



ANSWER BOOKLET

Student: <u>Digital</u>	Number: <u>3</u>
Course: Department:	Number:
Division:	Instructor:
Date: Day Month Year	

Dr. Abdellatif Abu Issa

For Instructor's Use

Question	Grade
1	
2	
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Total	

Examples of Algebraic Manipulation

$$\textcircled{1} x(x' + y) = xx' + xy = 0 + xy = xy$$

$$\textcircled{2} x + x'y = (x + x'y)(x + y) = 1(x + y) = x + y$$

$$\textcircled{3} (x + y)(x + y')$$

$$= x + xy + xy' + yy' = x(1 + y + y') = x$$

$$\begin{aligned}\textcircled{4} xy + x'z + yz &= xy + x'z + yz(x + x') \\ &= xy + x'z + xyz + x'y z \\ &= xy(1 + z) + x'z(1 + y) \\ &= xy + x'z\end{aligned}$$

$$\textcircled{5} (x + y)(x' + z)(y + z) = (x + y)(x' + z)$$

$$= (xx' + xz + x'y + yz)(y + z)$$

$$= xyz + xz + x'y + x'yz + yz$$

$$= yz(x + x') + xz + x'y + yz$$

$$= yz + xz + x'y$$

$$= yz + x'y$$

$$= (x + y)(x' + z)$$

$$~~(x + y)(y + z)(x + z)(z + z) + x'y~~$$

$$~~x + x'~~$$

$$= yz(x + x') + xz + x'y =$$

$$xyz + x'yz + xz + x'y = xz(y + 1) + x'y(z + 1)$$

$$= xz + x'y = (x + x')(x + y)(x' + z)$$

$$~~(x + y)(x' + z)~~$$

⑧ Complement of a function

Demorgan's theorems :-

$$(A+B+C+\dots+Z)' = A'B'C'D'\dots Z'$$

$$(ABCDE\dots Z)' = A'+B'+C'+D'+E'+\dots+Z'$$

Example

$$F = x'yz' + x'y'z$$

$$\begin{aligned}\rightarrow F' &= (x'yz' + x'y'z)' \\ &= (x+y'+z)(x+y+z')\end{aligned}$$

Ex

$$F = x(x'z' + yz)$$

$$\begin{aligned}F' &= (x(x'z' + yz))' \\ &= x' + (y'z' + yz)' \\ &= x' + (y'z')' \cdot (yz)' \\ &= x' + (y+z)(y'+z')\end{aligned}$$

2-5

⑧ Canonical and Standard Forms

- a binary variable (x) may ~~be~~ appear in its normal form (x) or in its complement form (x').
- for two binary variable (x, y) ^{combined with AND operation} \rightarrow there are 2^2 forms: $x'y'$, $x'y$, xy' , xy .
- these terms are called minterms or a standard product.
- for n variable $\rightarrow 2^n$ forms (minterms).
- for two binary variables (x, y), Combined with OR operation \rightarrow there are 2^n forms: $x+y$, $x+y'$, $x'+y$, $x'+y'$.
- these terms are called maxterms or standard sum.

		<u>Minterm</u>		<u>Maxterm</u>	
x	y	Term	Designation	Term	Designation
0	0	$x'y'$	m_0	$x+y$	M_0
0	1	$x'y$	m_1	$x+y'$	M_1
1	0	xy'	m_2	$x'+y$	M_2
1	1	xy	m_3	$x'+y'$	M_3

① sum of all minterms (for 2 variables x, y)

$$\Rightarrow x'y' + x'y + xy' + xy = 1 \text{ (always)}$$

② product of all maxterms (for 2 variables x, y)

$$\Rightarrow (x+y) \cdot (x+y') \cdot (x'+y) \cdot (x'+y') = 0 \text{ (always)}$$

Ex.

