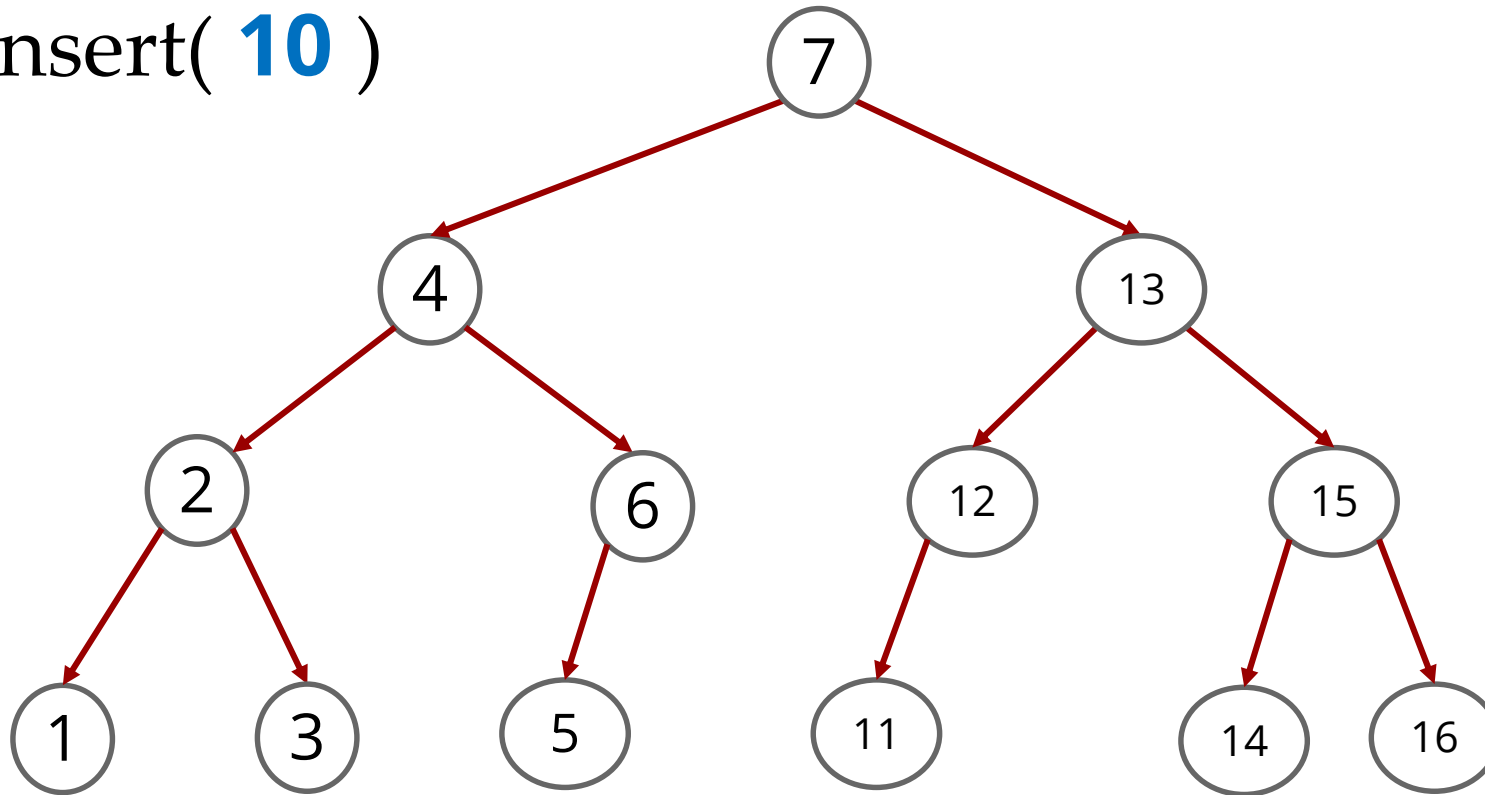
# Example

- Insert( **11** )



