

Distance Between Words for spellchecking





Errors are close to correct words

Need a measure of distance between strings (words/phrases):

- Jackard coefficient/similarity (done)
- Edit distance: effort to convert one into another
- Soundex: phonetic closeness (sound: تشابه النطق)
- More likely to mistake with adjacent keyboard characters, or without language switch: Really ^{١٨}نشمغ really in Arabic keyboard!
- Confusion letters: {ء ي و ا آ إ}, {ى ي}, {ة ه} less distance within
- **Final distance: A weighted sum!**



How similar are two strings?

- Spell correction

- The user typed “graffe”

Which is closest?

- graf
 - graft
 - grail
 - giraffe

- Computational Biology

- Align two sequences of nucleotides

```

AGGCTATCACCTGACCTCCAGGCCGATGCCC
TAGCTATCACGACCGCGGTCGATTGCCCCGAC
    
```

- Resulting alignment:

```

-AGGCTATCACCTGACCTCCAGGCCGA--TGCCC---
TAG-CTATCAC--GACCGC--GGTCGATTGCCCCGAC
    
```

- Also for Machine Translation, Information Extraction, Speech Recognition



Edit Distance

- The minimum edit distance between two strings
- Is the minimum number of editing operations
 - Insertion
 - Deletion
 - Substitution
- Needed to transform one into the other



Minimum Edit Distance

- Two strings and their **alignment**:

I	N	T	E	*	N	T	I	O	N
*	E	X	E	C	U	T	I	O	N



Minimum Edit Distance

I	N	T	E	*	N	T	I	O	N
*	E	X	E	C	U	T	I	O	N
d	s	s		i	s				

- If each operation has cost of 1
 - Distance between these is 5
- If substitutions cost 2 (Levenshtein)
 - Distance between them is 8



Alignment in Computational Biology

- Given a sequence of bases

AGGCTATCACCTGACCTCCAGGCCGATGCCC
TAGCTATCACGACCGCGGGTCGATTGCCCCGAC

- An alignment:

–AGGCTATCACCTGACCTCCAGGCCGA–TGCCC––
TAG–CTATCAC–GACCGC–GGTCGATTGCCCCGAC

- Given two sequences, align each letter to a letter or gap



Other uses of Edit Distance in NLP

- Evaluating Machine Translation and speech recognition

R Spokesman confirms senior government adviser was shot

H Spokesman said the senior adviser was shot dead

S

I

D

I

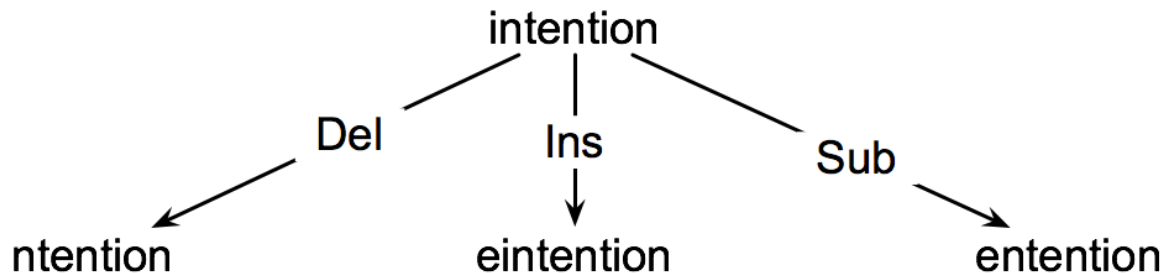
- Named Entity Extraction and Entity Coreference

- IBM Inc. announced today
- IBM profits
- Stanford President John Hennessy announced yesterday
- for Stanford University President John Hennessy



How to find the Min Edit Distance?

- Searching for a path (sequence of edits) from the start string to the final string:
 - **Initial state:** the word we're transforming
 - **Operators:** insert, delete, substitute [Levenshtein] exchange[Demerau]
 - **Goal state:** the word we're trying to get to
 - **Path cost:** what we want to minimize: the number of edits





