

chapter 2 : Application layer

Internet protocol stack: application \Rightarrow IMAP, SMTP, HTTP
↓
FTP

- server-client paradigm :

server \Rightarrow always on, fixed IP

client \Rightarrow intermittently connected
, dynamic IP

ex: HTTP, IMAP, FTP

Transport \Rightarrow TCP, UDP

↓
Network \Rightarrow IP, routing protocols

↓
Link \Rightarrow Ethernet, WiFi, PPP

↓
physical \Rightarrow wire

- peer-peer architecture

α no always on server α end systems communicate directly

α peers request & provide services from each other

ex: P2P file sharing

- process : program running within a host.

- inter-process communication : how to processes in the same host communicate.

- processes on different hosts communicate by exchanging messages

- client processes : process that initiates communication.

- server process : process that waits to be contacted.

} P2P has both

- socket : process sends / receives messages to / from its socket.

- identifier of a process is both IP address & port number
 \hookrightarrow 32 bit

- port of HTTP server is 80 & mail server is 25

	TCP	UDP
reliability	✓	X
flow control	✓	X
congestion control	✓	X
connection oriented	✓	X
timing, minimum throughput, security	X	X

• then why is there UDP ?? because its simple & doesn't need setup, so when we don't need reliable transport we use it

• Vanilla TCP & UDP sockets:

- No encryption
- clear text passwords sent into socket through internet in clear text

• TLS (transport layer security)

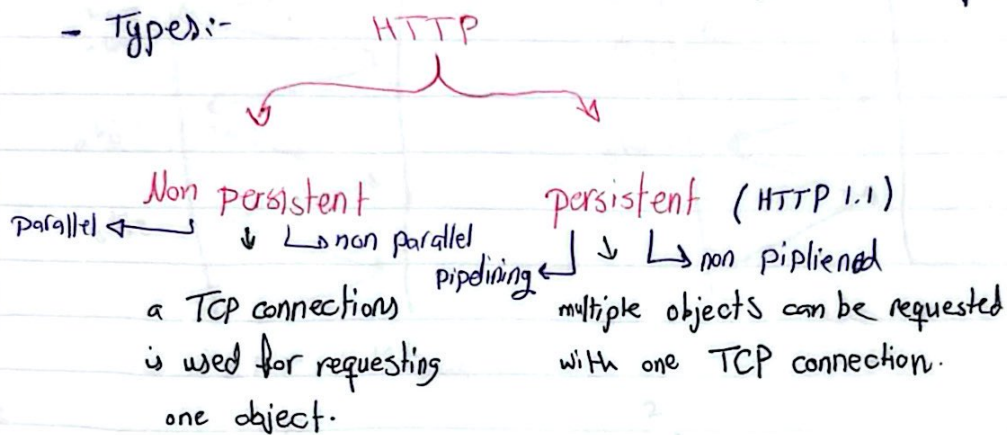
- encrypted TCP con.
 - data integrity
 - end point authentication
- in application layer (apps use TLS libraries)

HTTP (hypertext transfer protocol)

