

- Fixed point . $P = g(P)$

- Fixed point Iteration theorem :

Assume $g \in C[a,b]$

- If $g(x) \in [a,b]$ for all $x \in [a,b]$, then g has a fixed point in $[a,b]$.

and • If $|g'(x)| < K < 1$ for all $x \in [a,b]$, then g has a unique fixed point in $[a,b]$.

- Fixed point Iteration theorem :

Assume the following $g, \bar{g} \in C[a,b]$, K is constant positive, $P_0 \in [a,b]$ is the starting point $g(x) \in [a,b]$ for all $x \in [a,b]$.

- If $|g'(x)| \leq K < 1$ for all $x \in [a,b]$, then the FPI converge to the unique FP, P
(P is called attractive fixed point).

- If $|g'(x)| > 1$ for all $x \in [a,b]$, the FPI diverge.
(P is called repelling fixed point).

- Corollary :

Assume g satisfies conditions of FPI Theorem 2↑. If P_n is used to approximate the fixed point P , then an upper bounds for the error are:

A. $|P - P_n| \leq K^n |P - P_0|$ for all $n \geq 1$.

B. $|P - P_n| \leq \frac{K^n}{1-K} |P_1 - P_0|$ for all $n \geq 1$.

* K upper bound of $|g'(x)|$

... \Rightarrow if $n \rightarrow \infty$, $O \rightarrow 0$

- Bisection Method of Bolzano:**
- This method uses to solve $f(x)=0$, and depends on Bolzano theorem.
 - conditions required : $f \in C[a,b]$ and $f(a)f(b) < 0$.
 - First iteration $c_0 = \frac{a_0 + b_0}{2}$ (midpoint).
 - Now find $f(c_0)$, we have three cases:
 - if $f(c_0) = 0$, then c_0 is the root and we are done.
 - if $f(c_0)f(a_0) < 0$, then $[a_1, b_1] = [a_0, c_0]$ and $c_1 = \frac{a_1 + b_1}{2} = \frac{a_0 + b_0}{2}$.
 - if $f(c_0)f(b_0) < 0$, then $[a_1, b_1] = [c_0, b_0]$ and $c_1 = \frac{a_1 + b_1}{2} = \frac{c_0 + b_0}{2}$.

In general, $c_n = \frac{a_n + b_n}{2}$ is Bisection method.

Bisection Theorem :

- Assume that $f \in C[a,b]$, \exists a number $r \in [a,b]$ s.t $f(r)=0$, $f(a)f(b) < 0$ and $\{c_n\}_{n=0}^{\infty}$ represents the sequence of midpoints generated by the bisection method.
- Then an upper bound of the error is $|r - c_n| \leq \frac{b-a}{2^{n+1}}$.
- The sequence $\{c_n\}_{n=0}^{\infty}$ converges to the zero r .

Remark :

The Bisection theorem provides a starting to find the number of iteration for a given accuracy δ :

$$\frac{b-a}{2^{n+1}} < \delta \rightarrow n > \frac{\ln(\frac{b-a}{\delta})}{\ln 2} - 1$$

* False position Method. (converges faster).

- conditions to solve $f(x)=0$: $f \in C[a,b]$ and $f(a)f(b) < 0$.
- As in Bisection method , we have three cases
 - If $f(a_0)f(c_0) < 0$, then the zero $r \in [a_0,c_0]$.
 - If $f(c_0)f(b_0) < 0$, then the zero $r \in [c_0,b_0]$.
 - If $f(c_0)=0$, then the zero $r=c_0$.

In general : $c_n = b_n - \left(\frac{b_n - a_n}{f(b_n) - f(a_n)} \right) f(b_n)$.

* Newton's method iteration : $p_{n+1} = p_n - \frac{f(p_n)}{f'(p_n)}$, $n=0,1,2,\dots$ (Faster than Bisection and False position methods)

* Secant Method : $p_{n+2} = p_{n+1} - \left(\frac{p_{n+1} - p_n}{f(p_{n+1}) - f(p_n)} \right) f(p_{n+1})$. (as fast as Newton method)
 p_0, p_1 given.

* Multiplicity of Roots :

- we say $f(x)=0$ has a root of order N at $x=p$ (p has multiplicity M) iff
 $f(p)=0, f'(p)=0, f''(p)=0, \dots, f^{(M-1)}(p)=0, f^M(p) \neq 0$.

