

# COMP2421 – DATA STRUCTURES AND ALGORITHMS

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## Splay Trees

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# Splay Tree

- Splay tree is a binary search tree.
- It has one interesting difference: whenever an element is looked up in the tree, the splay tree reorganizes to **move** that element to the root of the tree, without breaking the binary search tree invariant.
- Meaning, it brings the recently accessed item to the root of the tree.

## Splay Tree (2)

- Splay trees provide excellent locality. Frequently accessed items are easy to find. Infrequent items are far away from the root.
- The performance of the splay tree is  $O(\log n)$ . For a sequence of  $M$  operations performed on an  $n$ -node, splay tree takes  $O(M \log n)$  time.

## Splay Tree (3)

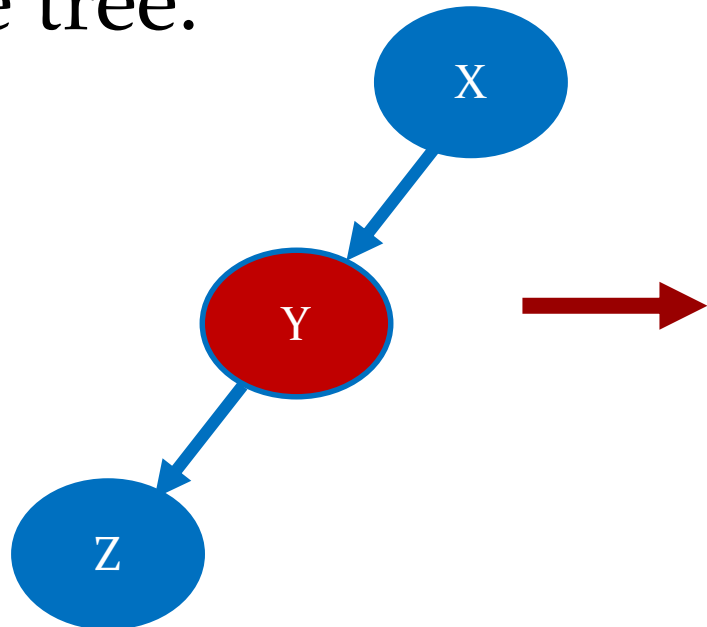
- The 90–10 rule states that 90% of the accesses are to 10% of the data items. However, balanced search trees do not take advantage of this rule.
- A cache stores in main memory the contents of some of the disk blocks. The goal is that when a disk access is requested, the block can be found in the main memory cache and thus save the cost of an expensive disk access.
- Browsers make use of the same idea: a cache stores locally the previously visited Web pages.

# Splay Rotations

- When an element is accessed in Splay tree, tree rotations are used to move that element to the top of the tree.
- Type of rotations:
  1. Simple rotation
  2. Zig-Zig & Zag-Zag rotation
  3. Zig-Zag rotation

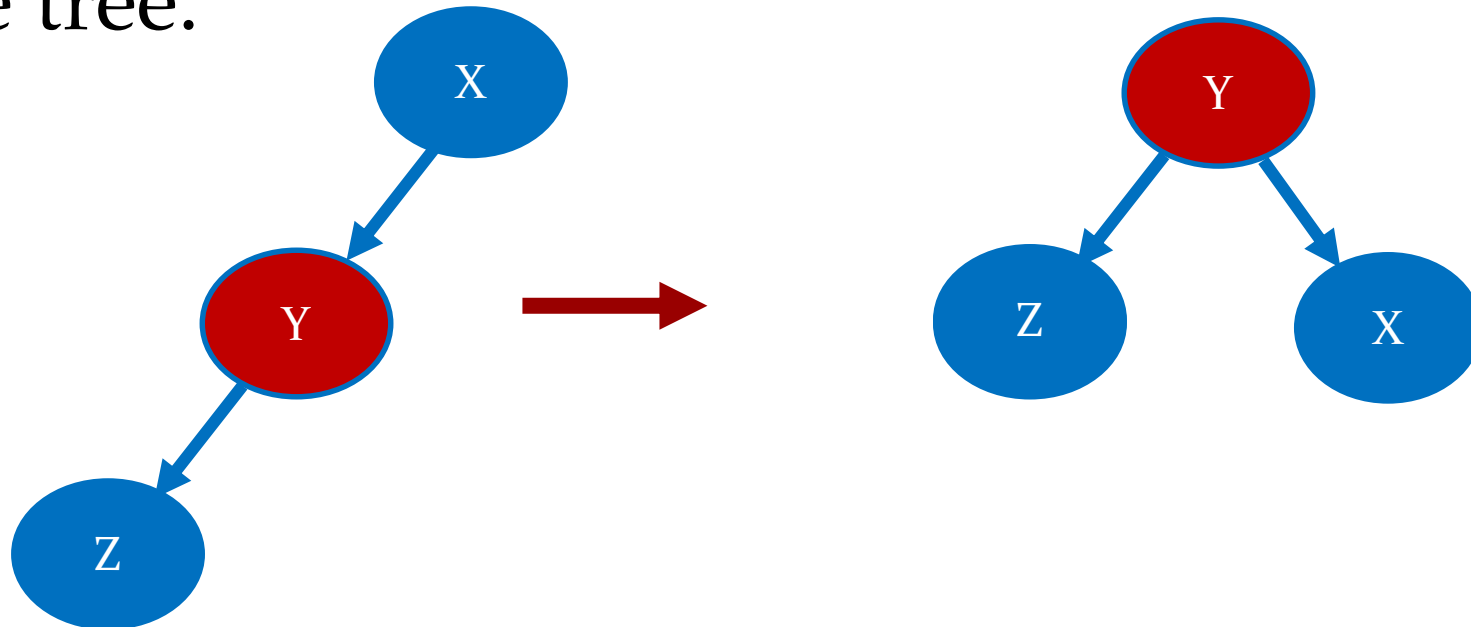
# Simple Rotation

- This is similar to the tree rotation used in AVL tree.
- In this case it is applied to the root of the splay tree, moving the splayed node Y up to become the new root of the tree.



