

Lists

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List

☐ What is a list?

An ordered sequence of elements A1, A2, ..., AN

- ☐ Two types of implementation:
 - 1.Array-Based
 - 2.Linked List

List

- **□** Common operations:
 - > Insert
 - > Find
 - Delete
 - Print list
 - > Print element
 - > IsEmpty
 - > IsLast
 - > FindPrevious
 - > First
 - > Kth
 - > Last



Lists: Array-Based Implementation

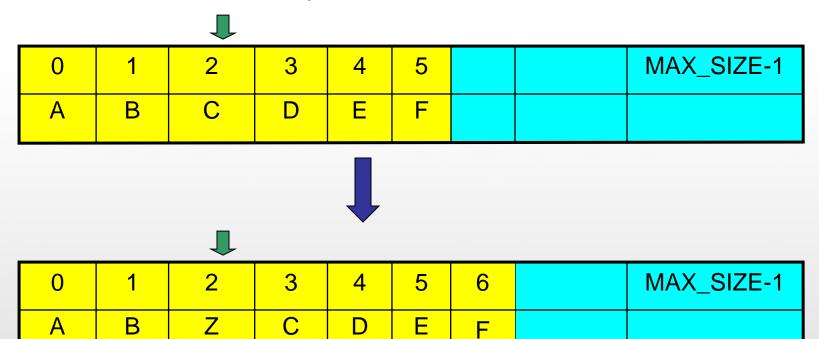
Basic Idea:

Pre-allocate a big array of size MAX_SIZE
Keep track of current size using a variable currentsize
Shift elements when you have to insert or delete

0	1	2	3	 currentSize-1	MAX_SIZE-1
A ₁	A ₂	A_3	A_4	 A _N	

Lists: Array-Based Implementation

Insert Z in kth position



Lists: Array-Based Implementation

Running time for N elements?

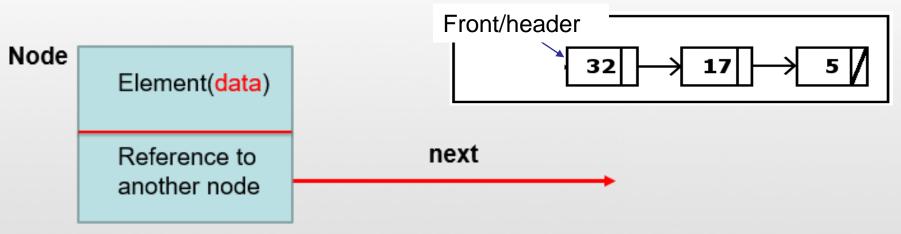
Worst case is insert at position 0. Must move all N items one position before the insert

This is O(N) running time. Probably too slow

Lists: Linked List Implementation

Definition: A **linked list** is a collection of **nodes** that together form a linear.

node: A compound object that stores the contents of the node (Element) and a pointer or reference to the next node in the list (next).

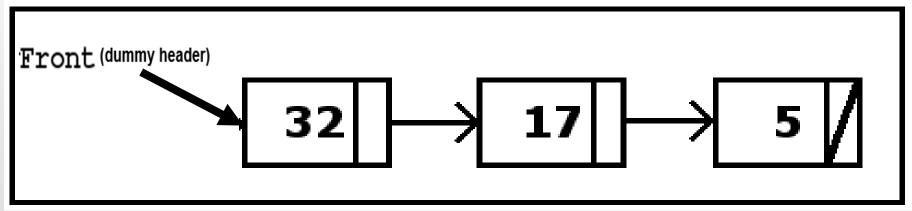


Node implementation

```
/* Stores one element of a linked list. */
public class Node {
  public Object element; // data
  public Node next;
  public Node(Object element) {
    this (element, null);
  public Node(Object element, Node next) {
    this.element = element;
    this.next = next;
```

linked list

- a list implemented using a linked sequence of nodes
- the list only needs to keep a reference to the first node (we might name it Front)
- in Java: java.util.LinkedList (but we'll write our own)

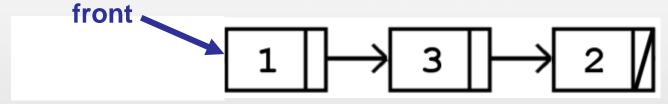


Let's examine sample chains of nodes together, and try to write the correct code for each.

Each Node stores an Integer object

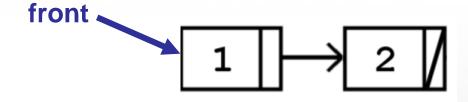


After:

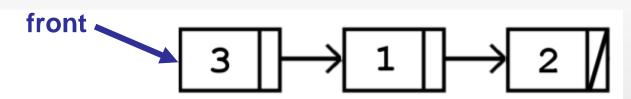


front.next = new Node(new Integer(3), front.next);

Before:



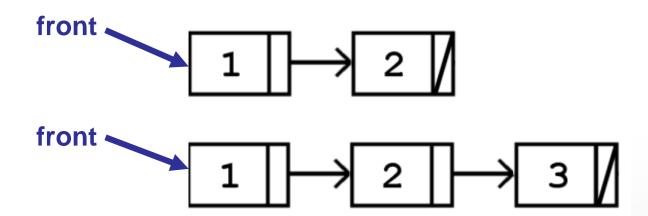
After:



front = new Node(new Integer(3), front);

