CHAPTER 3

Huffman Coding

Huffman code is a particular type of optimal prefix code that is commonly used for lossless data compression. The process of finding or using such a code proceeds by means of Huffman coding.

ASCII code → fixed length code, it take a lot of space.

ABCDEF

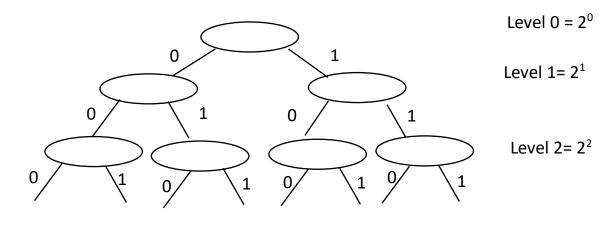
- 6 characters → each character need 8 bits
 1000 character → 8000 bits
- But we need 3 bit to represent 6 character
 1000 character → 3000 bit
- Some character occurs more frequency than the other.

Leaves→ given character set (ASCII = 256 Characters)

To code a character trace a path from root to leaves

Left child → 0

Right child → 1



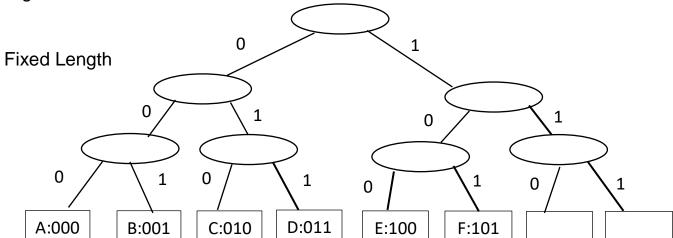
.....

Leaves→ given character set (6 characters = 3 bits)

To code a character trace a path from root to leaves

Left child → 0

Right child → 1



freq/1000

If we have file with 100,000 character

ABCDEF

Frequency 45 13 12 16 9 5

Type	Α	В	С	D	Е	F	Size
ASCII	1000001	1000010	1000011	1000100	1000101	1000110	800,000 bit = 100,000 byte = 100 Kbyte
Fixed Length	000	001	010	011	100	101	300,000 bit = 37,500 byte = 36.6 Kbyte
Variable Length	0	101	100	111	1101	1100	1 * 45,000 + 3 * 13,000 + 3 * 12,000 + 3 * 16,000 + 4 * 9000 + 4 * 5000 = 244,000 bit = 28,000 byte = 27.3 Kbyte

Example:

ACDFB

01001111100101

01001111100101

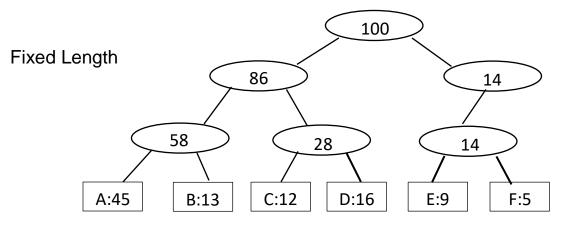
ACDF B

Leaves→ given character set

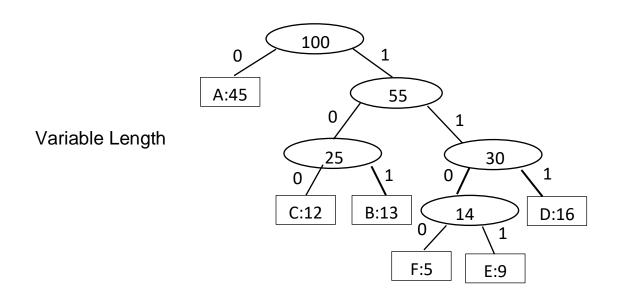
To code a character trace a path from root to leaves

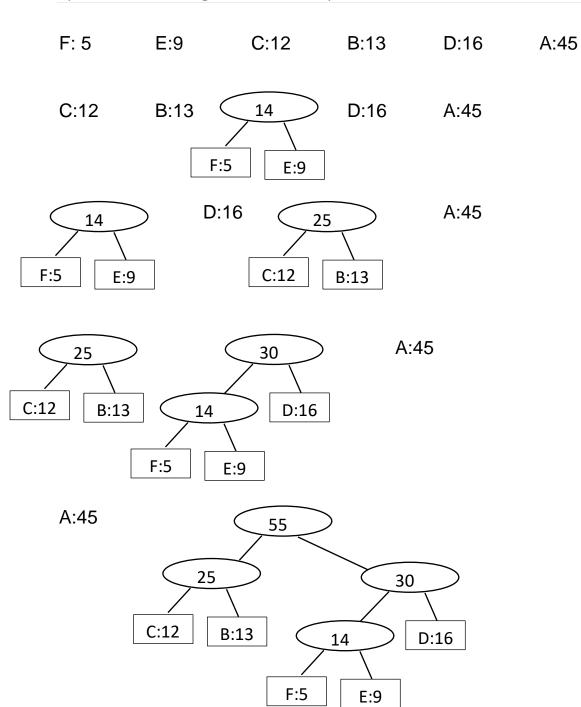
Left child → 0

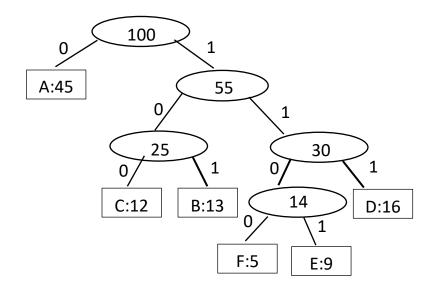
Right child → 1



Not necessary to be full tree







Algorithm

```
n = number of leaves
Q → Priority Queue
Insert ( Q, Character)
for ( i = 1; i < n; i++)
    z = allocate new node;
    x = deleteMin( Q );
    y = deleteMin( Q );
    z.leftChild ← x
    z.rightChild ← y
    freq( z ) = freq( x ) + freq( y );
    insert ( Q, z)
end for</pre>
```