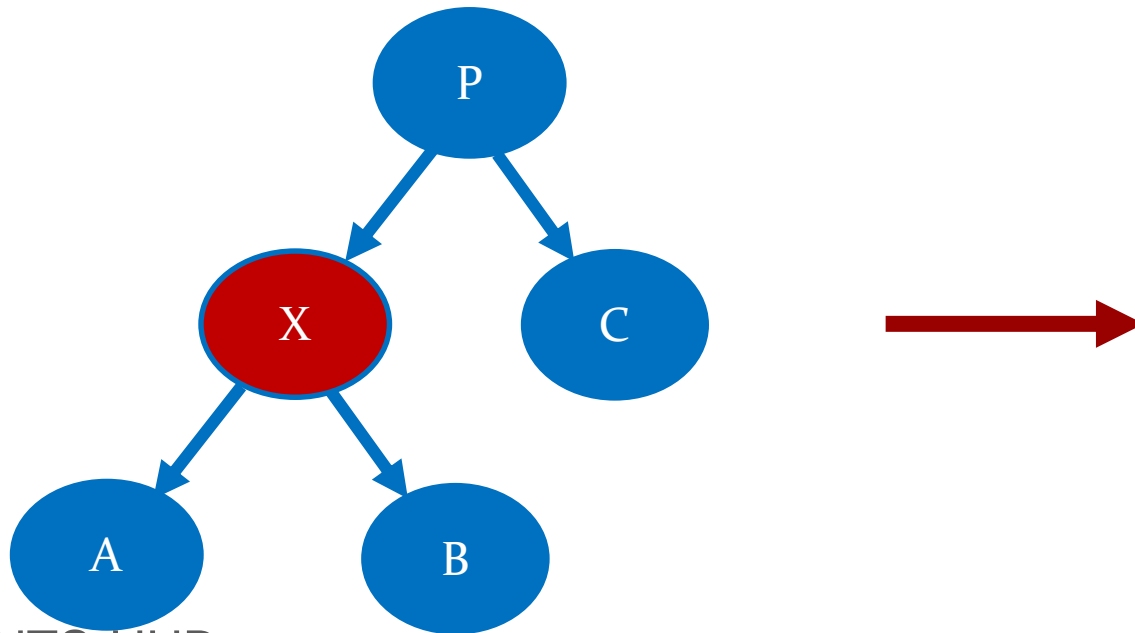
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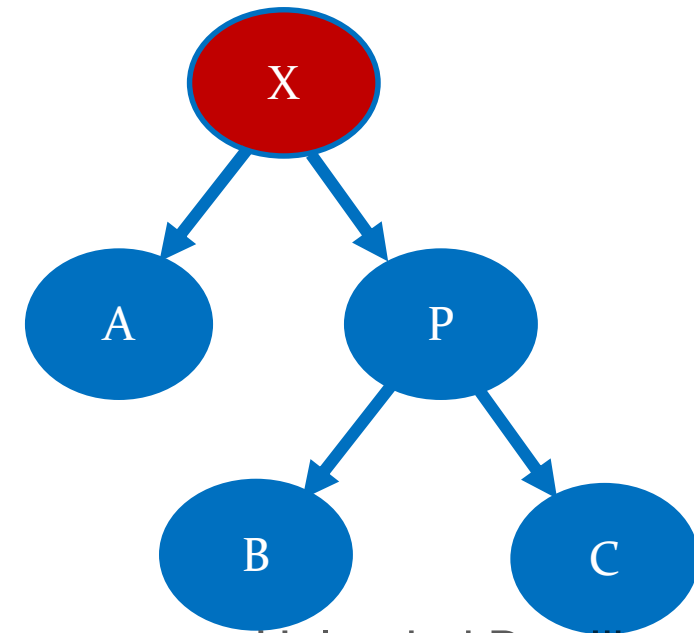
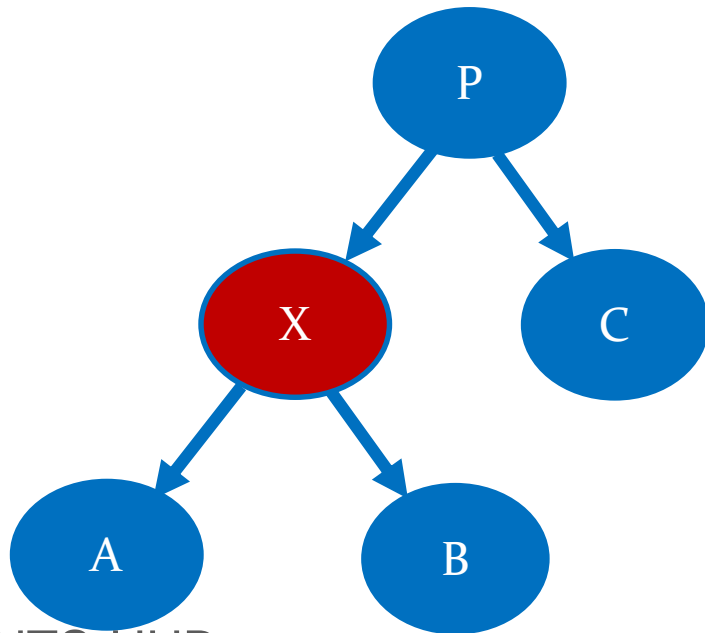
- The Zig-case and the Zag-case are the same as simple case
  - Let  $X$  be a non-root node on access path which we are rotating
  - If the parent of  $X$  is the root, then rotate  $X$  and the root





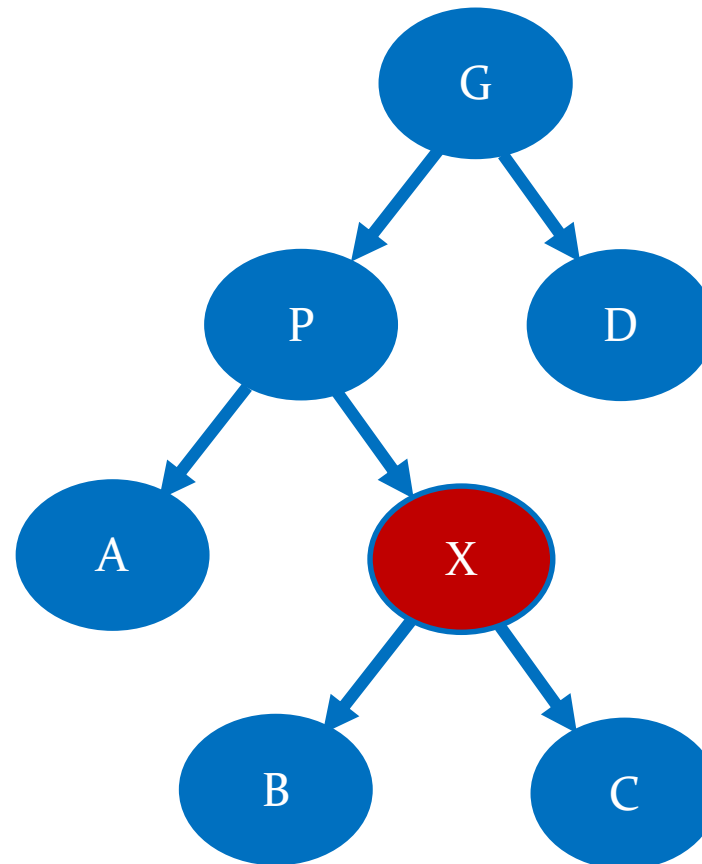
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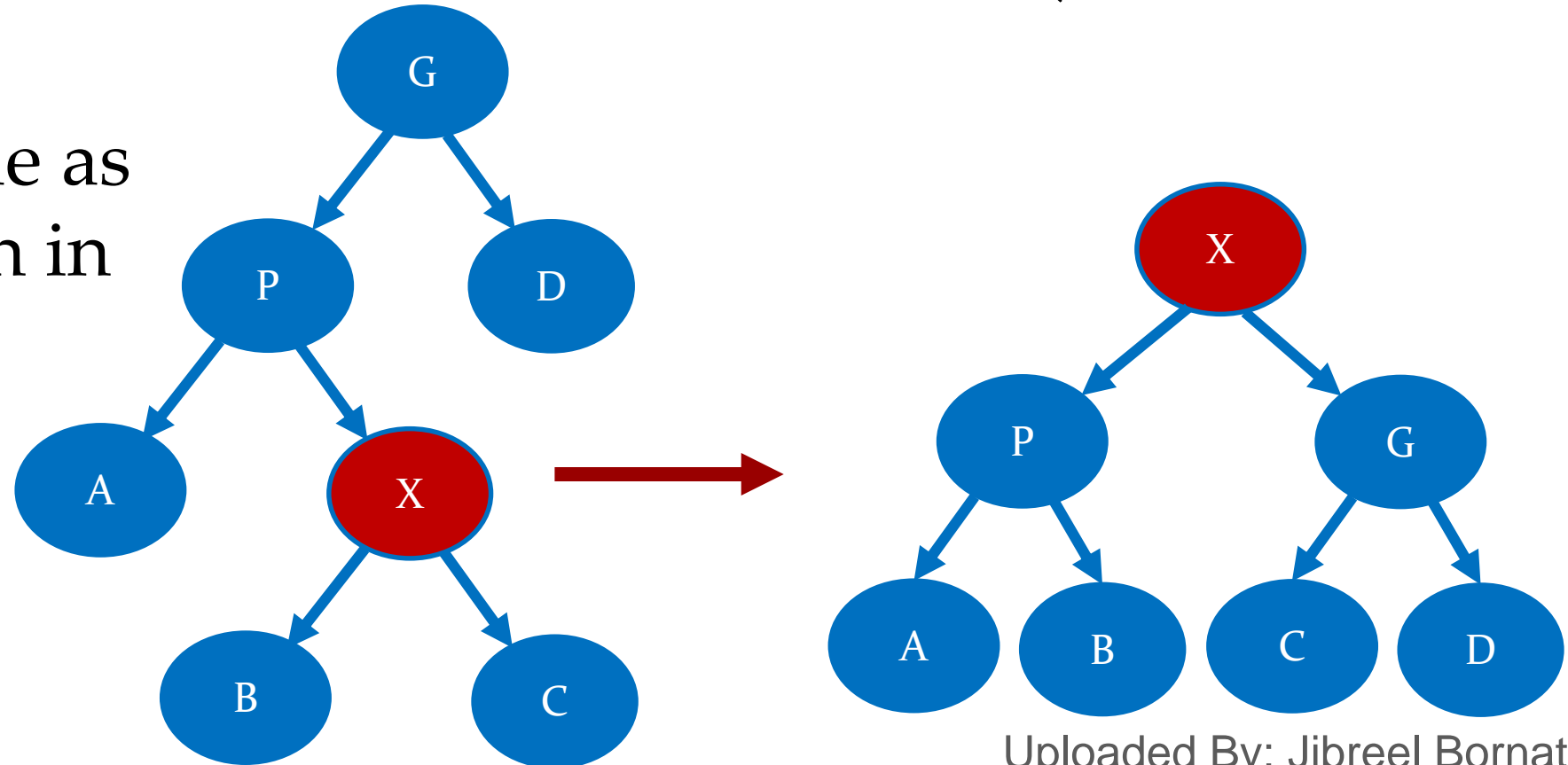
## Zig-Zag case

