



**Birzeit University**  
**Faculty of Engineering and Technology**  
**Department of Electrical and Computer Engineering**  
**First Semester – 2023/2024**  
**ENCS2340 - Digital Systems**  
**Homework # 2**

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**Section : 1**

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**Notes:**

- 1- Use this page as a cover for your homework.
- 2- Late homeworks will not be accepted (the system will not allow it).
- 3- Due date is Wednesday January 17, 2024 at 11:59 pm on ritaj.
- 4- Organize your solution for each question (Q1, Q2, etc.) and add them to one file. Then, name your file as (Assign2\_LastName\_FirstName\_StudentID.pdf).

**Q1 (10 points):** Design a combinational circuit with three inputs,  $x$ ,  $y$  and  $z$ , and the three outputs,  $A$ ,  $B$ , and  $C$ . when the binary input is 0, 1, 2, or 3, the binary output is one greater than the input. When the binary input is 4, 5, 6, or 7, the binary output is one less than the input.

① specification :-

- 1) There is 3-inputs ( $x, y, z$ ) and 3-outputs ( $A, B, C$ )
- 2) if the Binary input is  $[0, 1, 2, 3]$ , then the Binary output is one greater than the input.
- 3) if the Binary input is  $[4, 5, 6, 7]$ , then the Binary output is one less than the input.

② Formulation - By truth table

x	y	z	A	B	C
0	0	0	0	0	1
0	0	1	0	1	0
0	1	0	0	1	1
0	1	1	1	0	0
1	0	0	0	1	1
1	0	1	1	0	0
1	1	0	1	0	1
1	1	1	1	1	0

→  $A = \sum(3, 5, 6, 7)$

$B = \sum(1, 2, 4, 7)$

$C = \sum(0, 2, 4, 6)$

③ Logic Minimization :-

k-map for A

k-map for B

x \ yz	00	01	11	10
0			1	
1		1	1	1

x \ yz	00	01	11	10
0		1		1
1	1		1	

$A = xz + xy + yz$

$B = x'y'z + x'y'z' + xy'z' + xyz$   
 $= x'(y'z + y'z') + x(y'z' + yz)$   
 $= x'(y \oplus z) + x(y \oplus z)'$   
 $= x \oplus y \oplus z$

k-map for C

x \ yz	00	01	11	10
0	1			1
1	1			1

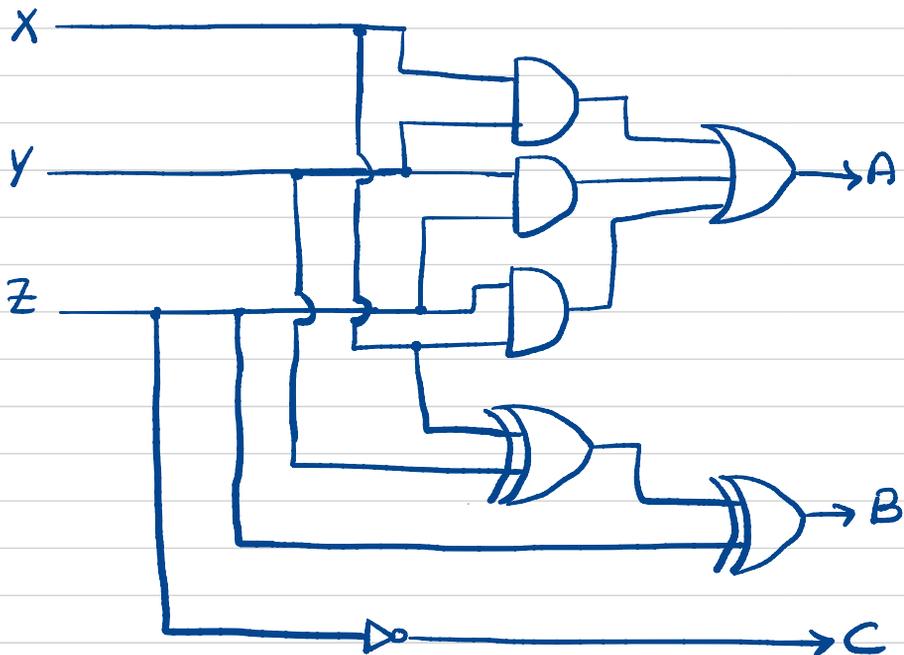
$C = z'$

#### ④ Technology Mapping :-

$$A = XY + XZ + YZ$$

$$B = X \oplus Y \oplus Z$$

$$C = Z'$$



Q2 (5 points): Implement the Boolean function  $F(A,B,C) = AB + A'C + A'B'$  Using a single 4x1 multiplexer.

$$F(A,B,C) = AB + A'C + A'B'$$

truth table:

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

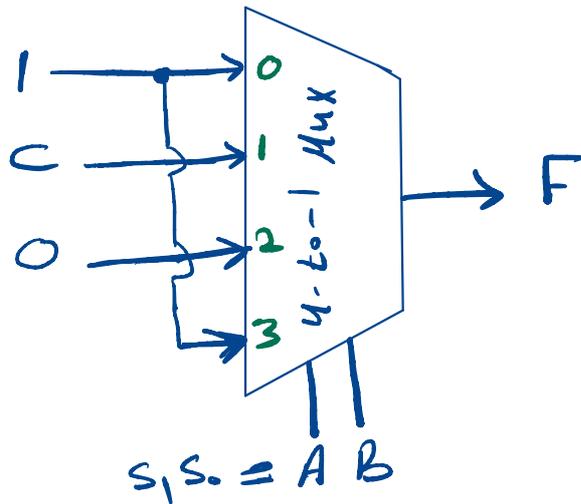
$F=1$

$F=C$

$F=0$

$F=1$

→ For (4x1) multiplexer we need 2-select then A is  $s_1$ , B is  $s_0$  and C is input



Q3 (5 points): Implement the same function in Q2 using the minimum number of 2x4 decoders with enable and a single NOR gate.

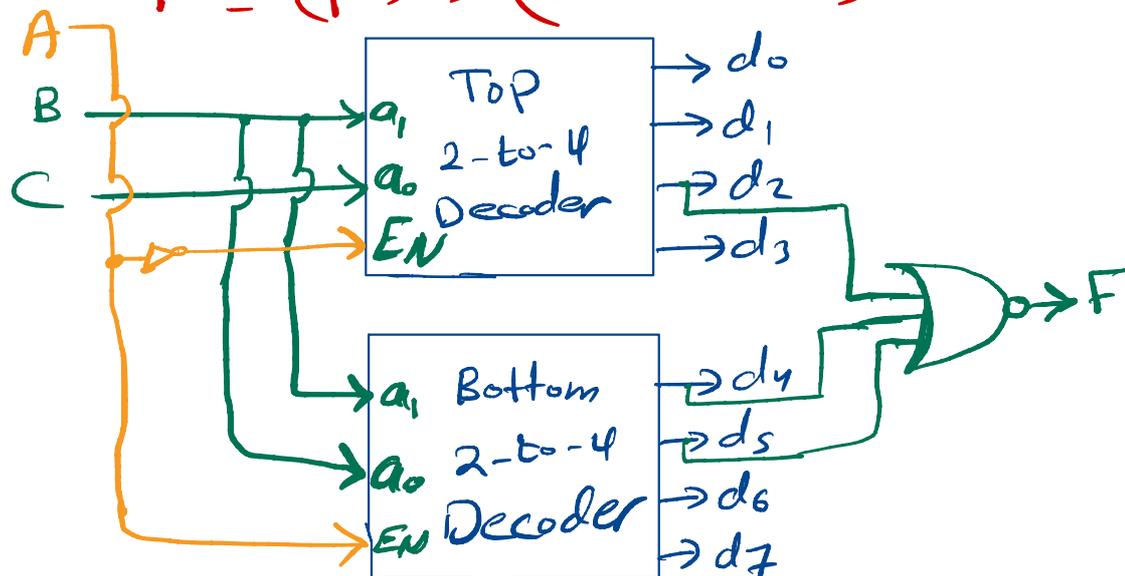
$$F(A,B,C) = AB + A'C + A'B' = \sum(0, 1, 3, 6, 7)$$

truth table:

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

$$F' = \sum(2, 4, 5) = d_2 + d_4 + d_5$$

$$F = (F')' = (d_2 + d_4 + d_5)'$$



Q4 (10 points): Implement the following function  $F(A,B,C,D) = \sum(0, 2, 4, 6, 8, 10)$  using

- Mux 4x1
- Decoders 3-to-8
- AND-OR
- NAND-NAND

$$F(A,B,C,D) = \sum(0, 2, 4, 6, 8, 10) = d_0 + d_2 + d_4 + d_6 + d_8 + d_{10}$$

Truth table:

A	B	C	D	F
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	1
0	1	1	1	0
1	0	0	0	1
1	0	0	1	0
1	0	1	0	1
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

(a) Mux 4x1

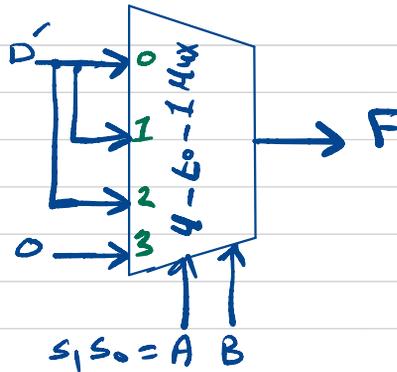
For Mux 4x1 we need 2-select then A will be  $S_1$  and B  $S_0$