- Lemma : If $f(x)=0$ has a root p and \exists a continuous function $h(x)$ s.t

$$f(x) = (x-p)^M h(x) \text{ where } h(p) \neq 0 \text{ then the root } p \text{ has multiplicity } M.$$

* Speed of Convergence :

- If there exist two positive constants $A \neq 0$ and $R > 0$ s.t $\lim_{n \rightarrow \infty} \frac{|p - p_{n+1}|}{|p - p_n|^R} = \lim_{n \rightarrow \infty} \frac{|E_{n+1}|}{|E_n|^R} = A$

then we say that $\{p_n\}$ converges to p with order of convergence R .

- We use the order of convergence R to measure the speed of convergence of any method.

- If $R=1$ then the convergence of $\{p_n\}$ is linear
- If $R=\frac{3}{2}$ then the convergence of $\{p_n\}$ is super linear
- If $R=2$ then the convergence of $\{p_n\}$ is quadratic
- If $R=3$ then the convergence of $\{p_n\}$ is cubic.

- When $R \uparrow \rightarrow$ speed $\uparrow \rightarrow$ error \downarrow .

- A is called the asymptotic error constant.

* Speed of Newton's Method.

Assume Newton iteration $P_{n+1} = P_n - \frac{f(P_n)}{f'(P_n)}$, produces a sequence $\{P_n\}$ that converges to the root P of the function $f(x)$. Then :

1. If P is simple root ($M=1$), then Newton iteration $\{P_n\}$ converges to P

quadratically ($R=2$) with $A = \left| \frac{\bar{f}(P)}{2\bar{f}'(P)} \right|$ and $\lim_{n \rightarrow \infty} \frac{|E_{n+1}|}{|E_n|^2} = A$.

2. If P is a multiple root (of order $M \geq 1$) then Newton's iteration $\{P_n\}$ converges to P linearly ($R=1$) with $A = \frac{M-1}{M}$ and $\lim_{n \rightarrow \infty} \frac{|E_{n+1}|}{|E_n|} = A$.

* Accelerated Newton-Raphson iteration.

- Assume Newton-Raphson iteration $P_{n+1} = P_n - \frac{f(P_n)}{f'(P_n)}$, produces a sequence $\{P_n\}$ that converges to the root P .

- Assume P is a multiple root (of order $M \geq 1$). Then

1. Newton's iteration $\{P_n\}$ converges linearly ($R=1$) with $A = \frac{M-1}{M}$ and $\lim_{n \rightarrow \infty} \frac{|E_{n+1}|}{|E_n|} = A$.

2. The Modification of Newton's iteration $P_{n+1} = P_n - \frac{M f(P_n)}{\bar{f}'(P_n)}$, given P_0 , $n=0, 1, 2, \dots$

converges quadratically ($R=2$) to P and $A = \lim_{n \rightarrow \infty} \frac{|E_{n+1}|}{|E_n|^2}$.

* Speed of convergence for secant Method :

$$P_{n+2} = P_{n+1} - \left(\frac{P_{n+1} - P_n}{f(P_{n+1}) - f(P_n)} \right) f(P_{n+1}).$$

1. If P is simple root ($M=1$), then secant Method iteration $\{P_n\}$ converges to P

with $R = 1.618$ and $A = \lim_{n \rightarrow \infty} \frac{|E_{n+1}|}{|E_n|^{1.618}} = \left| \frac{\bar{f}(P)}{2\bar{f}'(P)} \right|^{0.618}$

2. If P is multiple root ($M \geq 1$), then secant Method iteration converges to P
with $R=1$ and A depends on $f(x)$.

* speed of convergence for Bisection Method :

$$R=1 \quad \text{and} \quad A = \frac{1}{2} \quad \rightarrow \quad \frac{|E_{n+1}|}{|E_n|} \approx \frac{1}{2}$$

* speed of convergence for False position Methods :

$$R=1, \quad A \text{ depends on } f(x), \quad \frac{|E_{n+1}|}{|E_n|} \approx A$$

* speed of convergence for Fixed point Iteration.

$$R=M, \quad A = \left| \frac{g^{(K)}(\rho)}{K!} \right|$$

Chapter 3:

3.3 : Solving $n \times n$ linear and nonlinear systems.

- Linear system : $Ax = b$

- Cost (complexity) : is the number of operations $+, -, \times, /$ required to complete a certain calculation.