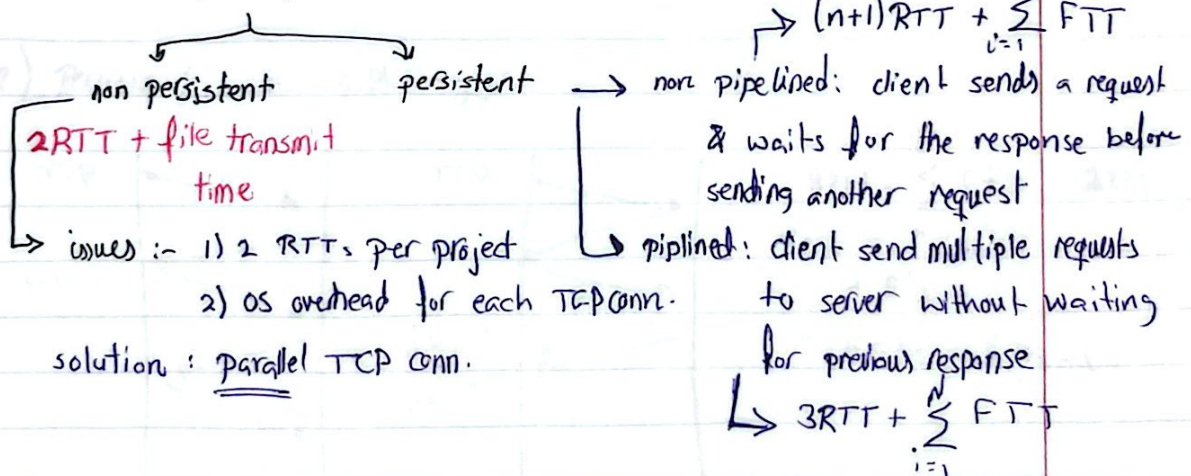
- web's application layer protocol
- client/server
- TCP used \Rightarrow ① client initiates TCP conn. & creates socket to port 80.
② server accepts TCP
③ HTTP messages exchanged between browser & webserver.
④ TCP connection closed

- stateless \rightarrow server has no history about past client requests

- Types:-



- RTT: time for a packet to travel from client to server & back.
- HTTP response time (per project)



HTTP messages

① HTTP request message

ASCII

- request line → Get html
- header lines start

end

⇒ \r\n in a line alone mean the end of header lines.

- body
types:

- post method:- user input
- Get method:- (for data sending to a server after url & ?)
- head method:- requests only headers that would be returned if specified URL were requested with an HTTP Get method.
- put method:- uploads new file (object) to server

② HTTP response message:-

- status line (protocol, status code, status phrase)
- header lines ⇒ end with line \r\n
- data requested

status codes:-

- | | | |
|---|---|--|
| I. 200 OK
object request
succeeded its
in this message | II. 301 moved permanently
requested object moved,
new location in this
message (field) | III. 400 Bad Request
request message
not understood
by servers. |
| IV. 404 Not Found
requested object not
on this server | V. 505 HTTP version not supported | |

- stateful protocol :: client makes to changes to x or none at all

- cookies are used to maintain user/server state between transactions.

cookie components

- ⇒
- 1) cookie header line of HTTP response message.
 - 2) cookie = = in next HTTP request message
 - 3) cookie file kept on user's host managed by user's browsers.
 - 4) back-end database at website.

- use of cookies :-

• authorization • shopping carts • recommendations • user session state.

* Web caches (proxy servers)

⇒ satisfy client request without involving origin server.

⇒ they act like both client & server, client when they request sth from origin server & server when they response to an existing response.

⇒ an object allowable of caching is in response header:

cache-control: max-age = <seconds>
no-cache

⇒ cache is installed by ISP

* conditional get is used so an object is not sent if cache has up to date version

cache:- if modified since: (date)