- In this case, X and both a parent P and a grandparent G, X is the right child of P and P is a left child of G (or vice versa).
- This is the same as double rotation in AVL tree



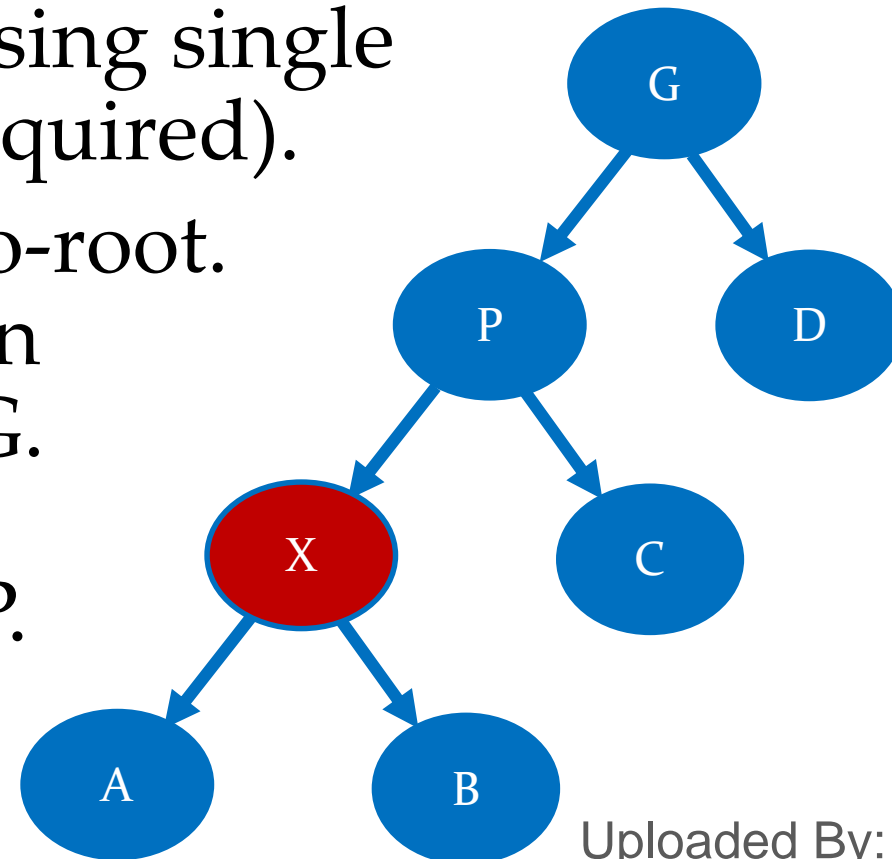
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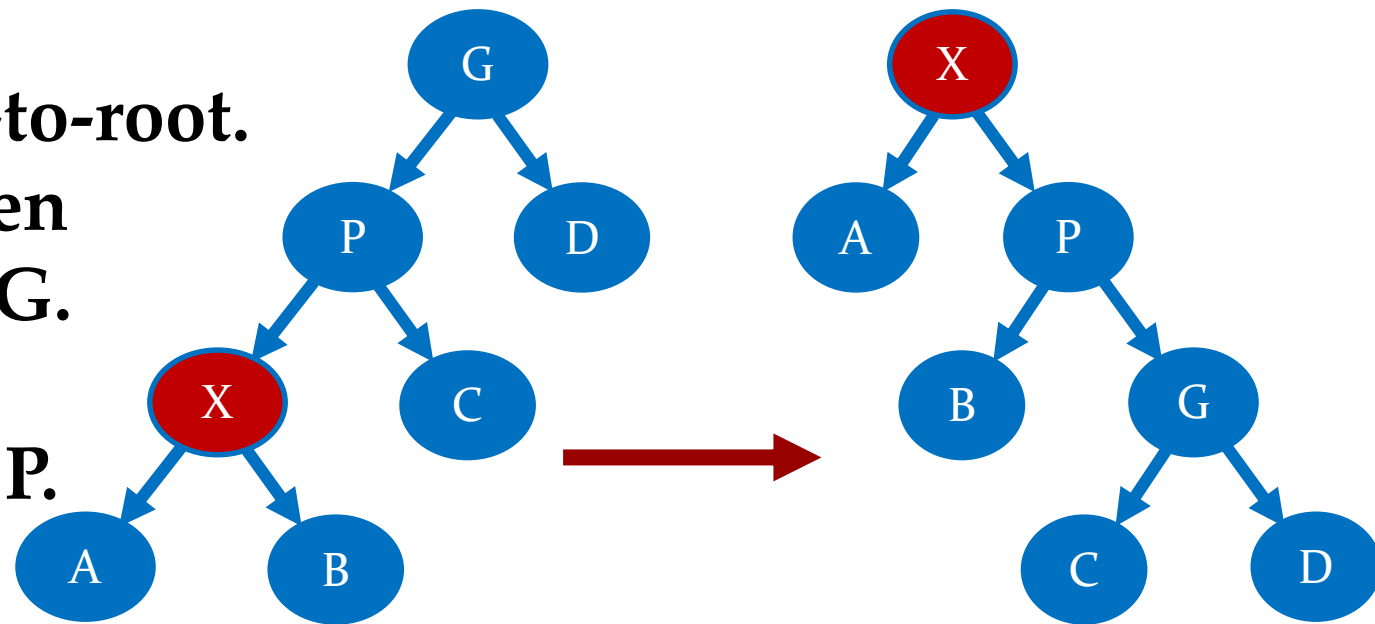
## Zig-Zig case

- Here X and P are either both left children or both right children.
- The transformation is done using single rotations (twice or more as required).
- This is different than rotate-to-root. Rotate-to-root rotates between X and P and between X and G. The zig-zig splay rotates between P and G and X and P.
- This case is not in AVL Trees.