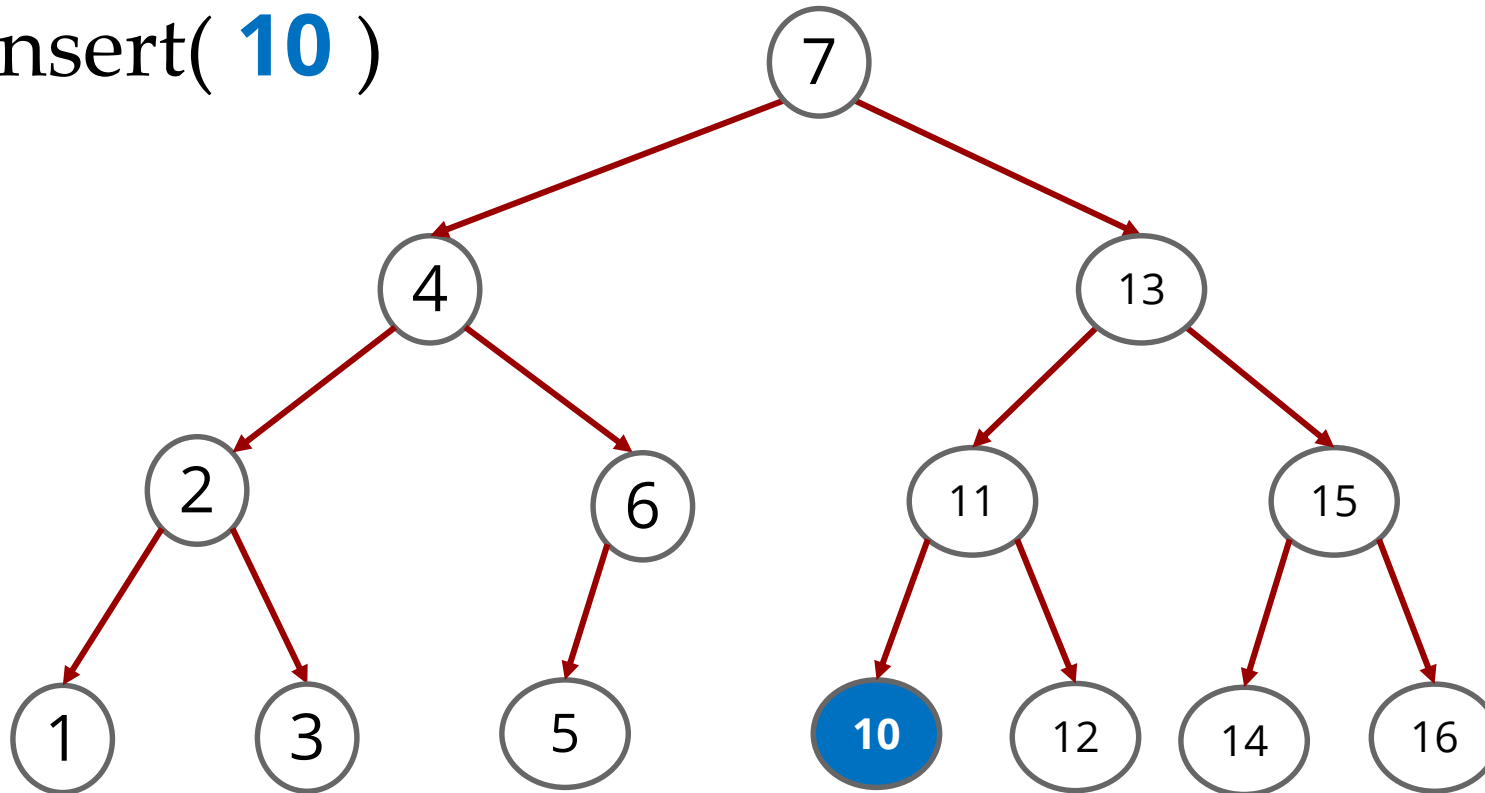
# Example

- Insert( **10** )



