

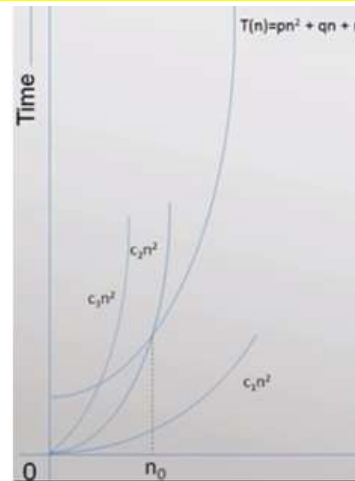


The Big-O Notation

Assume the order of time of an algorithm is a **quadratic** time as displayed in the graph. Our job is to find an **upper bond** for this function $T(n)$. Consider a function c_1n^2 ← never over take $T(n)$

c_2n^2 such that its greater than $T(n)$ for $n > n_0$. In this case we say that c_2n^2 is an upper bond of $T(n)$

But we can come up with many functions satisfy this condition. We need to be precise.



OBSERVATIONS:

$$\forall n > n_0$$

$$c_2n^2 \geq pn^2 + qn + r$$

Big Oh $O(n^2)$: $f(n)$: there exist positive constants c and n_0 such that $0 \leq f(n) \leq cn^2$ for all $n \geq n_0$

In general

$O(g(n))$: $f(n)$: there exist positive constants c and n_0 such that $0 \leq f(n) \leq cg(n)$ for all $n \geq n_0$

Example 1:

$$5n^2 + 6 \in O(n^2) \quad ??? \quad \checkmark$$

$$\text{Find } cn^2 \rightarrow c=6 \text{ and } n_0=3$$

$$\rightarrow c=5.1 \quad n_0=8$$

Example 2:

$$5n + 6 \in O(n^2) \quad ??? \quad \checkmark$$

$$\text{Find } cn^2 \rightarrow c=11 \text{ and } n_0=1$$

Example 3:

$$n^3 + 2n^2 + 4n + 8 \in O(n^2) \quad ??? \quad \times$$

$$\text{Find } cn^2 \geq n^3 + 2n^2 + 4n + 8 \quad ??? \quad \times$$

$$a_m n^m + a_{m-1} n^{m-1} + \dots + a_0 \in O(n^m)$$

$$\log n \leq \sqrt{n} \leq n \leq n \log n \leq n^2 \leq n^3 \leq 2^n \leq n!$$

What does it mean?