Minimum Edit as Search

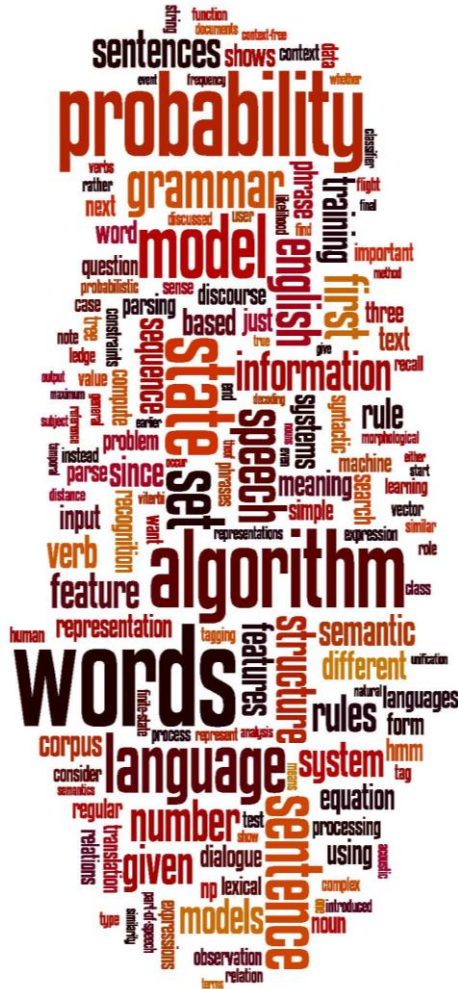
- But the space of all edit sequences is huge!
 - We can't afford to navigate naïvely
 - Lots of distinct paths wind up at the same state.
 - We don't have to keep track of all of them
 - Just the shortest path to each of those revisited states.



Defining Min Edit Distance

- For two strings
 - X of length n
 - Y of length m
- We define $D(i,j)$
 - the edit distance between $X[1..i]$ and $Y[1..j]$
 - i.e., the first i characters of X and the first j characters of Y
 - The edit distance between X and Y is thus $D(n,m)$

Minimum Edit Distance





Dynamic Programming for Minimum Edit Distance

- **Dynamic programming:** A tabular computation of $D(n, m)$
- Solving problems by combining solutions to subproblems.
- Bottom-up
 - We compute $D(i, j)$ for small i, j
 - And compute larger $D(i, j)$ based on previously computed smaller values
 - i.e., compute $D(i, j)$ for all $i (0 < i < n)$ and $j (0 < j < m)$



Defining Min Edit Distance (Levenshtein)

- Initialization

$$D(i, 0) = i$$

$$D(0, j) = j$$

- Recurrence Relation:

For each $i = 1 \dots M$

For each $j = 1 \dots N$

$$D(i, j) = \min \begin{cases} D(i-1, j) + 1 \\ D(i, j-1) + 1 \\ D(i-1, j-1) + \begin{cases} 2; & \text{if } X(i) \neq Y(j) \\ 0; & \text{if } X(i) = Y(j) \end{cases} \end{cases}$$

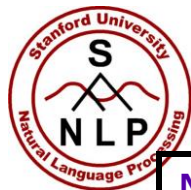
- Termination:

$D(N, M)$ is distance



The Edit Distance Table

N	9									
O	8									
I	7									
T	6									
N	5									
E	4									
T	3									
N	2									
I	1									
#	0	1	2	3	4	5	6	7	8	9
	#	E	X	E	C	U	T	I	O	N



The Edit Distance Table

N	9									
O	8									
I	7									
T	6									
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N	2									
I	1									
#	0	1	2	3	4	5	6	7	8	9
	#	E	X	E	C	U	T	I	O	N

$$D(i,j) = \min \begin{cases} D(i-1,j) + 1 \\ D(i,j-1) + 1 \\ D(i-1,j-1) + \begin{cases} 2; & \text{if } S_1(i) \neq S_2(j) \\ 0; & \text{if } S_1(i) = S_2(j) \end{cases} \end{cases}$$



Edit Distance

$$D(i,j) = \min \begin{cases} D(i-1,j) + 1 \\ D(i,j-1) + 1 \\ D(i-1,j-1) + \begin{cases} 2; & \text{if } S_1(i) \neq S_2(j) \\ 0; & \text{if } S_1(i) = S_2(j) \end{cases} \end{cases}$$

N	9									
O	8									
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I	1									
#	0	1	2	3	4	5	6	7	8	9
	#	E	X	E	C	U	T	I	O	N



The Edit Distance Table

N	9	8	9	10	11	12	11	10	9	8
O	8	7	8	9	10	11	10	9	8	9
I	7	6	7	8	9	10	9	8	9	10
T	6	5	6	7	8	9	8	9	10	11
N	5	4	5	6	7	8	9	10	11	10
E	4	3	4	5	6	7	8	9	10	9
T	3	4	5	6	7	8	7	8	9	8
N	2	3	4	5	6	7	8	7	8	7
I	1	2	3	4	5	6	7	6	7	8
#	0	1	2	3	4	5	6	7	8	9
	#	E	X	E	C	U	T	I	O	N