- The cost of finding $|A_{nn}|$ for $n \geq 2$ is :

$$1. \text{cost} = n! \sum_{k=2}^n \frac{2k-1}{k!} \quad \text{or}$$

$$2. \text{cost} = [n/e - 2] \quad \text{where } [\frac{n}{e}] \text{ is the greatest integer function and } e \approx 2.718.$$

- If A is $n \times n$ matrix, then the cost of calculating A^2 is $2n^3 - n^2$.

- Backward substitution :

Used to solve a linear system of equations that has an upper-triangle coefficient matrix.

$$\text{cost of backward substitution} = n^2$$

- Cost of Gaussian Elimination - GE

1. The cost of reducing $n \times n$ linear system to an upper-triangle system using G.E is

$$\frac{4n^3 + 3n^2 - 7n}{6}$$

upper triangle system

2. The cost of solving $n \times n$ linear system using G.E is $\frac{4n^3 + 9n^2 - 7n}{6}$ (total cost)

* Gaussian Elimination and pivoting :

→ is two types :

1. Trivial pivoting : - If $a_{pp} \neq 0$, then do not switch rows.

- If $a_{pp} = 0$, then switch row P by the first row K below P s.t. $a_{Kp} \neq 0$.

2. Partial pivoting : - If $a_{pp} \neq 0$ or $a_{pp} = 0$, choose the pivot - $|a_{kp}| = \max \{ |a_{pp}|, |a_{p+1,p}|, \dots, |a_{n,p}| \}$.

- reducing the error being propagated when using a first digit arithmetic.

* LU-Factorization :

To solve the linear system $AX=b$:

- write $A=LU$ where L is lower triangle matrix. and U is upper triangle matrix.
- let $y=UX \Rightarrow AX=b$ becomes $LUX=b \rightarrow Ly=b$.
- Now solve $Ly=b$ by F.S and find y .
- Then solve $UX=y$ by B.S and find x .

⇒ cost of $LU = \text{cost of G.E.}$

Remark : The total cost of solving any linear $n \times n$ system $AX=b$ using LU-Factorization.

$$= \text{cost of } LU + \text{cost of F.S} + \text{cost of B.S}$$

$$= \frac{4n^3 - 3n^2 - n}{6} + n^2 - n + n^2 \xrightarrow{\text{cost on B.S}}$$

Cost of A=LU → *Cost of F.S (Ly=b)*

$$= 4n^3 + 9n^2 - 7n$$

$$= \frac{4n^3 + 9n^2 - 7n}{6}$$

$$= \text{cost of G.E.}$$

* Cramer's Rule .

→ To solve the linear $n \times n$ system $Ax = b$:

- Find $|A|$, $|A| \neq 0$.

- The solution $x = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$ is obtained as follow :

$x_i = \frac{|A_{ii}|}{|A|}$ where A_{ii} is obtained by replacing the i^{th} column of A by column b for all $i=1,2,3,\dots$.

→ Cost of Cramer Rule :

$$(n+1)D_n + n, \text{ where } D_n \text{ is the cost of } |A|.$$

* Solving nonlinear systems :

II Newton's Method .

- Given 2×2 nonlinear system

$$f(x,y) = 0, \quad g(x,y) = 0 \quad \text{with initial point } (x_0, y_0).$$

- This method find a sequence of points $(x_1, y_1), (x_2, y_2), \dots$ that approximate (x, y) .

- We will only find (x_1, y_1) as follow : $\begin{pmatrix} x_1 \\ y_1 \end{pmatrix} = \begin{pmatrix} x_0 \\ y_0 \end{pmatrix} - J^{-1}_{(x_0, y_0)} \begin{pmatrix} f(x_0, y_0) \\ g(x_0, y_0) \end{pmatrix},$ where

J is the Jacobian matrix given by

$$J(x, y) = \begin{pmatrix} f_x & f_y \\ g_x & g_y \end{pmatrix} \rightarrow J^{-1}(x, y) = \frac{1}{f_x g_y - f_y g_x} \begin{pmatrix} g_y & -f_y \\ -g_x & f_x \end{pmatrix}.$$

$$\text{To find } (x_2, y_2) = \begin{pmatrix} x_1 \\ y_1 \end{pmatrix} - J^{-1}(x_1, y_1) \begin{pmatrix} f(x_1, y_1) \\ g(x_1, y_1) \end{pmatrix}.$$

2 Fixed Point Iteration (F.P.I)

- This method can be used to solve 2×2 or 3×3 nonlinear systems.