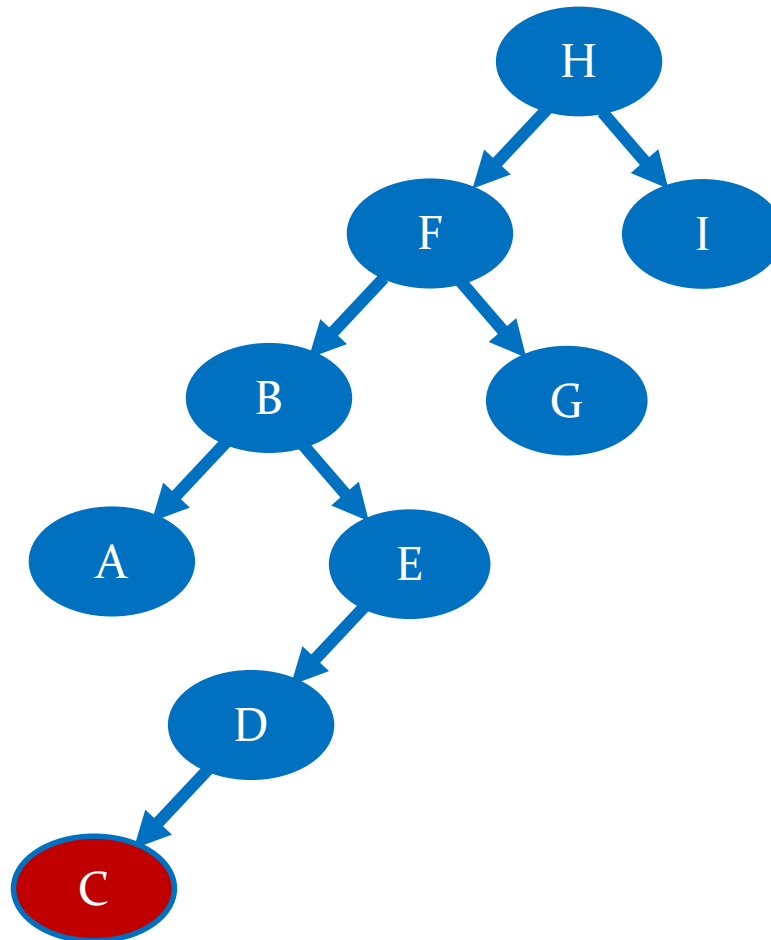
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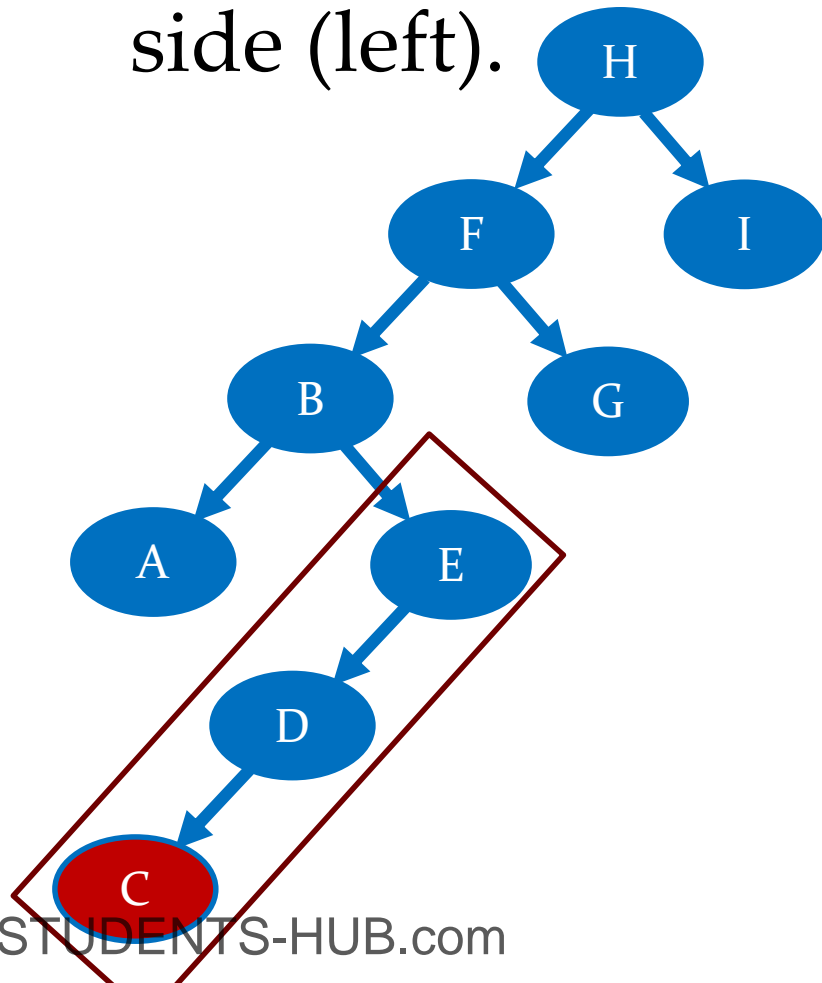
# Example

- Splay the node C in the following tree.



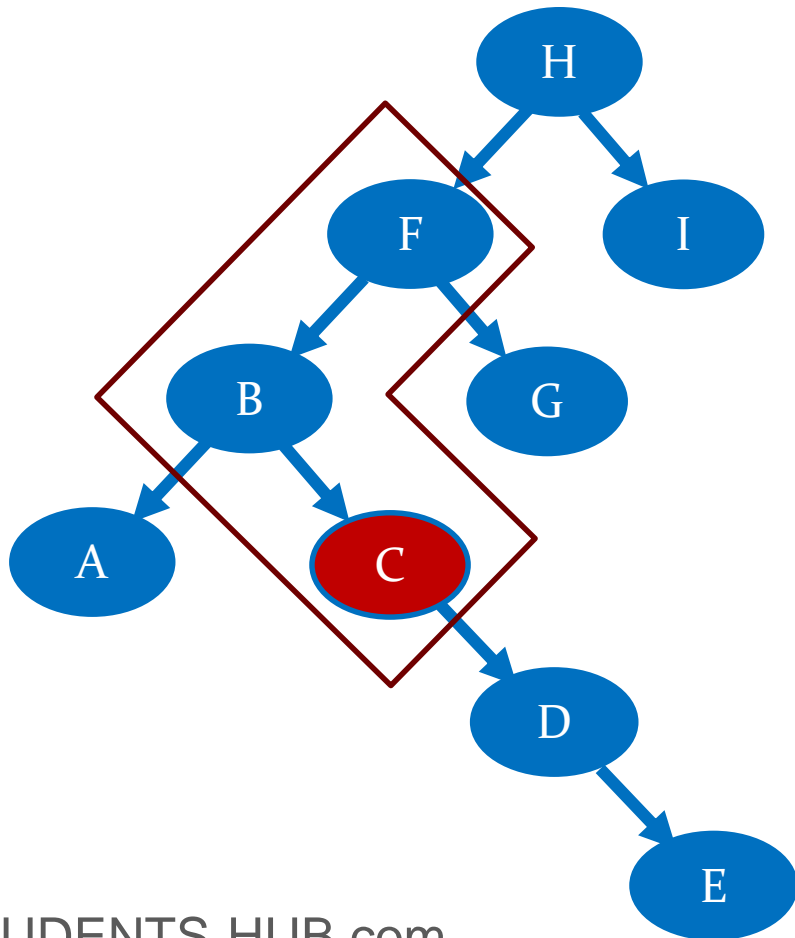
# Example

- Step1: Zig-Zig because C and its ancestors are on the same side (left).