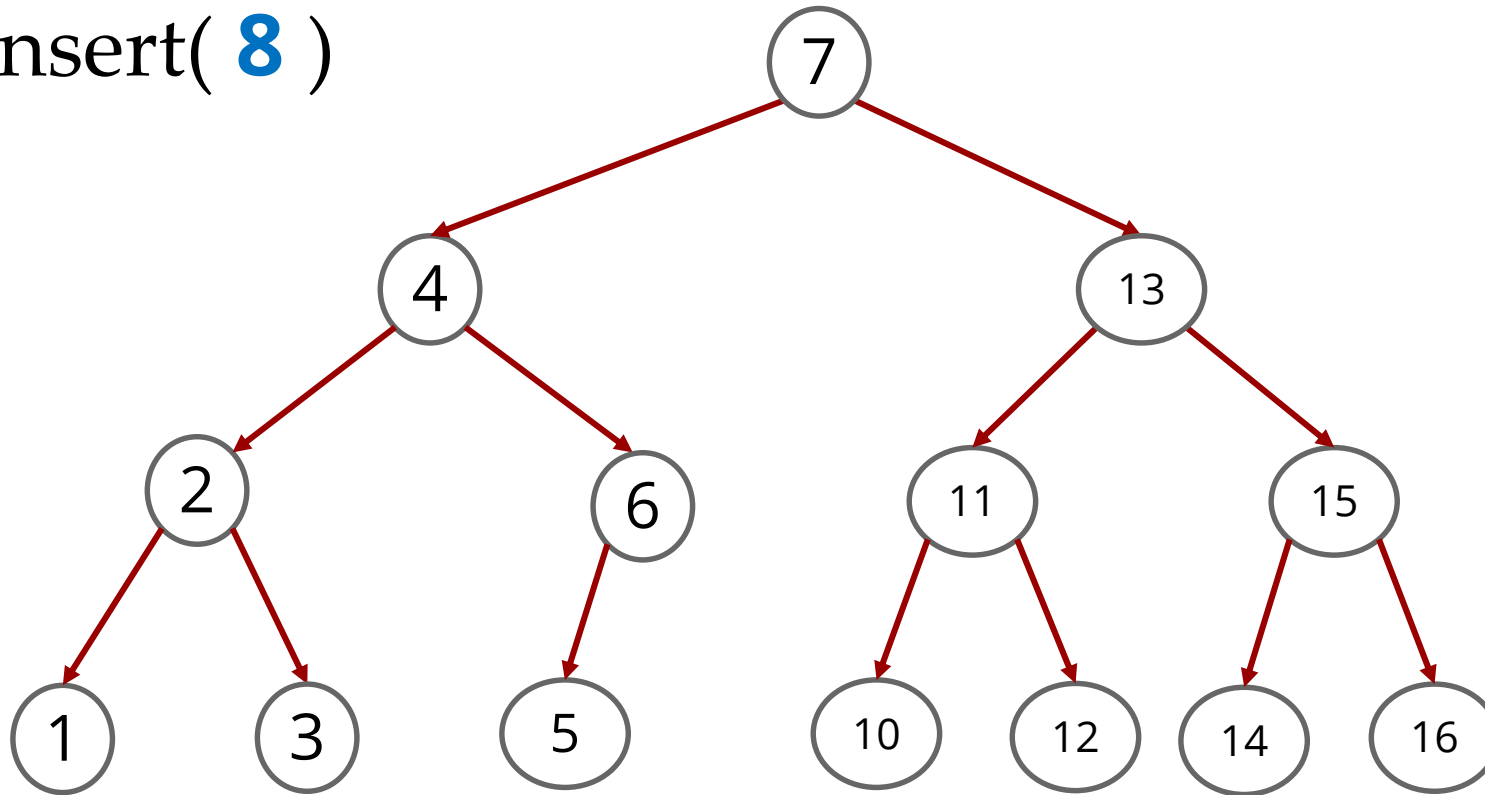
# Example

- Insert( **10** )