Minimum Edit Distance

Weighted Minimum Edit Distance





Levenshtein vs Demerau

- Levenshtein: three ops: add, delete, substitute.
- Demerau: Add, delete, substitute, exchange:
 - Resaerch → research 1 Demerau, 2 or more Levenshtein



Weighted Edit Distance

- Why would we add weights to the computation?
 - Spell Correction: some letters are more likely to be mistyped than others
 - Biology: certain kinds of deletions or insertions are more likely than others



Confusion matrix for spelling errors

sub[X, Y] = Substitution of X (incorrect) for Y (correct)

X	Y (correct)																									
	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
a	0	0	7	1	342	0	0	2	118	0	1	0	0	3	76	0	0	1	35	9	9	0	1	0	5	0
b	0	0	9	9	2	2	3	1	0	0	0	5	11	5	0	10	0	0	2	1	0	0	8	0	0	0
c	6	5	0	16	0	9	5	0	0	0	1	0	7	9	1	10	2	5	39	40	1	3	7	1	1	0
d	1	10	13	0	12	0	5	5	0	0	2	3	7	3	0	1	0	43	30	22	0	0	4	0	2	0
e	388	0	3	11	0	2	2	0	89	0	0	3	0	5	93	0	0	14	12	6	15	0	1	0	18	0
f	0	15	0	3	1	0	5	2	0	0	0	3	4	1	0	0	0	6	4	12	0	0	2	0	0	0
g	4	1	11	11	9	2	0	0	0	1	1	3	0	0	2	1	3	5	13	21	0	0	1	0	3	0
h	1	8	0	3	0	0	0	0	0	0	2	0	12	14	2	3	0	3	1	11	0	0	2	0	0	0
i	103	0	0	0	146	0	1	0	0	0	0	6	0	0	49	0	0	0	2	1	47	0	2	1	15	0
j	0	1	1	9	0	0	1	0	0	0	0	2	1	0	0	0	0	0	5	0	0	0	0	0	0	0
k	1	2	8	4	1	1	2	5	0	0	0	0	5	0	2	0	0	0	6	0	0	0	4	0	0	3
l	2	10	1	4	0	4	5	6	13	0	1	0	0	14	2	5	0	11	10	2	0	0	0	0	0	0
m	1	3	7	8	0	2	0	6	0	0	4	4	0	180	0	6	0	0	9	15	13	3	2	2	3	0
n	2	7	6	5	3	0	1	19	1	0	4	35	78	0	0	7	0	28	5	7	0	0	1	2	0	2
o	91	1	1	3	116	0	0	0	25	0	2	0	0	0	0	14	0	2	4	14	39	0	0	0	18	0
p	0	11	1	2	0	6	5	0	2	9	0	2	7	6	15	0	0	1	3	6	0	4	1	0	0	0
q	0	0	1	0	0	0	27	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
r	0	14	0	30	12	2	2	8	2	0	5	8	4	20	1	14	0	0	12	22	4	0	0	1	0	0
s	11	8	27	33	35	4	0	1	0	1	0	27	0	6	1	7	0	14	0	15	0	0	5	3	20	1
t	3	4	9	42	7	5	19	5	0	1	0	14	9	5	5	6	0	11	37	0	0	2	19	0	7	6
u	20	0	0	0	44	0	0	0	64	0	0	0	0	2	43	0	0	4	0	0	0	0	2	0	8	0
v	0	0	7	0	0	3	0	0	0	0	0	1	0	0	1	0	0	0	8	3	0	0	0	0	0	0
w	2	2	1	0	1	0	0	2	0	0	1	0	0	0	0	7	0	6	3	3	1	0	0	0	0	0
x	0	0	0	2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	9	0	0	0	0	0	0	0
y	0	0	2	0	15	0	1	7	15	0	0	0	2	0	6	1	0	7	36	8	5	0	0	1	0	0
z	0	0	0	7	0	0	0	0	0	0	0	7	5	0	0	0	0	2	21	3	0	0	0	0	3	0



Weighted Min Edit Distance

- Initialization:

$$D(0, 0) = 0$$

$$D(i, 0) = D(i-1, 0) + \text{del}[x(i)]; \quad 1 < i \leq N$$

$$D(0, j) = D(0, j-1) + \text{ins}[y(j)]; \quad 1 < j \leq M$$

- Recurrence Relation:

$$D(i, j) = \min \begin{cases} D(i-1, j) & + \text{del}[x(i)] \\ D(i, j-1) & + \text{ins}[y(j)] \\ D(i-1, j-1) & + \text{sub}[x(i), y(j)] \end{cases}$$

- Termination:

$D(N, M)$ is distance



Keyboarding Errors: more for mobiles





Where did the name, dynamic programming, come from?