- For 2×2 nonlinear system:

- Write $x = g_1(x, y)$... ①
 $y = g_2(x, y)$

- Given $(x_0, y_0) = (p_0, q_0)$

- The FPI is $p_{n+1} = g_1(p_n, q_n)$, $q_{n+1} = g_2(p_n, q_n)$.

- For 3×3 nonlinear system:

- Write $x = g_1(x, y, z)$
 $y = g_2(x, y, z)$... ②
 $z = g_3(x, y, z)$

- Given $(x_0, y_0, z_0) = (p_0, q_0, r_0)$.

- The FPI is $p_{n+1} = g_1(p_n, q_n, r_n)$, $q_{n+1} = g_2(p_n, q_n, r_n)$, $r_{n+1} = g_3(p_n, q_n, r_n)$.

Def :

- The point (p, q) is fixed point of the system ① if $p = g_1(p, q)$ and $q = g_2(p, q)$.

- The point (p, q, r) is fixed point of the system if $p = g_1(p, q, r)$ and $q = g_2(p, q, r)$ and $r = g_3(p, q, r)$.

Convergence of FPI - Two dimensions.

- Assume (P, Q) is fixed point of $x = g_1(x, y)$ and $y = g_2(x, y)$.
- If (P_0, Q_0) is sufficiently close to (P, Q) and if
 - A. $|g_{1x}(P, Q)| + |g_{1y}(P, Q)| < 1$ and
 - B. $|g_{2x}(P, Q)| + |g_{2y}(P, Q)| < 1$, Then the FPI convergence to the fixed point (P, Q) .

Remark :

1. Convergence of FPI for three dimensions follows similarly to Theorem above by adding Z -component.

2.

Div. g_i converge to λ (where $\lambda \neq 1$)

3 Seidel Iteration (improvement of FPI)

This method can be used to solve 2×2 or 3×3 nonlinear system.

For 2×2 nonlinear system :

- Write $x = g_1(x, y)$, $y = g_2(x, y)$.
- Given $(x_0, y_0) = (P_0, Q_0)$.
- The seidel Iteration is $P_{n+1} = g_1(P_n, Q_n)$, $Q_{n+1} = g_2(P_n, Q_n)$.

For 3×3 nonlinear system :

- Write $x = g_1(x, y, z)$, $y = g_2(x, y, z)$, $z = g_3(x, y, z)$.
- Given $(x_0, y_0, z_0) = (P_0, Q_0, R_0)$.
- The seidel Iteration is $P_{n+1} = g_1(P_n, Q_n, R_n)$,
 $Q_{n+1} = g_2(P_n, Q_n, R_n)$,
 $R_{n+1} = g_3(P_n, Q_n, R_n)$.

Chapter 4 :

- Interpolations means polynomial approximation.
- given a function $f(x)$ on $[a, b] = [x_0, x_n]$.
- given $n+1$ points $(x_0, f(x_0)), (x_1, f(x_1)), \dots, (x_n, f(x_n))$ through the partition
$$a = x_0 < x_1 < x_2 < \dots < x_n = b$$
- We need approximate $f(x)$ by a polynomial of order at most n passing through these $n+1$ points (nodes) on $[a, b]$.
- That is, $f(x) \approx P_n(x) + E_n(x)$ on $[a, b]$.
is called interpolation polynomial \leftarrow \rightarrow truncation error.

Th. Given $n+1$ points : $(x_0, y_0), (x_1, y_1), \dots, (x_n, y_n)$, Then \exists unique polynomial $P_n(x)$ with degree $\leq n$ that passes through these point.

Remark: If $f(x)$ is given and analytic at x_0 it has continuous derivatives of all derivatives of all orders and can be represented as Taylor series in an interval about x_0 , then we can use Taylor polynomial Approximation to estimate $f(x)$ by a Taylor polynomial.

Theorem : Taylor polynomial Approximation :

- Assume $f \in C^{n+1}[a,b]$ and $x_0 \in [a,b]$ is fixed.
- If $x \in [a,b]$, then $f(x) \approx P_n(x) + E_n(x)$ where $P_n(x)$ is the Taylor polynomial of degree n that estimates $f(x)$ on $[a,b]$ given by $P_n(x) = \sum_{k=0}^n \frac{f^{(k)}(x_0)}{k!} (x-x_0)^k$.

and $E_n(x)$ is the truncation error given by $E_n(x) = \frac{f^{(n+1)}(c)}{(n+1)!} (x-x_0)^{n+1}$.