# Example

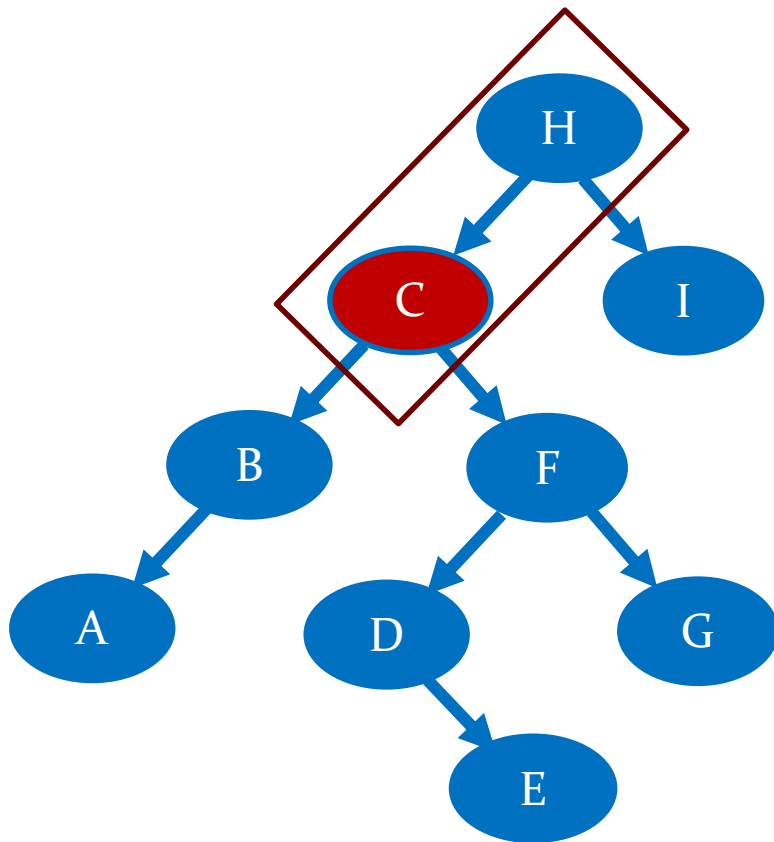
- Step2: Zig-Zag because C is right to B and B is left to F.





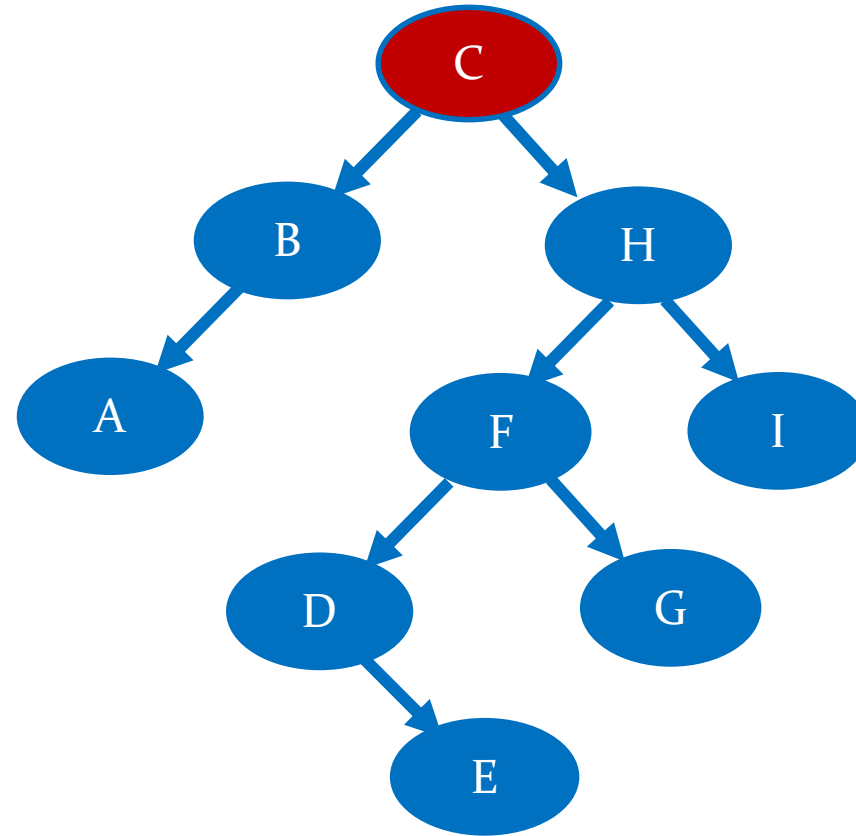
# Example

- Step3: Zig (single rotation) with the root H.



