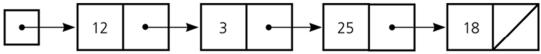


Linked List

Algorithm - abstract way to perform computation tasks **Data Structure** - abstract way to organize information



Linked List:

head

Node:



Node code:

```
public class Node<T> {
    private T data;
    private Node<T> next;

public Node(T data) { this.data = data; }

public void setData(T data) { this.data = data; }

public T getData() { return data; }

public Node<T> getNext() { return next; }

public void setNext(Node<T> next) { this.next = next; }
}
```

Linked List Code:

```
public class LinkedList<T> {
      private Node<T> head;
}
```

Inserting a new node:

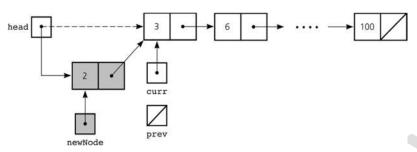
Inserting a Node into a Specified Position of a Linked List:

Three steps to insert a new node into a linked list

- Determine the point of insertion
- Create a new node and store the new data in it
- Connect the new node to the linked list by changing references

Case 1: To insert a node at the beginning of a linked list: (curr == head)

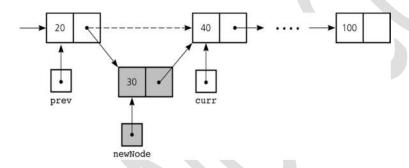
newNode.next = head; head = newNode;



What's the time complexity of inserting an item to the head?? \rightarrow **O(1)**

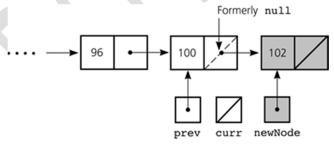
Case 2: To insert a node between two nodes:

newNode.next = curr; prev.next = newNode;



Case 3: Inserting at the end of a linked list is a special case if curr is null:

newNode.next = curr; prev.next = newNode;



Time Complexity \rightarrow O(n)

H.W. → implement insert into a sorted linked list

Determining curr and prev

Determining the point of insertion or deletion for a sorted linked list of objects

for (prev = null , curr = head;

(curr != null) && (newValue.compareTo(curr.item) > 0);

```
prev = curr, curr = curr.next); // end for
```

Create a driver class to test linked list classes.

Override the *toString* methods first

Node toString:

@Override

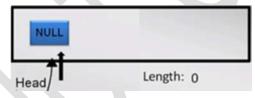
public String toString() { return data.toString(); }

LinkedList toString:

@Override

```
public String toString() {
    String res = "→";
    Node<T> curr = head;
    while (curr != null) {
        res += curr + "→";
        curr = curr.next;
    }
    return res + "NULL";
}
```

Length of Linked List?



Case 1: If it's empty:

Case 2: If not: Make a pointer and move over all the nodes and maintain a counter



Length code: Time Complexity \rightarrow O(n)

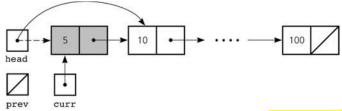
```
public int length() {
    int length = 0;
    Node<T> curr = head;
    while (curr != null) {
        length++;
        curr = curr.next;
    }
    return length;
}
```



Data Structure: Lectures Note

Deleting Nodes:

Case 1: Deleting the head node:



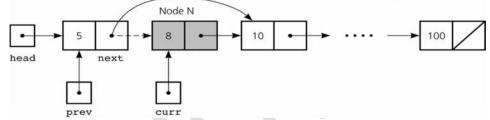
Simply move the **head** to the **head.next**: head = head.next;

Now first Node has no reference to it → Garbage

Time Complexity \rightarrow **O(1)**

```
Delete at head code: // make sure linked list is not empty
public Node<T> deleteAtStart() {
         Node<T> toDel = head;
         head = head.next;
         return toDel;
```

Case 2: Delete node N which curr references:



Set **next** in the node that precedes **N** to reference the node that follows **N**prev.next = curr.next; // prev.next = prev.next.next;

Searching for an Item in a Linked List:



Time Complexity: linear growth \rightarrow O(n)

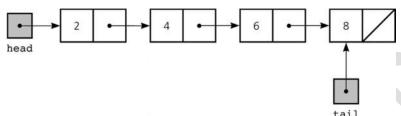
Find code:



Variations of the Linked List:

- 1- Tail References (Doubly Ended Linked List)
 - Remembers where the end of the linked list is.
 - Therefore, we can add and delete at both ends.
 - To add a node to the end of a linked list

```
tail.next = new Node(request, null);
```



```
public class DoubleEndedList<T> extends LinkedList<T> {
    private Node<T> tail;
    public Node<T> getTail() {            return tail; }

    public void addAtEnd(T data) {
        Node<T> newNode = new Node<T>(data);
        if (head == null) { // empty
            head = newNode;
            tail = newNode;
        }
        else {
            tail.setNext(newNode);
            tail = newNode;
        }
    }
}
```