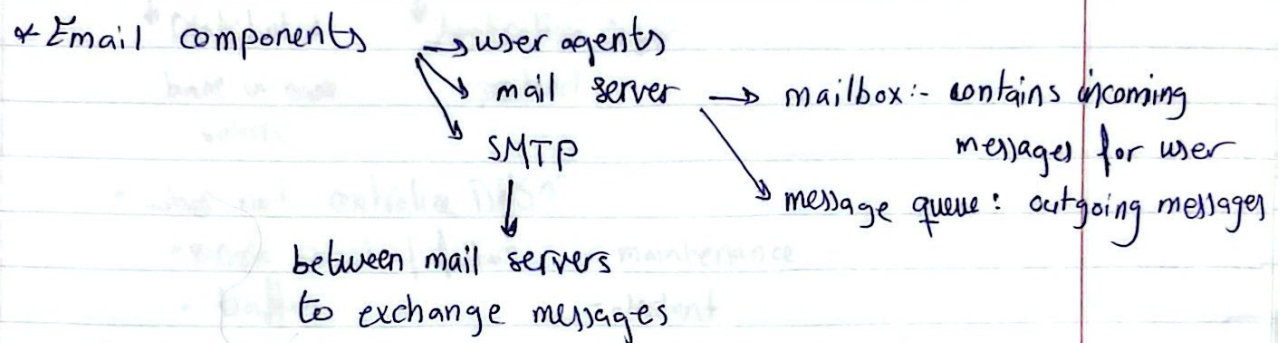
server:- no object if up to date

HTTP/1.0 304 Not Modified

HTTP 1.1 ⇒ multiple pipelined gets over a TCP connection.
 ⇒ FCFS responds by server to get requests.
 ⇒ loss recovery stalls object transmission.

HTTP 2 ⇒ methods, status codes, header fields same as HTTP 1.1
 ⇒ transmission order of objects is by priority
 → divides objects to frames to lessen HOL Blocking
 → same as HTTP 1.1 in loss recovery & no security over vanilla TCP connection.

HTTP 3 → added security



- ① client sends their message to their mail server.
- ② client mail server of SMTP opens TCP conn. with server mail.
- ③ SMTP client sends message over TCP
- ④ message is put in SMTP server in mailbox.

SMTP RFC (5321) ⇒ uses TCP, port 25

① TCP ③ transfer

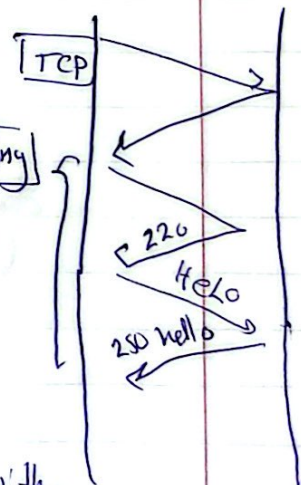
② handshaking ③ closure of SMTP handshaking

- commands: ASCII

- response: status code & phrase.

↳ 220 SMTP ready
 221 service closing
 250 Request completed
 354 start message input & end with =

MAIL FROM DATA
 RCPT TO QUIT

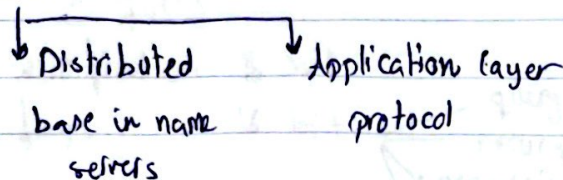


- SMTP: protocol for exchanging email messages defined in RFC #5821
- HTTP is defined in RFC 7281
- RFC 2822 defines syntax for email like html defines syntax web browsers.
- IMAP (Mail access protocol) :- retrieval from server (RFC 3501)

* DNS

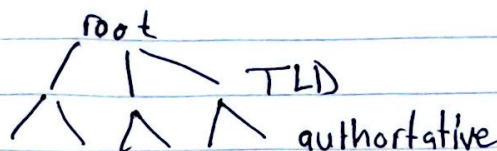
- IP 32 bit
- name

DNS is used to map between IP address & name.



• why not centralize DNS??

- single point of failure
- maintenance.
- traffic
- elastant



• Local DNS : \Rightarrow Does not belong to hierarchy
 \Rightarrow each ISP has one

DNS: distributed database storing resource records.
 RR format (name, value, type, ttl)

- RR types :-

① A

(host name, IP, A, TTL)

② NS

(domain, hostname of authoritative, NS, TTL)
name server for this domain

③ CNAME

(alias, canonical, CNAME, TTL)

④ MX

(~~alias~~ domain, canonical, MX, TTL)

* DNS

→ reply
→ query

identification 16 bit

flag 16 bit

query 0 reply 1
recursion desired
recursion available
reply is authoritative