→ However, f is usually not known or hard to compute. So How to find $P_n(x)$?

1. Lagrange Interpolation : $P_n(x)$ is the Lagrange polynomial.

2. Newton Interpolation : $P_n(x)$ is the Newton polynomial.

II Lagrange polynomial :

- Linear interpolation uses ~~a line~~ a line segment that passes through two points.
- In general, the Lagrange polynomial of degree at most n passes through $n+1$ points : $(x_0, y_0), (x_1, y_1), \dots, (x_n, y_n)$ has the form $P_n(x) = \sum_{k=0}^n y_k L_{n,k}(x)$.

$$L_{n,k} = \frac{(x-x_0)(x-x_1)\dots(x-x_{k-1})(x-x_{k+1})\dots(x-x_n)}{(x_k-x_0)(x_k-x_1)\dots(x_k-x_{k-1})(x_k-x_{k+1})\dots(x_k-x_n)}$$

Def : Uniform partition :

- The partition of $[a,b] = [x_0, x_n]$ is uniform if the nodes $x_0, x_1, x_2, \dots, x_n$ are equally spaced.
- That is, $x_k = x_0 + kK$ for $k=0, 1, 2, \dots, n$.

$$h = \frac{b-a}{n}$$

2] Newton's polynomial :

Th (Newton polynomial) :

- Given $x_0, x_1, x_2, \dots, x_n$, $n+1$ distinct numbers in $[a, b]$.
- Then, \exists a unique polynomial $P_n(x)$ (called Newton polynomial) of degree at most n s.t $f(x_i) = P(x_i)$ for $i=1, 2, \dots, n$.
- Furthermore, Newton polynomial is given by

$$P_n(x) = a_0 + a_1(x-x_0) + a_2(x-x_0)(x-x_1) + a_3(x-x_0)(x-x_1)(x-x_2) + \dots + a_n(x-x_0)(x-x_1)\dots(x-x_{n-1}).$$

Where the coefficients of Newton's polynomial are given by the divided difference

$$a_k = f[x_0, x_1, \dots, x_k] \text{ for } k=0, 1, \dots, n.$$

→ That is :

$$a_0 = f[x_0] = f(x_0) = y_0 : \text{Zero divided differences}.$$

$$a_1 = f[x_0, x_1] = \frac{f(x_1) - f(x_0)}{x_1 - x_0} = \frac{y_1 - y_0}{x_1 - x_0} : \text{First divided differences}$$

$$a_2 = f[x_0, x_1, x_2] = \frac{f[x_1, x_2] - f[x_0, x_1]}{x_2 - x_0} = \frac{\frac{y_2 - y_1}{x_2 - x_1} - \frac{y_1 - y_0}{x_1 - x_0}}{x_2 - x_0} : \text{Second divided differences}.$$

* Divided Difference Table for $y=f(x)$:

x_k	$f(x_k) = y_k$	1^{st} D.D	2^{nd} D.D	3^{rd} D.D
x_0	$y_0 = a_0$			
x_1	y_1	$f[x_0, x_1] = a_1$		
x_2	y_2	$f[x_1, x_2]$	$f[x_0, x_1, x_2] = a_2$	
x_3	y_3	$f[x_2, x_3]$	$f[x_1, x_2, x_3]$	$f[x_0, x_1, x_2, x_3] = a_3$
x_4	y_4	$f[x_3, x_4]$	$f[x_2, x_3, x_4]$	$f[x_1, x_2, x_3, x_4]$

↳ Error terms and error bounds

* Error terms $E_n(x)$:

$$E_n(x) = f(x_k) - P_n(x_k) = y_k - p_k$$

Theorem ($E_n(x)$):

Assume $f \in C^{n+1}[a,b]$ and $x_0, x_1, \dots, x_n \in [a,b]$ are $n+1$ nodes.

Then $f(x) = P_n(x) + E_n(x)$ where $E_n(x)$ is the error term given by

$$E_n(x) = (x-x_0)(x-x_1)(x-x_2)\dots(x-x_n) \frac{f^{(n+1)}(c)}{(n+1)!} \text{ for some } c = c(x) \text{ lies in } [a,b].$$

$$\rightarrow P_n(x) = y_0 L_{n,0}(x) + y_1 L_{n,1}(x) + \dots + y_n L_{n,n}(x) \quad (\text{Lagrange polynomial})$$

$$\rightarrow P_n(x) = q_0 + q_1(x-x_0) + q_2(x-x_0)(x-x_1) + \dots + q_n(x-x_0)(x-x_1)\dots(x-x_{n-1}) \quad (\text{Newton polynomial})$$

How to find an upper bound for the error term $E_n(x)$?