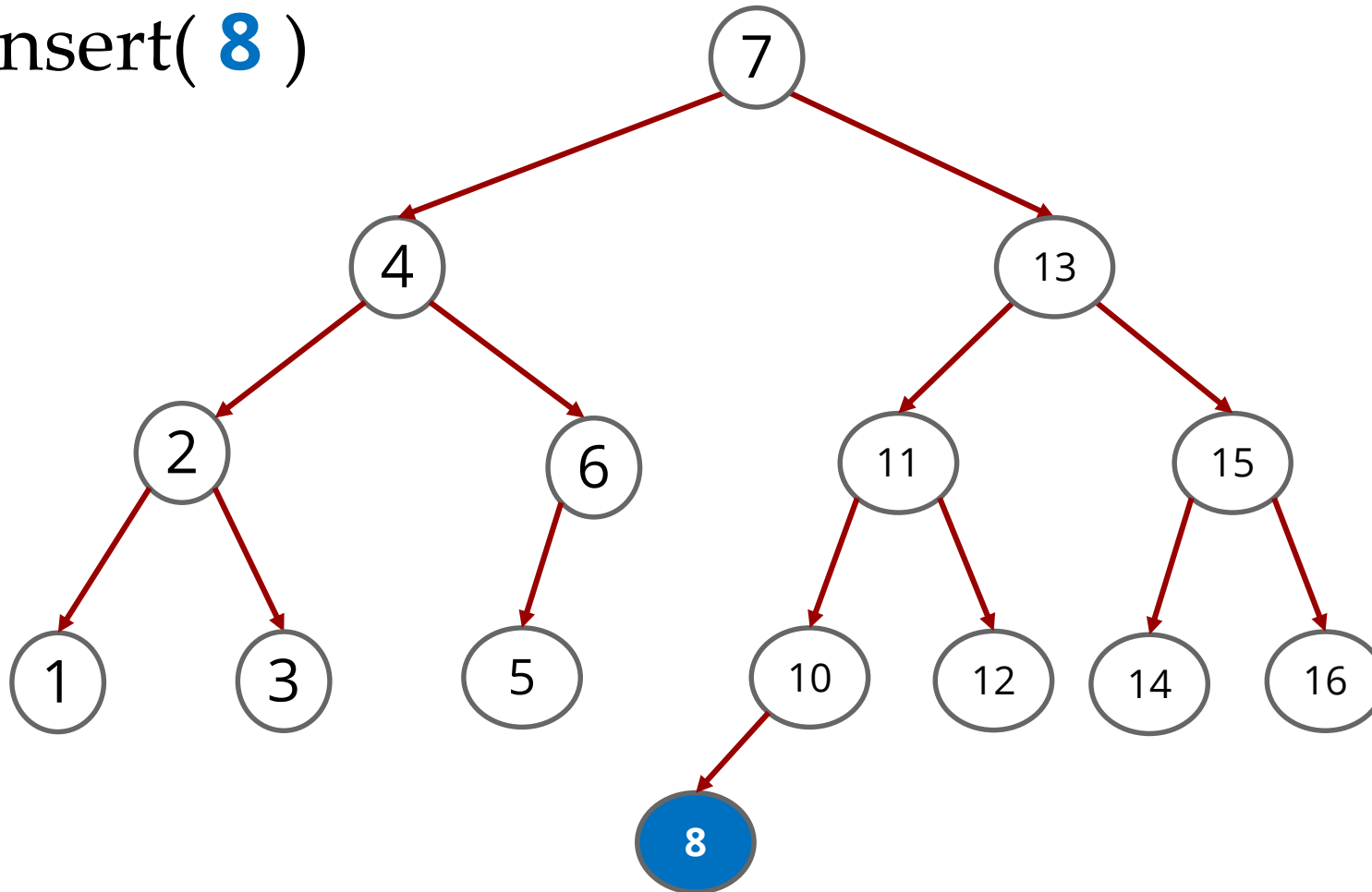
# Example

- Insert( 8 )