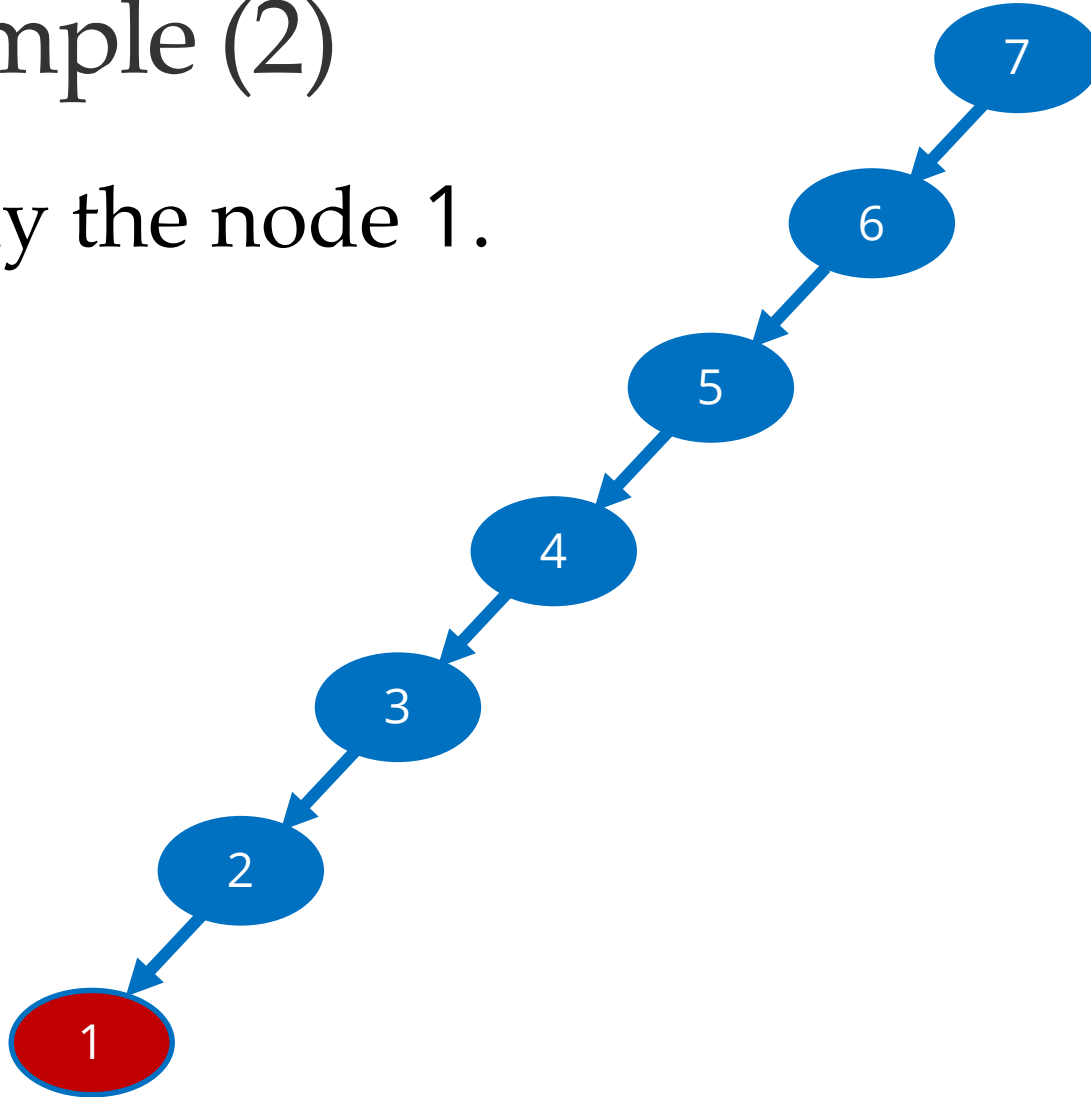
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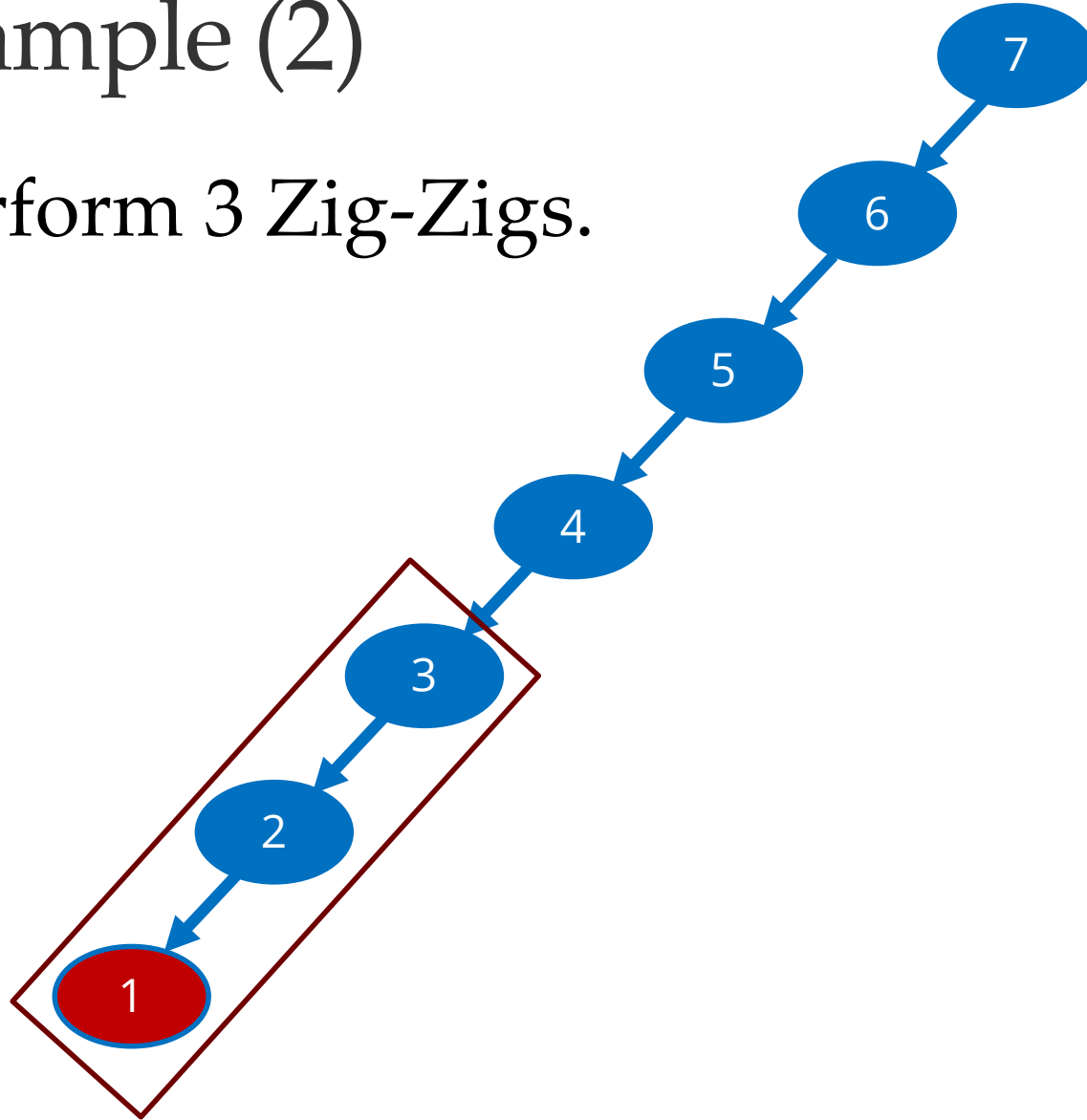
## Example (2)

- Splay the node 1.



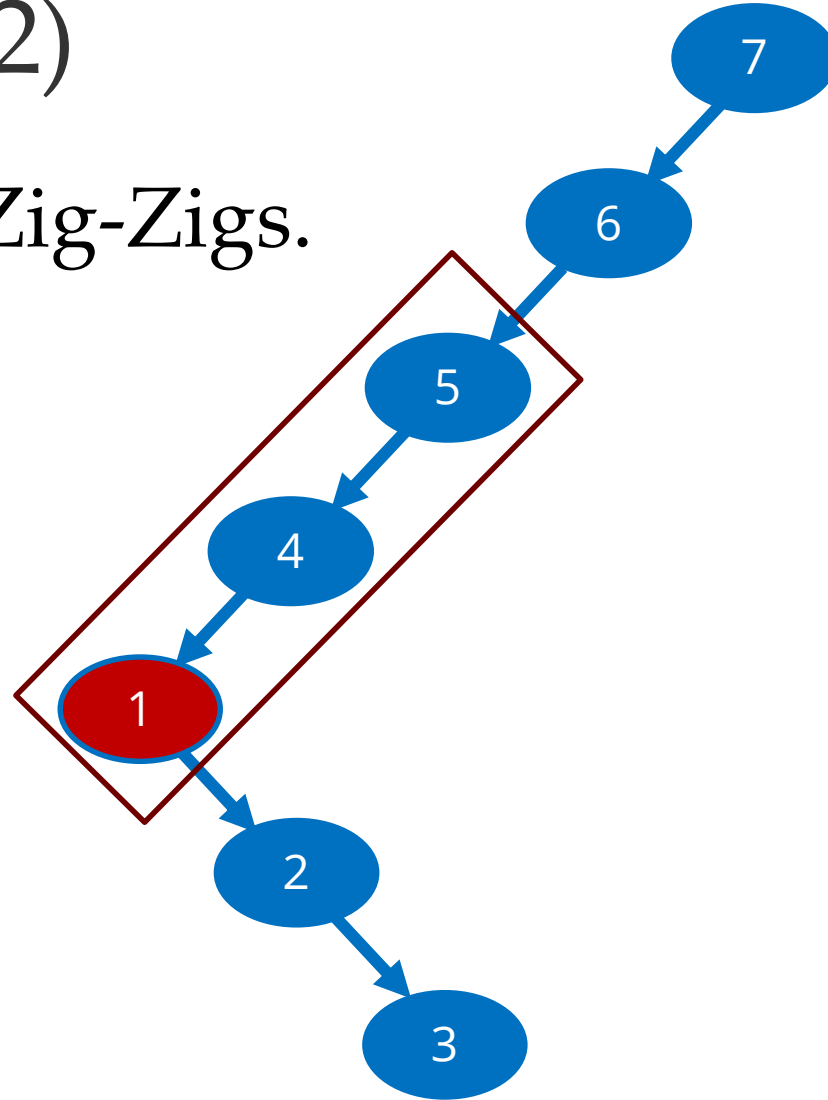
## Example (2)

- Perform 3 Zig-Zigs.