- That is, we need to find some constant s.t. $|E_n(x)| \leq \text{constant}$.

- Finding the upper bound depends on whether the nodes are equally spaced (uniform partition) or not (not uniform partition).

* Th 1 (upper bound of the error term for Interpolation - uniform partition).

$n=1$ ① Given $(x_0, y_0), (x_1, y_1), (n=1)$ with $E_1(x) = \frac{(x-x_0)(x-x_1)}{2!} \hat{f}(c)$, Then

$$|E_1(x)| \leq \frac{h^2 M_2}{8} \text{ where } M_2 = \max |f'(x)|. \quad h = x_1 - x_0$$

$n=2$ ② Given $(x_0, y_0), (x_1, y_1), (x_2, y_2)$ with $E_2(x) = \frac{(x-x_0)(x-x_1)(x-x_2)}{3!} \hat{f}(c)$ Then

$$|E_2(x)| \leq \frac{h^3 M_3}{9\sqrt{3}} \text{ where } M_3 = \max |f''(x)|. \quad h = \frac{x_2 - x_0}{2}$$

$n=3$ ③ Given $(x_0, y_0), (x_1, y_1), (x_2, y_2), (x_3, y_3)$ with $E_3(x) = \frac{(x-x_0)(x-x_1)(x-x_2)(x-x_3)}{4!} f^{(4)}(c)$

$$\text{Then } |E_3(x)| \leq \frac{h^4 M_4}{24} \text{ where } M_4 = \max |f^{(4)}(x)|. \quad h = \frac{x_3 - x_0}{3}$$

CURVE FITTING.

* Upper bound of the error term for interpolation - Non uniform partition.

Given $(x_0, y_0), (x_1, y_1), (x_2, y_2)$ with $E_2(x) = \frac{(x-x_0)(x-x_1)(x-x_2)}{3!} \tilde{f}(c)$ Then

$$|E_2(x)| \leq \frac{|\phi_2(x)| |\tilde{f}(c)|}{6} \text{ where } .$$

$|\phi_2(x)|$ is an upper bound of $\phi_2(x) = (x-x_0)(x-x_1)(x-x_2)$ and

$|\tilde{f}(c)|$ is an upper bound of $|\tilde{f}(c)|$.

Chapter 4 done.

Chapter 5 : Curve fitting.

- To handle the errors, we use norms to measure how far the curve $y=f(x)$ lies from the data.

$$|e_k| = |f(x_k) - j_k| , \quad k=1, 2, \dots, n$$

- We consider the following norms :

1. Maximum error : $E_\infty(f) = \max |e_k| , \quad k \in \{1, 2, \dots, n\}$.

2. Average error : $E_1(f) = \frac{1}{n} \sum_{k=1}^n |e_k|$.

3. Root-mean-square error : $E_2(f) = \sqrt{\frac{1}{n} \sum_{k=1}^n |e_k|^2}$.

* Finding the Least-squares line :

- Given n distinct points : $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)$.

- The least square line $y = f(x) = Ax + B$ is the line that minimizes the RMSE, $E_2(f)$.

$$E_2(f) = \sqrt{\frac{\sum (f(x_i) - j_i)^2}{n}} \Rightarrow n E_2^2 = \sum_{i=1}^n (f(x_i) - j_i)^2$$

- E_2 is minimized iff $E(A, B) = \sum_{i=1}^n (Ax_i + B - y_i)^2$ is minimized.

$$\rightarrow \frac{\partial E}{\partial A} = 2 \sum_{i=1}^n (Ax_i + B - y_i) x_i = 0 \Leftrightarrow \sum_{i=1}^n (Ax_i^2 + Bx_i - y_i x_i) = 0$$

$$A \sum_{i=1}^n x_i^2 + B \sum_{i=1}^n x_i = \sum_{i=1}^n x_i y_i \quad \text{... (1)}$$

$$\rightarrow \frac{\partial E}{\partial B} = 2 \sum_{i=1}^n (Ax_i + B - y_i) = 0 \Leftrightarrow \sum_{i=1}^n (Ax_i + B - y_i) = 0$$

$$A \sum_{i=1}^n x_i + nB = \sum_{i=1}^n y_i \quad \text{... (2)}$$

equation (1) and (2) are called the linear normal equations and used to find the coefficients A and B .