x	y	z	minterms		maxterms	f ₁	f ₂
			f ₁	f ₂			
0	0	0	$x'y'z' (m_0)$		$x+y+z$	0	0
0	0	1	$x'y'z (m_1)$		$x+y+z'$	1	0
0	1	0	$x'y z' (m_2)$		$x'+y+z$	0	0
0	1	1	$x'y z (m_3)$		$x'+y+z'$	0	1
1	0	0	$x y' z' (m_4)$		$x'+y+z$	1	0
1	0	1	$x y' z (m_5)$		$x'+y+z'$	0	1
1	1	0	$x y z' (m_6)$		$x'+y'+z$	0	1
1	1	1	$x y z (m_7)$		$x'+y'+z'$	1	1

express f₁ in term of minterms

$$\Rightarrow f_1 = x'y'z + xy'z' + xyz$$

$$= m_1 + m_4 + m_7 = \sum (1, 4, 7)$$

in the same way f₂ = m₃ + m₅ + m₆ + m₇

$$= \sum (3, 5, 6, 7)$$

express f₁ in term of maxterms

$$f_1 = (x+y+z) \cdot (x+y'+z) \cdot (x'+y+z')$$

$$= M_0 \cdot M_2 \cdot M_3 \cdot M_5 \cdot M_6 = \prod (0, 2, 3, 5, 6)$$

① Any Boolean function can be expressed as a sum of ~~its~~ minterms or as a product of max terms.

Ex. Express $f = A + B'c$ in a sum of minterms

$$\begin{aligned}
 \Rightarrow f &= A(B+B') + B'c(A+A') \\
 &= AB + AB' + AB'c + A'B'c \\
 &= AB(c+c') + AB'(c+c') + AB'c + A'B'c \\
 &= \underline{ABC} + ABC' + AB'c + AB'c' + \underline{AB'c} + A'B'c \\
 &= m_7 + m_6 + m_5 + m_4 + m_1
 \end{aligned}$$

$$\rightarrow F(A, B, c) = \sum (1, 4, 5, 6, 7)$$

Ex. Express $f = xy + x'z$ in a product of maxterms
—change it into OR terms using distributive law

$$\begin{aligned}
 \Rightarrow f &= xy + x'z \\
 &= (xy + x')(xy + z) \\
 &= (x + x')(y + x')(x + z)(y + z) \\
 &= (x' + y)(x + z)(y + z) \\
 &= (x' + y + zz')(x + z + yy')(y + z + xx') \\
 &= (x' + y + z)(x' + y + z')(x + z + y)(x + z + y') \\
 &\quad (y + z + x)(y + z + x') \\
 &= (x' + y + z)(x' + y + z')(x + y + z)(x + y' + z) \\
 &= M_4 \cdot M_5 \cdot M_0 \cdot M_2 = \prod (0, 2, 4, 5)
 \end{aligned}$$

④ Conversion between Canonical forms

- complement of a function:

e.g. $F(A, B, C) = \sum(1, 4, 5, 6, 7)$

$$\Rightarrow F'(A, B, C) = \sum(0, 2, 3) = m_0 + m_2 + m_3$$

- The complement of a function expressed as the sum of minterms equals the sum of minterms missing from the original function.

* Use De Morgan's theorem \Rightarrow

$$\begin{aligned} F'(A, B, C) &= (m_0 + m_2 + m_3)' \\ &= m_0' \cdot m_2' \cdot m_3' = M_0 \cdot M_2 \cdot M_3 \\ &= \prod(0, 2, 3) \end{aligned}$$

$m_j' = M_j$

- Using truth table instead of algebraic theorems ($F = xy + x'z$)

x	y	z	x'	xy	x'z	xy + x'z	
0	0	0	1	0	0	0	
0	0	1	1	0	1	1	$\Rightarrow F = m_1 + m_3 + m_6 + m_7$
0	1	0	1	0	0	0	$= \sum(1, 3, 6, 7)$
0	1	1	1	0	1	1	
1	0	0	0	0	0	0	$\Rightarrow F = \prod(0, 2, 4, 5)$
1	0	1	0	0	0	0	
1	1	0	0	1	0	1	(Same result as using algebra)
1	1	1	0	1	0	1	