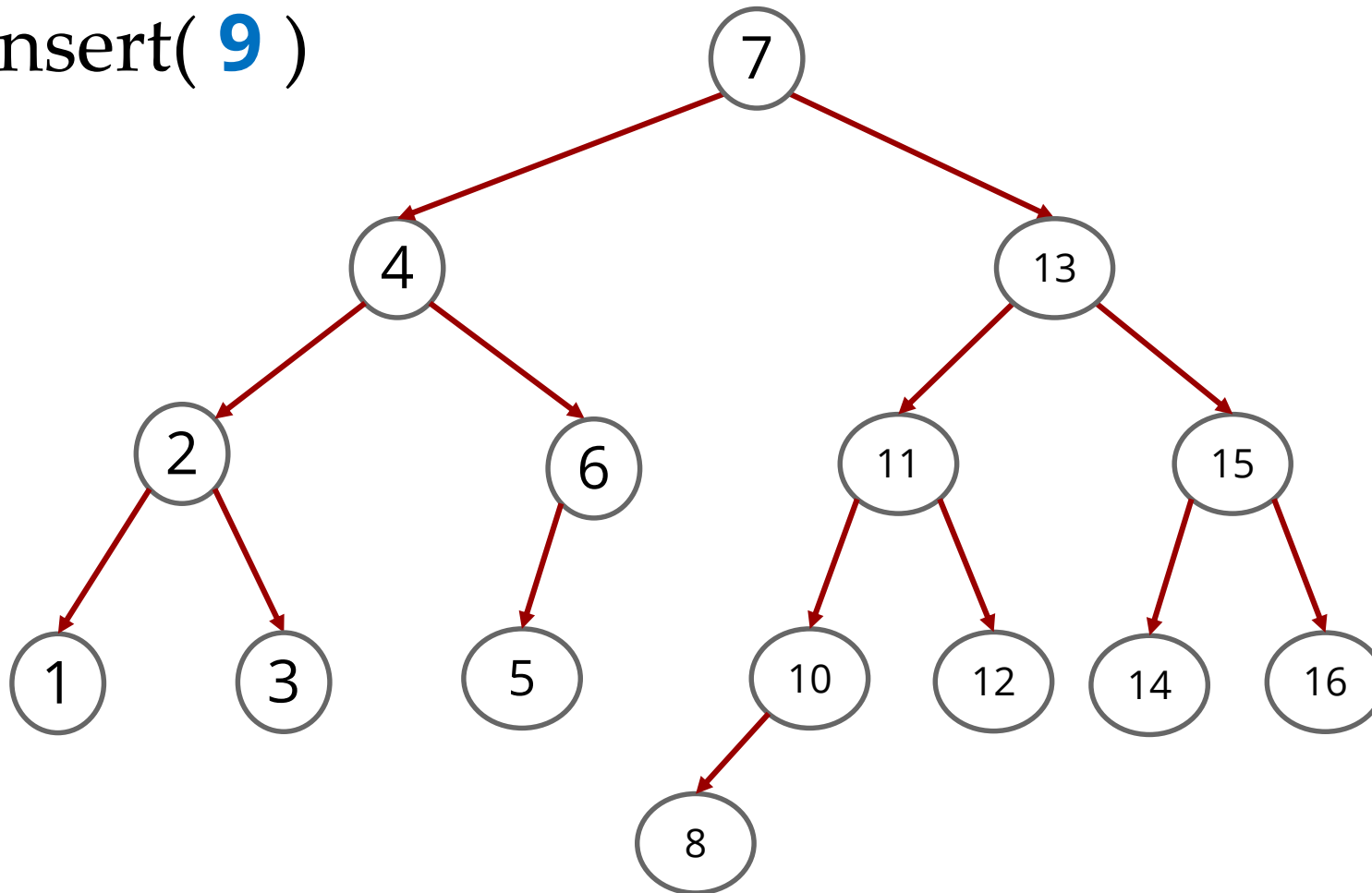
# Example

- Insert( 8 )