5.3 → ملخص المنهج

Chapter 6 : Numerical Differentiation.

* Notation :

- $f_k = f(x_0 + kh)$, $k=0, \pm 1, \pm 2, \pm 3$.

- That is, $f_0 = f(x_0)$

$$f_1 = f(x_0 + h)$$

$$f_{-1} = f(x_0 - h)$$

$$f_2 = f(x_0 + 2h)$$

* Difference Formula for $\bar{f}(x_0)$

- C.D.F of order $O(h^2)$: $\bar{f}(x_0) \approx \frac{f(x_0 + h) - f(x_0 - h)}{2h} + E_{h^2}(h)$

$$\bar{f}(x_0) \approx \frac{f_1 - f_{-1}}{2h} + \frac{-h^2 \bar{F}(c)}{6}$$

- F.D.F of order $O(h^2)$: $\bar{f}(x_0) \approx \frac{-3f_0 + 4f_1 - f_2}{2h} + \frac{h^2 \bar{F}(c)}{3}$

- B.D.F of order $O(h^2)$: $\bar{f}(x_0) \approx \frac{3f_0 - 4f_{-1} + f_{-2}}{2h} + \frac{h^2 \bar{F}(c)}{3}$

- C.D.F of order $O(h^4)$: $\bar{f}(x_0) \approx \frac{-f_2 + 8f_1 - 8f_{-1} + f_{-2}}{12h} + \frac{h^4 \bar{F}^{(4)}(c)}{30}$

* Difference Formula for $\bar{\bar{f}}(x_0)$:

- C.D.F of order $O(h^2)$: $\bar{\bar{f}}(x_0) \approx \frac{f_1 - 2f_0 + f_{-1}}{h^2} + \frac{-h^2 \bar{f}^{(4)}(c)}{12}$

- F.D.F of order $O(h^2)$: $\bar{\bar{f}}(x_0) \approx \frac{2f_0 - 5f_1 + 4f_2 - f_3}{h^2} + \frac{11h^2 \bar{f}^{(4)}(c)}{12}$

- B.D.F of order $O(h^2)$: $\bar{\bar{f}}(x_0) \approx \frac{2f_0 - 5f_{-1} + 4f_{-2} - f_{-3}}{h^2} + \frac{11h^2 \bar{f}^{(4)}(c)}{12}$

- C.D.F of order $O(h^4)$: $\bar{\bar{f}}(x_0) \approx \frac{-f_2 + 16f_1 - 30f_0 + 16f_{-1} - f_{-2}}{12h^2} + \frac{h^4 \bar{f}^{(4)}(c)}{90}$

6.2 : Derivation of D.F's .

- Recall Taylor expansion of $f(x)$ about x_0 :

$$f(x) = f(x_0) + \bar{f}(x_0)(x-x_0) + \frac{\bar{f}'(x_0)}{2!} (x-x_0)^2 + \frac{\bar{f}''(x_0)}{3!} (x-x_0)^3 + \dots$$

- clearly \rightarrow when $x = x_0 + h \rightarrow x - x_0 = h$:

$$f_1 = f(x_0+h) = f(x_0) + h \bar{f}(x_0) + \frac{h^2}{2!} \bar{f}'(x_0) + \frac{h^3}{3!} \bar{f}''(x_0) + \dots$$

$$f_{-1} = f(x_0-h) = f(x_0) - h \bar{f}(x_0) + \frac{h^2}{2!} \bar{f}'(x_0) - \frac{h^3}{3!} \bar{f}''(x_0) + \dots$$

$$f_2 = f(x_0+2h) = f(x_0) + 2h \bar{f}(x_0) + \frac{(2h)^2}{2!} \bar{f}'(x_0) + \frac{(2h)^3}{3!} \bar{f}''(x_0) + \dots$$

$$f_{-2} = f(x_0-2h) = f(x_0) - 2h \bar{f}(x_0) + \frac{(2h)^2}{2!} \bar{f}'(x_0) - \frac{(2h)^3}{3!} \bar{f}''(x_0) + \dots$$

$$f_k = f(x_0+kh) = f_0 + kh \bar{f}(x_0) + \frac{(kh)^2}{2!} \bar{f}'(x_0) + \frac{(kh)^3}{3!} \bar{f}''(x_0) + \dots$$

ch. 6 done.