* Standard forms

- two types of standard forms: sum of products and products of sums

- sum of products

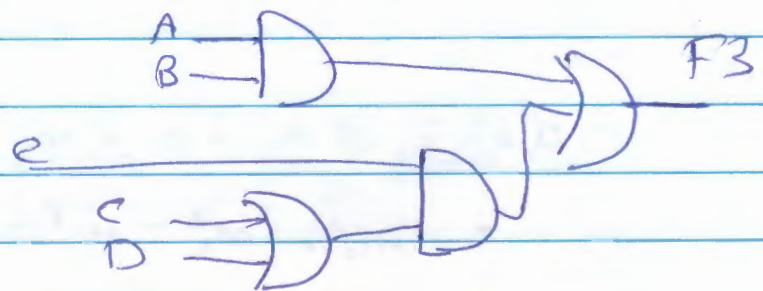
$$F1 = y' + xy + x'y'z'$$

- products of sums

$$F2 = x(y' + z)(x' + y + z')$$

- nonstandard form

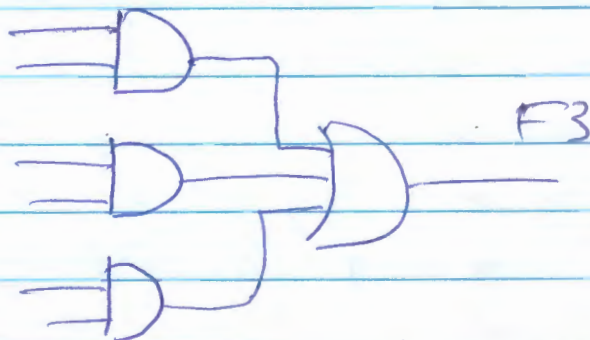
$$F3 = AB + C(D + E)$$



3-level implementation

→ change nonstandard form to a standard form

$$\Rightarrow F3 = AB + CD + CE$$




2-level implementation
(better).