Before:

After:



front.next.next = new Node(new Integer(3));

Before:

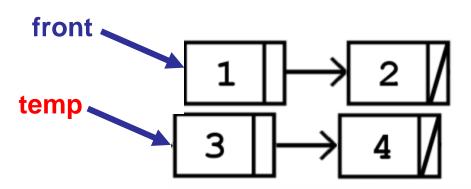
temp 2

After:

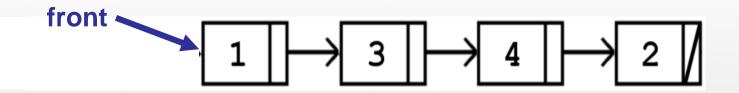
front
$$1 \rightarrow 3 \rightarrow 2$$

temp.next = front.next; // 3 -> 2 front.next = temp; // 1 -> 3

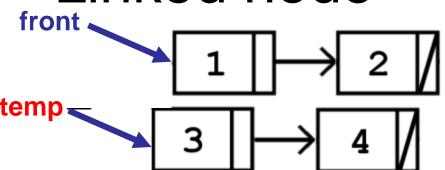
Before:



After:



temp.next.next = front.next; // 4 -> 2 front.next = temp; // 1 -> 3



Before:

• After: 1 3 2 4

front.next.next = temp.next; // 2 -> 4 temp.next = front.next; // 3 -> 2 front.next = temp; // 1 -> 3

Before:

front $2 \rightarrow 3$ front $2 \rightarrow 3$

After:

```
Node temp = front; // temp -> 1
front = front.next; // front -> 2
temp.next = front.next; // 1 -> 3
front.next = temp; // 2 -> 1
```

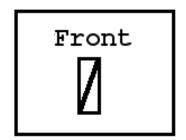
Linked list implementation

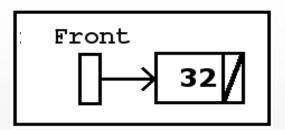
```
/* Models an entire linked list. */
public class LinkedList {
  private Node Front; // Dummy header
  public LinkedList() {
    Front = null;
  /* Methods go here */
```

You have to know

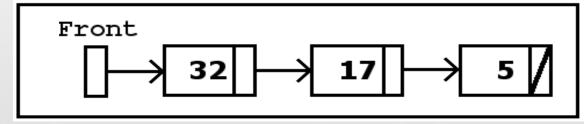
empty list

list with one element





 list with many elements



Analysis of LinkedList runtime

<u>OPERATION</u>	RUNTIME (Big-Oh)
add to start of list	0(1)
add to end of list	O(n)
add at given index	O(n)
clear	0(1)
get	O(n)
find index of an object	O(n)
remove first element	0(1)
remove last element	O(<i>n</i>)
remove at given index	O(<i>n</i>)
set	O(<i>n</i>)
size	O(<i>n</i>)
toString	O(<i>n</i>)

List Implementation (Array VS Linked List)

Operation	Array Complexity	Singly-linked list Complexity		
Insert at beginning	O(n)	O(1)		
Insert at end	O(1)	O(n)		
Insert at middle*	O(n)	O(n)		
Delete at beginning	O(n)	O(1)		
Delete at end	O(1)	O(n)		
Delete at middle*	O(n): O(1) access followed by O(n) shift	O(n): O(n) search, followed by O(1) delete		
Search	O(n) linear search	O(n)		
Indexing: What is the element at a given position k?	O(1)	O(n)		
Print list	O(n)	O(n)		
size	O(1)	O(n)		

Is there way to optimize the operations that in red color?

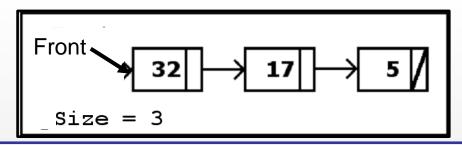
* middle: neither at the beginning nor at the end

List Implementation (Array VS Linked List) optimization

problem: array list has a O(1) size method, but the linked list needs O(n) time

solution: add a Size field to the linked list

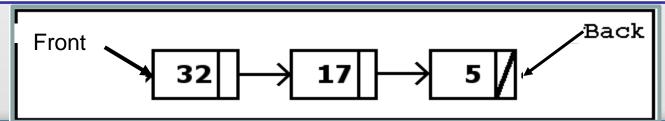
Big-Oh runtime for size improve to O(1)



problem: array list has O(1) get/remove of last element, but the linked list needs O(n)

solution: add a Back pointer to the last node

Big-Oh runtime improve the get/remove of last element to O(1)



Linked list implementation

```
/* Models an entire linked list. */
public class LinkedList {
  private Node Front, Back;
  private int Size;
  public LinkedList() {
    Front = Back = null;
    Size=0;
  /* Methods go here */
```

Linked list implementation

```
void addFirst(Object o)
  Inserts the given element at the beginning of the list. */
public void addFirst(Object element) {
  Node newNode;
  newNode=new Node(element);
  if (Size==0) {// Empty List
    Front=Back=newNode;
  else{
    newNode.next=Front;
    Front=newNode;
   Size++;// update Size
```