Make sure to override addAtStart to set the tail pointer correctly:

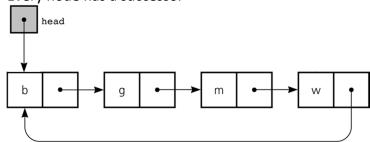
```
@Override
public void addAtStart(T data) {
    Node<T> newNode = new Node<T>(data);
    if (head == null) { // empty
        head = newNode;
        tail = newNode;
    }
    else{
        newNode.setNext(head);
        head = newNode;
    }
}
```



Data Structure: Lectures Note

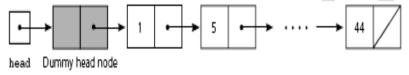
2- Circular Linked List

- Last node references the first node
- Every node has a successor



3- Dummy Head Nodes

- Always present, even when the linked list is empty
- Insertion and deletion algorithms initialize **prev** to reference the dummy head node, rather than **null**



Processing Linked Lists Recursively:

- Traversal
 - Recursive strategy to display a list
 Write the first node of the list
 Write the list minus its first node

```
public static void traversList(Node curr) {
    if(curr == null)
        System.out.println("NULL");
    else {
            System.out.print("[" + curr + "]-->");
            traversList(curr.next);
        }
}
```

- Recursive strategies to display a list backward
 - writeListBackward strategy

Write the last node of the list

Write the list minus its last node backward

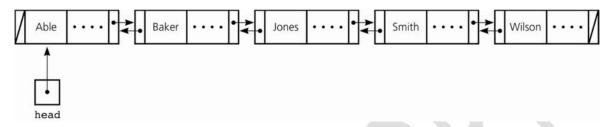
```
public static void traversListBackward(Node curr) {
    if(curr == null)
        System.out.print("NULL");
    else {
        traversListBackward(curr.next);
        System.out.print("<--[" + curr + "]");
    }
}</pre>
```

Doubly Linked List

Node:



Doubly Linked List: Each node references both its predecessor and its successor:



Doubly Node Code:

```
public class DNode <T extends Comparable<T>>{
    T data;
    DNode next;
    DNode prev;

public DNode(T data) { this.data = data; }
    public T getData() { return data; }
    public DNode getNext () { return next; }
    public DNode getPrev () { return prev; }

public void setNext(DNode next) { this.next = next; }
    public void setPrev(DNode prev) { this.prev = prev; }
    public String toString() { return this.data.toString(); }
}
```

Doubly Linked List code:

```
public class DLinkedList <T extends Comparable<T>>{
    DNode head;
}
```

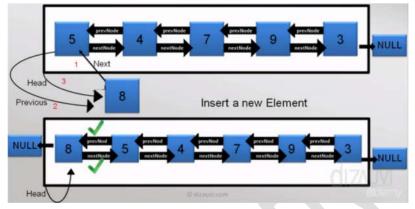
Override toString method code:

```
public String toString() {
   String res = "Head-->";
   DNode<T> curr = this.head;
   while (curr != null) {
     res += "["+curr + "]";
     curr = curr.getNext();
     if(curr!=null)
        res += "<=>";
   }
   return res + "-->NULL";
}
```



Data Structure: Lectures Note Insert a new node (not sorted)

Case 1: Insert at head:



```
public void insertAtHead(T data) {
  DNode<T> newNode = new DNode(data);
  if(head==null) // empty linkedlist
    head = newNode;
  else {
    newNode.setNext(this.head);
    head.setPrev(newNode);
    head = newNode;
 }
```

Case 2: Insert at end:

Student Activity: insert at last

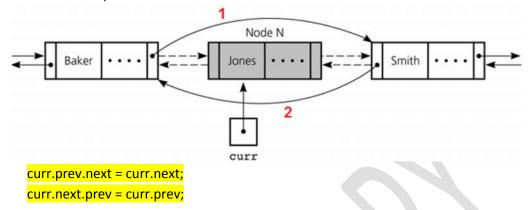
```
public void insertAtEnd(T data) {
  DNode<T> newNode = new DNode(data);
  if (head == null) // empty linkedlist
    head = newNode;
  else { // find last node
    DNode<T> last = head;
    while(last.getNext() != null)
      last = last.getNext();
    last.setNext(newNode);
    newNode.setPrev(last);
```

Length of a doubly linked list code:

```
public int length() {
  int length = 0;
  DNode<T> curr = this.head;
  while (curr != null) {
    length++;
    curr = curr.getNext();
  return length;
```

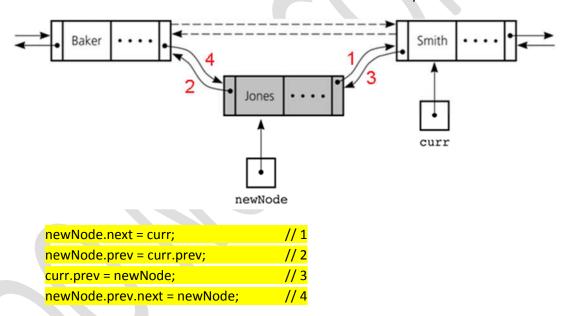
Delete a node:

• To delete the node that curr pointer references

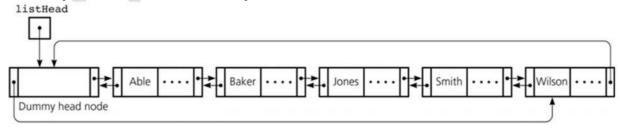


Insert a new Node (Sorted):

• To insert a new node that newNode references before the node referenced by curr



Circular doubly linked list with dummy head:



- Preceding reference of the dummy head node references the last node.
- next reference of the last node references the dummy head node.
- Eliminates special cases for insertions and deletions.

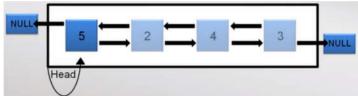


Case Study: Insertion Sort using doubly linked list (Using NO extra space):

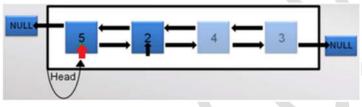
Review insertion sort logic and point to problem of insertion and time needed to shift the items

Worst case if the array is reverse sorted

Example: assume we need to sort the following doubly linked list:



Assumption: 1st node is sorted. We start from the 2nd element:



Here:

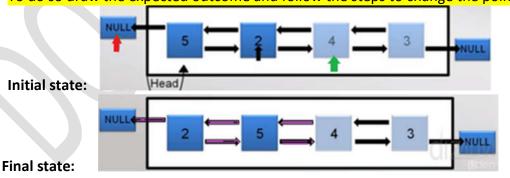
- The **black** pointer points to the **current** node to be sorted.
- The red pointer points to previous node of current node to be sorted.
- The green pointer points to next node of current node to be sorted.

Step 1: The **red** pointer keeps move backward until it reaches a node which has a value **smaller** than the **current** node **Or** reach **NULL**.

Step 2: the current item will be inserted after red pointer as follow:

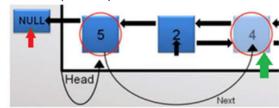
Make sure you maintain references correctly.

To do so draw the expected outcome and follow the steps to change the pointers:



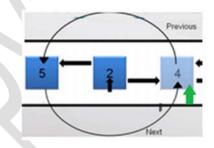
Case 1: insert to head

Step 2.0: make new green pointer = black.next

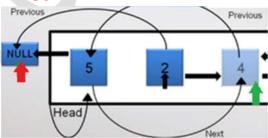


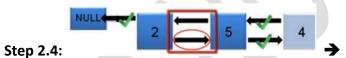


if (green != null) green.prev = black.prev

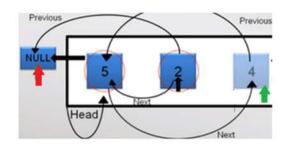








if(red==null) black.next = black.next.prev else black.next = red.next

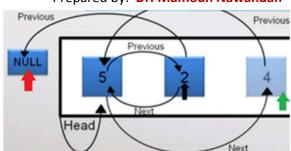




else red.next.prev = black

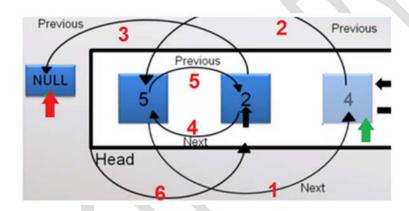






Step 2.6:

if (red == NULL) head = black
else red.next = black;

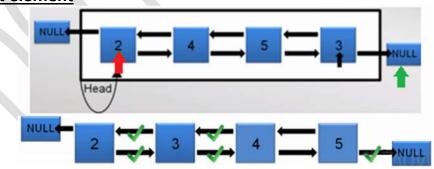


Step 2.7: black = green

Case 2: insert 4 in the middle

Practice yourself

Case 3: insert last element





Insertion Sort Code:

```
// Insertion Sort of a Doubly Linked List
public void sort() {
  DNode black = head.next;
  while (black != null) {
    DNode red = black.prev;
    while (red != null && (red.data.compareTo(black.data) > 0 )) // step 1.0
       red = red.prev;
    DNode green = black.next;
                                         // step 2.0
    if (red != null | | (head != black)) {
       black.prev.next = green;
                                        // step 2.1
       if (green!= null) {
                                        // step 2.2
         green.prev = black.prev;
       black.prev = red;
                                    // step 2.3
    if (red == null) { // set the black as head
       if (head != black) {
         black.next = head;
                                     // step 2.4
         black.next.prev = black;
                                      // step 2.5
         head = black;
                                   // step 2.6
       }
    } else { // red is not null
       black.next = red.next;
                                       // step 2.4
       red.next.prev = black;
                                      // step 2.5
       red.next = black;
                                    // step 2.6
    black = green; // step 2.7
  }
```