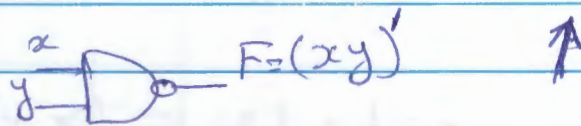
* Digital Logic Gates

① AND 

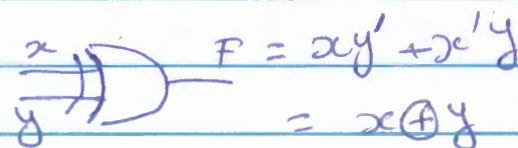
② OR 

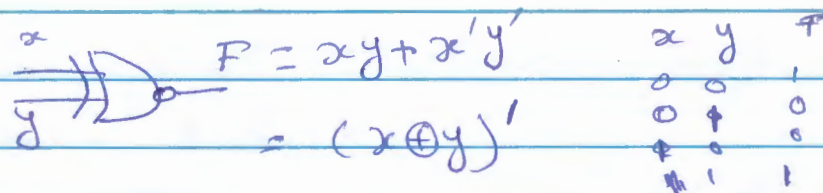
③ Inverter 

④ Buffer 

⑤ NAND  $F = (xy)'$ \uparrow

⑥ NOR  $F = (x+y)'$ \downarrow

⑦ Exclusive-OR
XOR  $F = xy' + x'y$
 $= x \oplus y$

⑧ Exclusive-Nor
or
equivalence  $F = xy + x'y'$
 $= (x \oplus y)'$

x	y	F
0	0	1
0	1	0
1	0	0
1	1	1

* notes

$$(x \downarrow y) \downarrow z \neq x \downarrow (y \downarrow z)$$

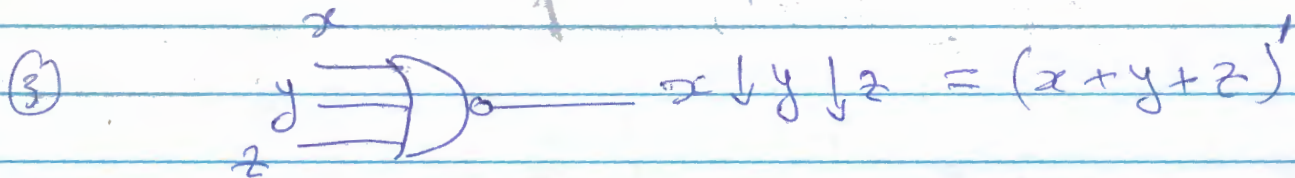
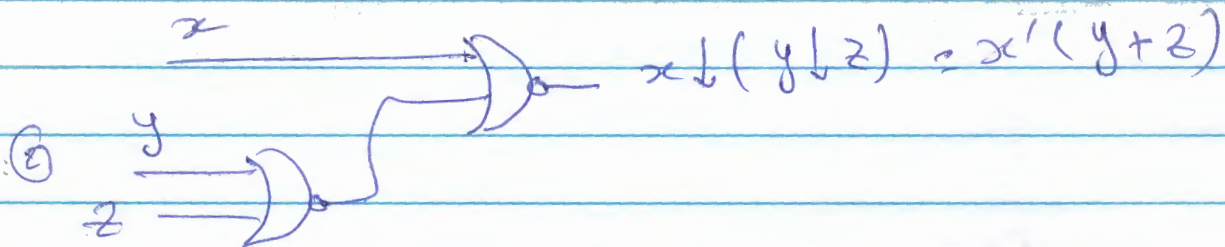
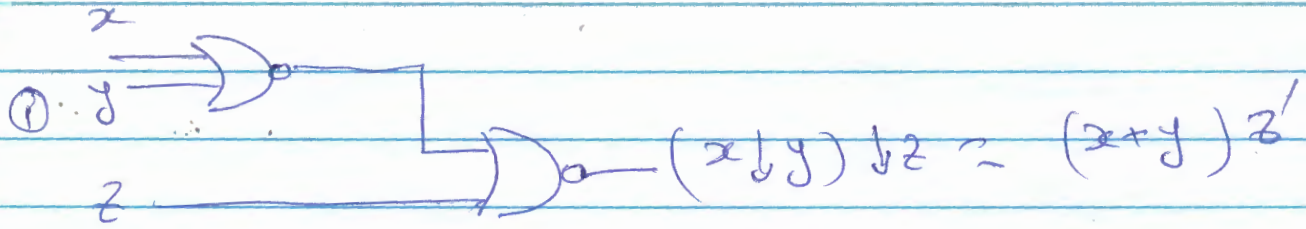
$$[(x \downarrow y)' + z]' \neq [x + (y + z)']'$$

$$xz' + yz' \neq x'y + x'z$$

\Rightarrow By definition

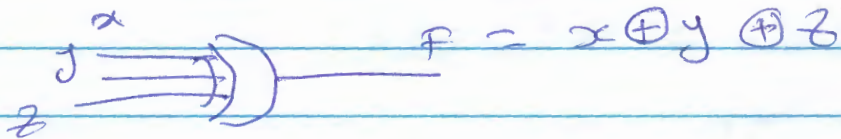
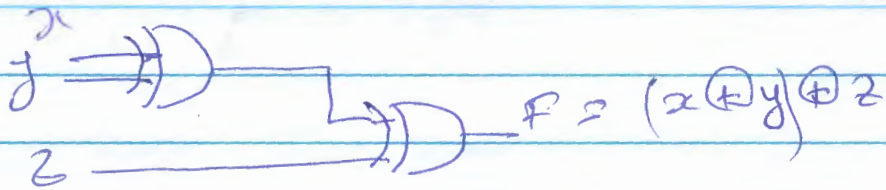
$$x \downarrow y \downarrow z = (x + y + z)'$$

$$x \uparrow y \uparrow z = (xyz)'$$



XOR

$$x \oplus (y \oplus z) = (x \oplus y) \oplus z = x \oplus y \oplus z$$



CHAPTER 3

⊕ Two variable map

⊕ three variable map

⊕ four variable map

$wx \backslash yz$	00	01	11	10
00				
01				
11				
10				

- If we have n variables
- Two adjacent squares represent a term of all variables except 1
- 4 adjacent squares ' ' ' ' variables except 2 \Rightarrow No. of literals = $n-2$
- 2^m adjacent squares ' ' ' ' variables except m \Rightarrow No. of literals = $n-m$