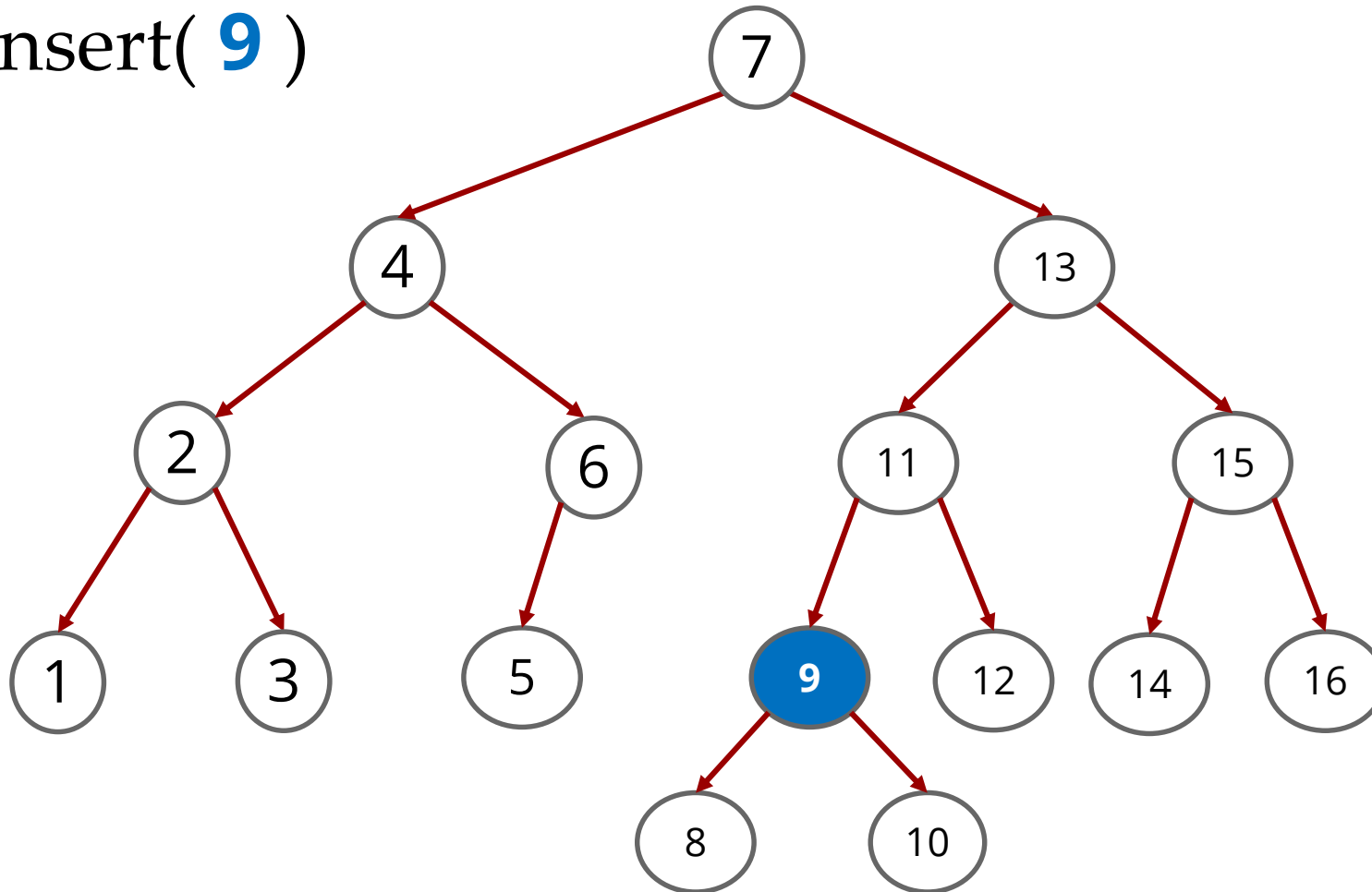
# Example

- Insert( 9 )



# Example

- Insert( **9** )





# AVL Tree Implementation

- The main structure is the same as the BST
- Differs in adding the Rotations methods
  - `SingleRotateWithLeft`
  - `SingleRotateWithRight`
  - `DoubleRotateWithRight`
  - `DoubleRotateWithLeft`
- Code will be sent (the textbook author's implementation)

# Applications of AVL Trees

- Used frequently for quick searching as it takes  $O(\log n)$  because the tree is balanced
- Used in Databases
  - Intensive look-up applications where insertion & deletion are not that frequent but search for items is performed frequently
  - The large cost of rebalancing is a limitation!
- Used for in-memory collections such as sets and dictionaries