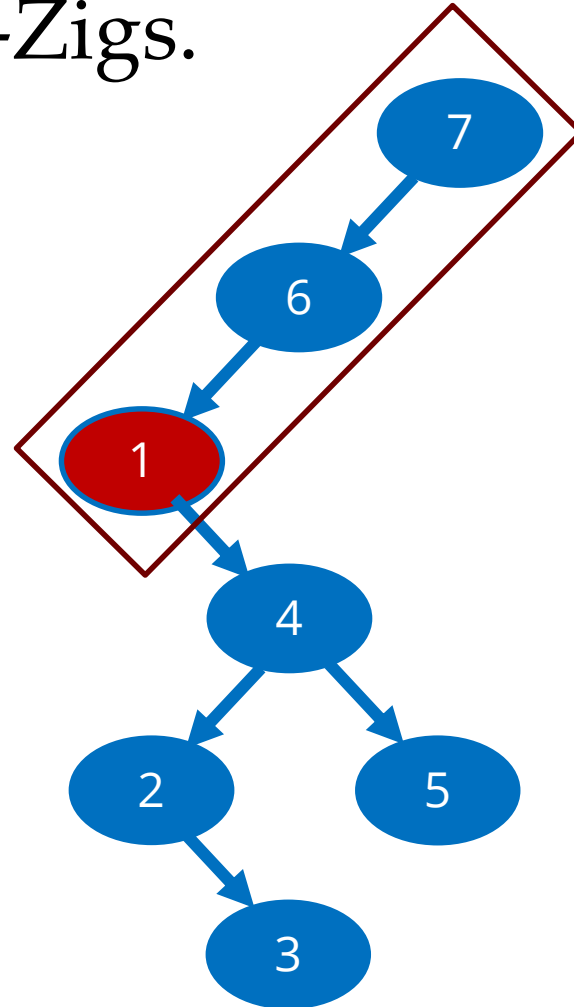
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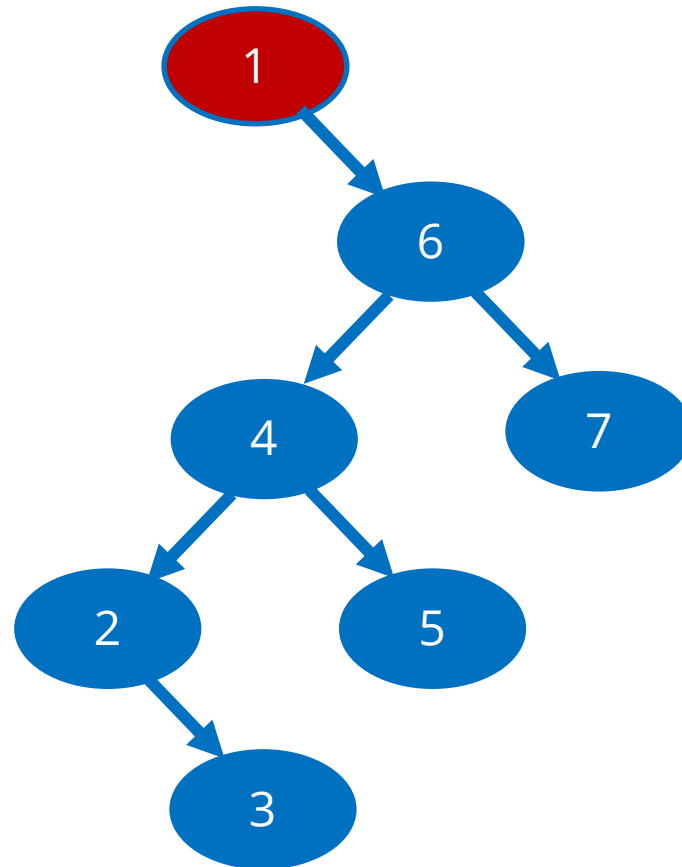
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# Basic Operations on Splay Trees

- Insertion: when an element is inserted at the right position, a splay is performed. As a result, the newly inserted node becomes the root of the tree.
- Find: the last node accessed during the search is splayed:
  - If the search is successful, the node found becomes the root of the tree.
  - Unsuccessful search, the last node accessed prior to reaching NULL, is splayed to become the root of the tree.
- FindMin & FindMax: perform a splay after the access.

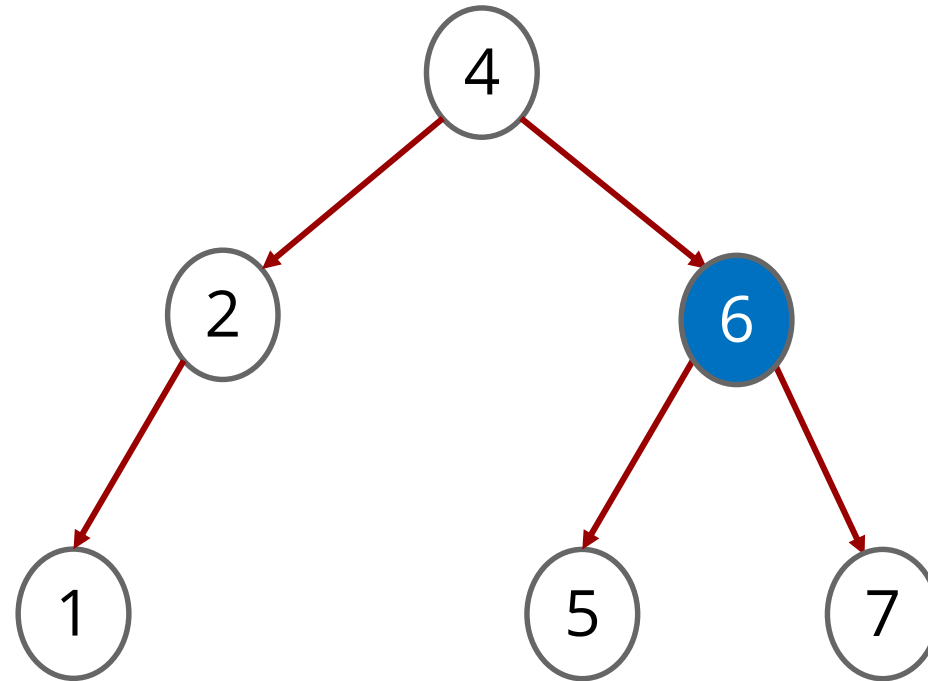


## Basic Operations on Splay Trees (2)

- Delete operation:
  - Access the node to be deleted bringing it to the root.
  - Delete the root leaving two subtrees L (left) and R (right).
  - Find the largest element in L using FindMax operation, thus the root of L will have no right child.
  - Make R the right child of L's root.