Ex: $F(w, x, y, z) = \sum (0, 1, 2, 4, 5, 6, 8, 9, 12, 13, 14)$

$w \backslash yz$	00	01	10	11
00	1	1	1	1
01	1	1	1	1
11	1	1	1	1
10	1	1	1	1

w	x	y	z	F
0	0	0	0	1
0	0	0	1	1
0	0	1	0	1

$$\Rightarrow F = y' + w'z' + xz'$$

~~Prime~~

$w \backslash yz$	00	01	11	10
00	1	1	1	1
01	1	1	1	1
11	1	1	1	1
10	1	1	1	1

$$F = y' + w'z' + xz'$$

④ Prime Implicants

- A prime implicant is a product term obtained by combining the maximum possible number of adjacent squares in the map.
- If a minterm in a square is covered by only one prime implicant, that prime implicant is said to be essential.

Ex.

$$F(A, B, C, D) = \sum (0, 2, 3, 5, 7, 8, 9, 10, 11, 13, 15)$$

CD \ AB	00	01	11	10
00	1		1	1
01		1	1	
11		1	1	
10	1	1	1	1

$$\begin{aligned} F &= BD + B'D' + CD + AD \\ &= BD + B'D' + CD + AB' \\ &= BD + B'D' + B'C + AD \\ &= BD + B'D' + B'C + AB' \end{aligned}$$

Ex. $F(A, B, C) = \sum (0, 1, 2, 5, 7)$

A \ Bc	00	01	11	10
0	1	1		1
1		1	1	

$$\begin{aligned} F &= A'C' + AC + A'B' \\ &= A'C' + AC + B'C \end{aligned}$$

⊕ Five Variables Map

AB \ CDE	000	001	011	010	110	111	101	100
00								
01								
11								
10								

Ex.

$$F(A, B, C, D, E) = \sum (0, 2, 4, 6, 9, 13, 21, 23, 25, 29, 31)$$

AB \ CDE	000	001	011	010	110	111	101	100
00	1			1	1			1
01		1					1	
11		1				1	1	
10						1	1	

$A = 0$

BC \ DE	00	01	11	10
00	1			1
01	1			1
11		1		
10		1		

$A = 1$

BC \ DE	00	01	11	10
00				
01		1	1	
11		1	1	
10		1		

$$F = ACE + A'B'E' + BD'E$$

⊕ Product of Sums Simplification

$$F(A, B, C, D) = \sum (0, 1, 2, 5, 8, 9, 10)$$

CD AB	00	01	11	10
00	1	1		1
01	1	1		1
11				
10	1	1		1

$$F = B'C' + B'D' + A'C'D$$

$$F' = AB + CD + BD'$$

$$\Rightarrow F = (A' + B')(C' + D')(B' + D)$$

⊕ Don't Care Conditions

Sometimes, for some combinations of the inputs, the output ^{has} not specified value, this is called don't care condition and designated by 'X'.

- Don't care conditions help to simplify the function.

Ex. $F(w, x, y, z) = \sum (1, 3, 7, 11, 15)$

$d(w, x, y, z) = \sum (0, 2, 5)$

$w \backslash yz$	00	01	11	10
00	X	1	1	X
01		X	1	
11			1	
10			1	

$$F = yz + w'x'$$

$$F = yz + w'z$$

* note: they don't have an identical truth table

Ex. $F(w, x, y, z) = \sum (1, 3, 7)$

$d(w, x, y, z) = \sum (0, 5)$

$w \backslash yz$	00	01	11	10
00	X	1	1	
01		X	1	
11				
10				