لطفاً

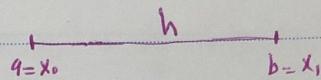
Ch 7 :

7.1 : Numerical Integration :

* closed Newton - Cotes quadrature formula:

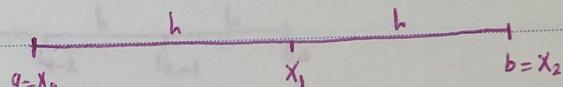
Th: Assume that $x_k = x_0 + kh$ are equally spaced nodes with $f_k = f(x_k)$, Then

[a] Trapezoidal Rule is :



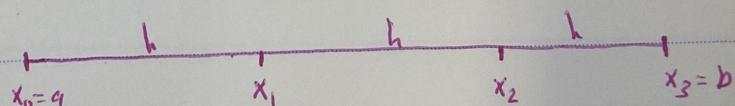
$$\int_a^b f(x) dx = \int_{x_0}^{x_1} f(x) dx \approx \frac{h}{2} (f_0 + f_1) \text{ with error } \frac{-h^3 f''(c)}{12}$$

[b] Simpson's $\frac{1}{3}$ Rule is :



$$\int_a^b f(x) dx = \int_{x_0}^{x_2} f(x) dx \approx \frac{h}{3} (f_0 + 4f_1 + f_2) \text{ with error } \frac{-h^5 f'''(c)}{90}$$

[c] Simpson's $\frac{3}{8}$ Rule is :



$$\int_a^b f(x) dx = \int_{x_0}^{x_3} f(x) dx \approx \frac{3h}{8} (f_0 + 3f_1 + 3f_2 + f_3) \text{ with error } \frac{-3h^5 f'''(c)}{80}$$

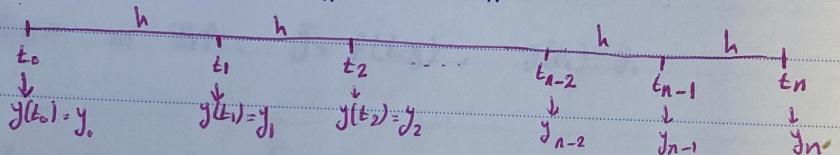
ch9: Numerical Approximation for IVP:

- Given the IVP: $\frac{dy}{dt} = f(t, y(t))$ with $y(t_0) = y_0$.
- To estimate the solution $y(t)$ numerically:
 - We find the values of y at different values of t .
 - Then approximate y using interpolation.

* General principle :

- Given an IVP: $\dot{y} = f(t, y)$, $y(t_0) = y_0$.
- To estimate the values of y on $[a, b] = [t_0, t_n]$:

$$\rightarrow \text{Find } h = \frac{t_n - t_0}{n} = \frac{b-a}{n}$$



$$\rightarrow \text{Find } y(t_k) = y_k, \text{ where } k=0, 1, \dots, n \text{ and } t_k = t_0 + kh.$$

- To find y_1, y_2, \dots .

II Euler's Method :

- Given an IVP: $\dot{y} = f(t, y)$, $y(t_0) = y_0$.
- $h = \frac{t_n - t_0}{n} = \frac{b-a}{n}$
- $y_1 = y_0 + h f(t_0, y_0)$.
- $y_2 = y_1 + h f(t_1, y_1)$.

$$\text{in general : } y_n = y_{n-1} + h f(t_{n-1}, y_{n-1})$$

② Taylor's Method of order 2 :

- Given an IVP: $\dot{y} = f(t, y)$, $y(t_0) = y_0$. with $[a, b] = [t_0, t_n]$, $h = \frac{b-a}{n}$

- $y_1 = y_0 + h f(t_0, y_0) + \frac{h^2}{2} \bar{f}(t_0, y_0)$.

- in general : $y_{k+1} = y_k + h f(t_k, y_k) + \frac{h^2}{2} \bar{f}(t_k, y_k)$, $k=0, 1, 2, \dots, n-1$.

③ Heun's Method :

- Given an IVP: $\dot{y} = f(t, y)$, $y(t_0) = y_0$.

- $y_1 = y_0 + \frac{h}{2} [f(t_0, y_0) + f(t_1, y_0 + hf(t_0, y_0))]$.

- $y_2 = y_0 + \frac{h}{2} [f(t_1, y_1) + f(t_2, y_1 + hf(t_1, y_1))]$,

ch. 9 done.