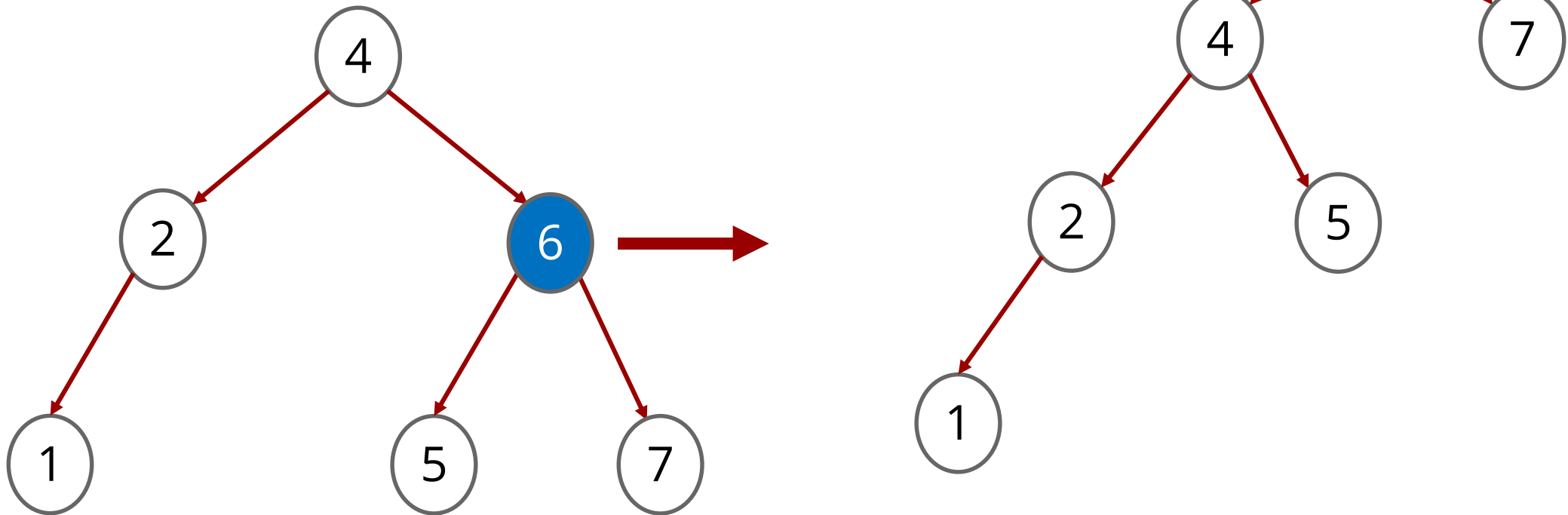
# Basic Operations on Splay Trees (3)

- Remove node 6.



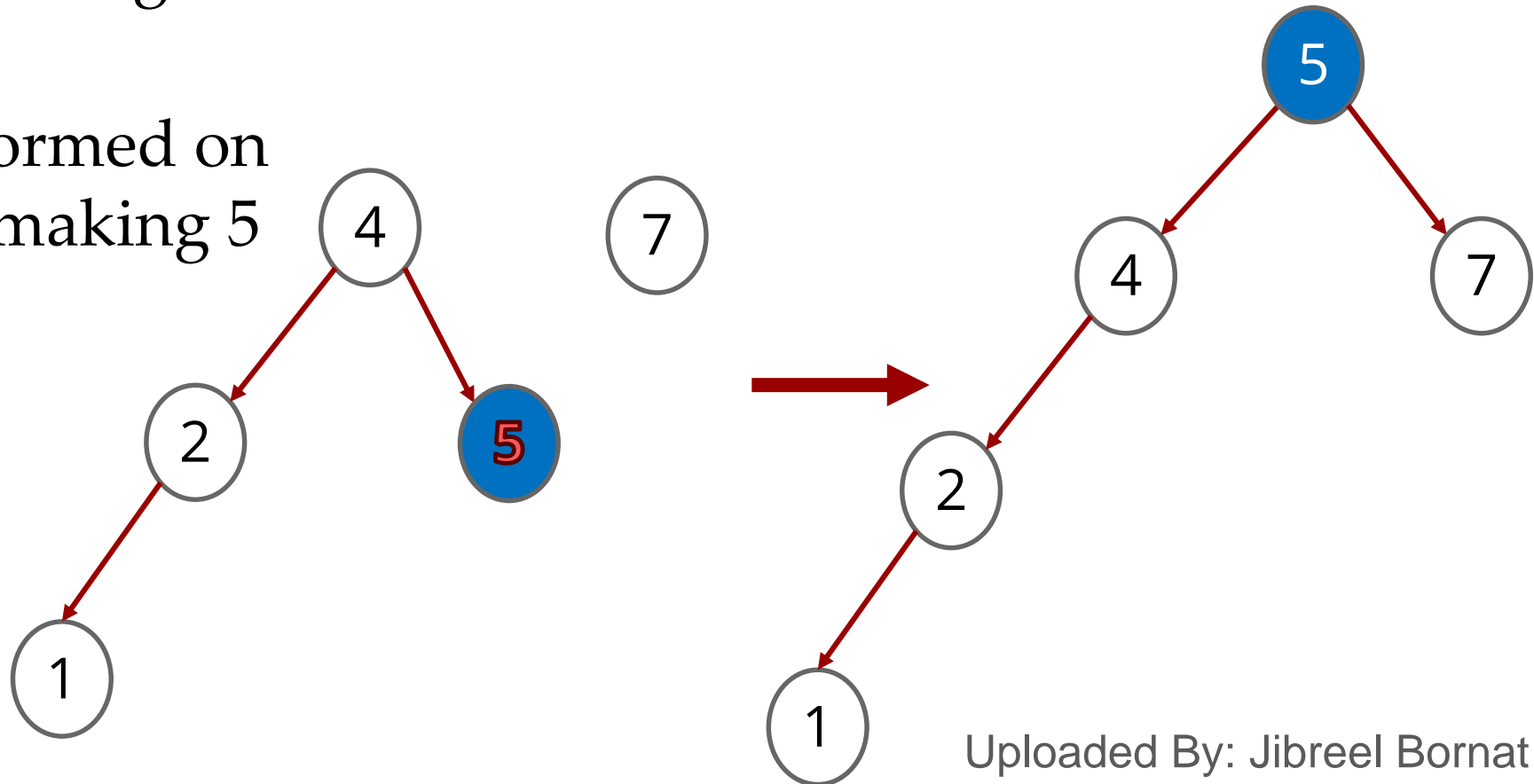
# Basic Operations on Splay Trees (3)

- Remove node 6.
  - Splay 6 to the root



# Basic Operations on Splay Trees (3)

- Remove node 6.
  - Delete node 6, leaving two subtrees.
  - FindMax is performed on the left subtree, making 5 the new root.



# Applications to Splay Tree

- Splay trees are the fastest balanced search trees.
- Network routers receiving packets & should decide where it should travel to. If an IP is used once, it is most likely to be used again. Thus, the look-up table is built using Splay tree.
- Implementation of cache, memory allocator, garbage collector, data compression, Intrusion Detection Systems.
- Unix *malloc*, GCC compiler, and sed string editor.

# Time Complexity

- Search:  $O(\log n)$
- Insert:  $O(\log n)$
- Delete:  $O(\log n)$