$F = D'$

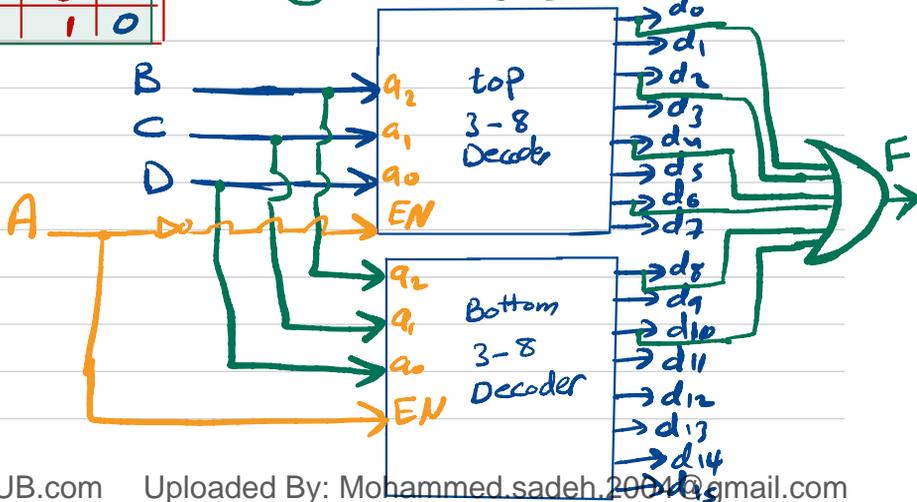
$F = D'$

$F = D'$

$F = 0$



(b) Decoders 3-to-8

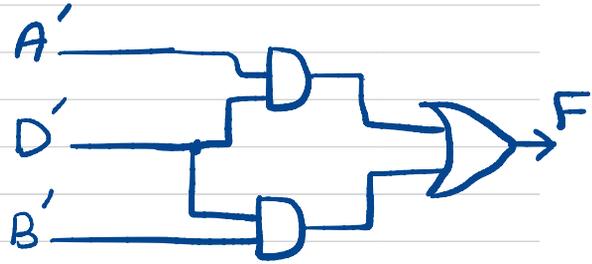


### ③ AND-OR

$$F(A, B, C, D) = \sum (0, 2, 4, 6, 8, 10)$$

AB \ CD	00	01	11	10
00	1			1
01	1			1
11				
10				1

$$F = A'D' + B'D'$$

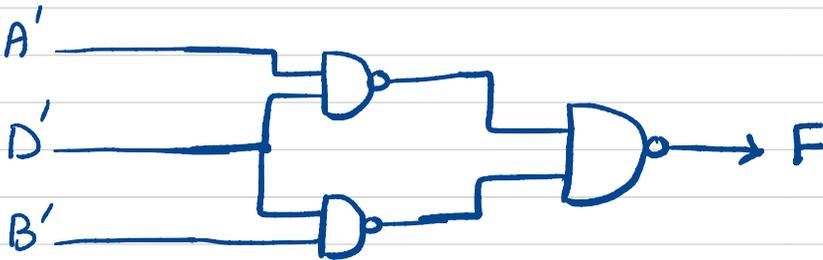


### ④ NAND-NAND

$$F = A'D' + B'D'$$

$$F' = (A'D' + B'D')' = (A'D')' \cdot (B'D')'$$

$$F = (F')' = ((A'D')' \cdot (B'D')')'$$



Q5 (6 points): In the following function determine the Essential prime implicant

$$F(A,B,C,D) = \Sigma (0,2,5,7,6,8,9,10,11,13,14,15)$$

$$F(A,B,C,D) = \Sigma (0,2,5,7,6,8,9,10,11,13,14,15)$$

CD \ AB	00	01	11	10
00	1			1
01		1	1	1
11		1	1	1
10	1	1	1	1

⇒ All prim implicants :-  
 $AC, CD', AB', BD, BC, AD, B'D'$

⇒ Essential prime implicants:  
 $BD, B'D'$

Q6 (4 points): Explain the concept of odd parity generator?

Odd Parity: one of the two parities we use in generators and checkers to add a parity bit to the  $n$ -bit code in generators and then the checkers receive  $(n+1)$  bit.  
 in odd parity:

- Count of 1's in the  $(n+1)$ -bit code is odd.
- use an even function to generate the odd parity bit
- use an even function to check the  $(n+1)$ -bit code

Example: For 4-bit code (1101) use odd parity to generate parity bit and check the message after receive we use even function to determine the parity bit →  $P=0$  then the sender transmits (01101)

then the receiver check it by even function to determine if there is an error, if received as (01101) →  $E=0$

But if one bit changes from 0 to 1 or 1 to 0 →  $E=1$  and there is an error.