

Time Comp. Recursion \rightarrow Recurrence

Ex. Factorial

```
long fact (int n) {
```

```
    if (n == 0)
        return 1;
```

```
    return n * fact(n-1);
}
```

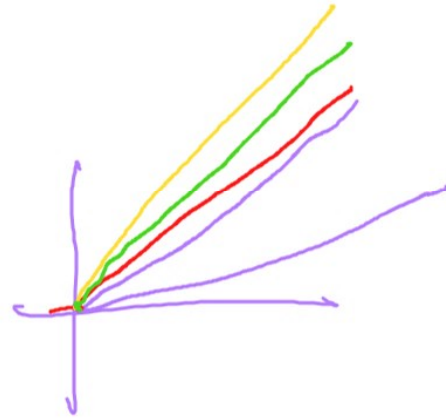
2g

✓

Rec.
Relation $T(n) = \begin{cases} d, & n=0 \\ T(n-1) + c, & n>0 \end{cases}$

$$\begin{aligned}
 T(n) &= T(n-1) + c \\
 T(n-1) &= T(n-2) + c \\
 T(n-2) &= T(n-3) + c \\
 &\vdots \\
 T(1) &= T(0) + c \\
 T(0) &= d
 \end{aligned}$$

$$\begin{aligned}
 T(n) &= nc + d \\
 &= O(n)
 \end{aligned}$$



MergeSort(A, L, H) {
 if (L < H)
 middle = (L + H) / 2

✓ MergeSort(A, L, middle);
 ✓ MergeSort(A, middle + 1, H);
 Merge(A, L, H);

}

A[0...n-1]
 (A[0, n])

R.R.

$$T(n) = 2T\left(\frac{n}{2}\right) + n \text{ merge}$$

$$d, n=1$$

$$T(n) = 2T\left(\frac{n}{2}\right) + n \quad \text{--- (1)}$$

$$T\left(\frac{n}{2}\right) = 2T\left(\frac{n}{2^2}\right) + \frac{n}{2} \quad \text{--- (2)}$$

$$T\left(\frac{n}{2^2}\right) = 2T\left(\frac{n}{2^3}\right) + \frac{n}{2^2} \quad \text{--- (3)}$$

$$T\left(\frac{n}{2^3}\right) = 2T\left(\frac{n}{2^4}\right) + \frac{n}{2^3} \quad \text{--- (4)}$$

Now (2) in (1)

$$T(n) = 2 \left[2T\left(\frac{n}{2^2}\right) + \frac{n}{2} \right] + n$$

$$= 2^2 T\left(\frac{n}{2^2}\right) + 2n \quad \text{--- (5)}$$

Sub. (3) in (5)

$$T(n) = 2^2 \left[2T\left(\frac{n}{2^3}\right) + \frac{n}{2^2} \right] + 2n$$

$$= 2^3 T\left(\frac{n}{2^3}\right) + 3n \quad \text{--- (6)}$$

After K of steps.

$$T(n) = 2^K T\left(\frac{n}{2^K}\right) + Kn$$

Let $\frac{n}{2^K} = 1$

$$n = 2^K$$

$$K = \log n$$

$$T(n) = nT(1) + n \log n$$

$$= n \cdot c + n \log n$$

$$= O(n \log n)$$

$$T(n) = \begin{cases} d, & n=1 \\ 2T(\frac{n}{2}) + 10, & n>1 \end{cases}$$

$$T(n) = 2T(\frac{n}{2}) + 10 \quad \text{--- (1)}$$

$$T(\frac{n}{2}) = 2T(\frac{n}{2^2}) + 10 \quad \text{--- (2)}$$

$$T(\frac{n}{2^2}) = 2T(\frac{n}{2^3}) + 10 \quad \text{--- (3)}$$

Sub. (2) in (1)

$$\begin{aligned} T(n) &= 2[2T(\frac{n}{2^2}) + 10] + 10 \\ &= 2^2 T(\frac{n}{2^2}) + 2^1 \cdot 10 + 2^0 \cdot 10 \quad \text{--- (4)} \end{aligned}$$

Sub (3) in (4)

$$\begin{aligned} T(n) &= 2^2 [2T(\frac{n}{2^3}) + 10] + 2^1 \cdot 10 + 2^0 \cdot 10 \\ &= 2^3 T(\frac{n}{2^3}) + 2^2 \cdot 10 + 2^1 \cdot 10 + 2^0 \cdot 10 \end{aligned}$$

at k^{th} step

$$T(n) = 2^k T(\frac{n}{2^k}) + 2^{k-1} \cdot 10 + 2^{k-2} \cdot 10 + \dots + 2 \cdot 10 + 2^0 \cdot 10$$

$$\frac{n}{2^k} = 1$$

$$n = 2^k$$

$$\begin{aligned} &10 \cdot (2^{k-1} + 2^{k-2} + \dots + 2^1 + 2^0) \\ &\frac{2^k - 1}{2 - 1} = (2^k - 1) \cdot 10 \end{aligned}$$

$$\begin{aligned} T(n) &= nT(1) + (n-1) \cdot 10 \\ &= d \cdot n + 10n - 10 \\ &= O(n) \end{aligned}$$

$$\frac{r^n - 1}{r - 1}$$

Uploaded By: anonymous

$$T(n) = 2T(n-1) + c \quad \text{--- (1)}$$

$$T(n-1) = 2T(n-2) + c \quad \text{--- (2)}$$

$$T(n-2) = 2T(n-3) + c \quad \text{--- (3)}$$

$$\begin{aligned} T(n) &= 2[2T(n-2) + c] + c \\ &= 2^2 T(n-2) + 2c + c \quad \text{--- (4)} \end{aligned}$$

$$\begin{aligned} T(n) &= 2^2 [2T(n-3) + c] + 2c + c \\ &= 2^3 T(n-3) + 2^2 c + 2c + c \end{aligned}$$

after k^{th} step.

$$T(n) = 2^k T(n-k) + 2^{k-1} c + 2^{k-2} c + \dots + c$$

$n-k=1$ will stop

$$n = k+1 \rightarrow k = n-1$$

$$T(n) = 2^{n-1} T(1) + 2^{n-2} \cdot c + 2^{n-3} \cdot c + \dots + c$$

Assume n is very large

$$n \approx n-1$$

$$\begin{aligned} T(n) &= 2^n T(1) + 2^{n-1} \cdot c + 2^{n-2} \cdot c + \dots + c \\ &= d \cdot 2^n + (n-1) \cdot c \\ &= O(2^n) \end{aligned}$$