...The 1950s were not good years for mathematical research. [the] Secretary of Defense ...had a pathological fear and hatred of the word, research...

I decided therefore to use the word, “**programming**”.

I wanted to get across the idea that this was dynamic, this was multistage... I thought, let's ... take a word that has an absolutely precise meaning, namely **dynamic**... it's impossible to use the word, **dynamic**, in a pejorative sense. Try thinking of some combination that will possibly give it a pejorative meaning. It's impossible.

Thus, I thought dynamic programming was a good name. It was something not even a Congressman could object to.”



Soundex Algorithm

Idea

- –Vowels are viewed as interchangeable in transcribing names
- Consonants with similar sounds (e.g., D and T) are put in equivalence classes
- related names often have the same soundex codes

Algorithm

- –Turn every term to be indexed into a four-character reduced form,
- build an inverted index from these reduced forms to the original terms called the soundex index
- Do the same with the query terms
- Search the soundex index



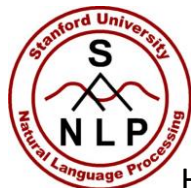
Soundex Algorithm: Four-Character Code

- The first character is a letter of the alphabet and the other three are digits between 0 and 9
- Algorithm
 - Retain the first letter of the term
 - Change all occurrences of the following letters to '0'
A, E, I, O U, H, W, and Y
 - Change letters to digits as follows
B, F, P, V $\rightarrow$ 1



Soundex Algorithm: Four-Character Code

- D, T $\rightarrow$ 3
- L $\rightarrow$ 4
- M, N $\rightarrow$ 5
- R $\rightarrow$ 6
- –Repeatedly remove one out of each pair of consecutive identical digits
- –Remove all 0's from the resulting string, pad the resulting string with trailing zeros and return the first four positions: a letter followed by three digits
- **Example:** Hermann $\rightarrow$ H655, Herman $\rightarrow$ H655, matched!



Arabic Soundex : Four-Character Code

Hold the first letter.

- Replace the characters (ا, أ, إ, آ, ح, ع, غ, ش, و, ي) with the value 0.
- Replace the characters (ب, ف) with the value 1.
- Replace the characters (ك, خ, ج, ز, س, ص, ظ, ق, ك) with the value 2.
- Replace the characters (ت, ث, د, ذ, ض, ط) with the value 3.
- Replace the character (ل) with the value 4.
- Replace the characters (م, ن) with the value 5.
- Replace the character (ر) with the value 6.
- <https://www.codeproject.com/Articles/26880/Arabic-Soundex>

Soundex Code	المحارف	المجموعة
1	b, f, p, v	شَقَهِيّ
2	c, g, j, k, q, s, x, z	حروف حَلَقِيّة و حُرُوف صَفِير
3	d, t	حرف نَطْعِي (ملفوظ بوضوح اللسان على مؤخر الاسنان الامامية العليا)
4	l	حرف صامت ملفوظ بلطف طَوِيل
5	m, n	حرف يلفظ من الانف
6	r	حرف صامت ملفوظ بلطف فَصِير

Form1

Word 1: Start

Word 2:

Select a word

id	Name	Value
1	أحمد	x530
2	أحمد	x530
3	أحمد	x530
4	أحمدان	x535
5	تمام	x550
6	تمام	x550
7	أحمدان	x535
8	تمام	x550
9	جامعة	x550
10	أحمدان	x354
11	أحمدان	x344
12	أحمدان	x354
13	أحمدان	x543
14	أحمدان	x540
15	تامر	x560
16	تامر	x560
17	محمدي	x540
18	محمدي	x540
19	عربي	x610
20	عربي	x610
21	أحمد	x530

Similar Words

id	Name	Value
1	أحمد	x530
2	أحمد	x530
3	أحمد	x530
21	أحمد	x530

☒ Arabic ☐ English

Create Sound List

Insert



Errors are close to correct words

Distance between error and correct words: a combination of:

- Jackard coefficient/similarity (distance)
- Edit distance (Levenshtein, Demerau): *not all equal*
- Soundex: phonetic closeness (sound: تشابه النطق): *less distance here*
- More likely to mistake with adjacent keyboard characters, or without language switch: Really ^ٲنشمغ really in Arabic keyboard!
- Confusion letters: {ء ي و ا آ إ}, {ى ي}, {ة ه} less distance within
- **Final distance: A weighted sum!**

